Midterm 2: Compsci 201 Form C

Prof. Astrachan

October 30, 2024

Name: _____

netid: _____

In submitting this test, I affirm that I have followed the Duke Community Standard.

Community standard acknowledgement (signature)

The definition of **ListNode** unless another is specified. This is the same definition as used in linked-list APT problems.

```
public class ListNode {
    int info;
    ListNode next;
    ListNode(int x) {
        info = x;
    }
    ListNode(int x, ListNode node) {
        info = x;
            next = node;
    }
}
```

This Exam is Form C, please mark your answer sheet accordingly

The method create below is intended for the call create(10) to return a list $1 \rightarrow 2 \rightarrow \ldots \rightarrow 9 \rightarrow 10$, but instead the call returns a list with a single node containing zero. In general, the call create(N) is intended to return a linked list of N nodes in order $1 \rightarrow 2 \rightarrow \ldots \rightarrow N$.

```
25      public ListNode create(int n){
26         ListNode list = null;
27         while (n >= 1){
28              list = new ListNode(0,null);
29              n -= 1;
30         }
31         return list;
```

PROBLEM 1:

As written, what is the complexity of the call create(N)?

- **A.** *O*(1)
- **B.** O(N)
- C. $O(N^2)$

PROBLEM 2:

What statement can replace the right hand side of the assignment statement on line 28 so that the code works as intended? The only modification you can make is to the right hand side of line 28.

Write the new code in the appropriate fill-in-the blank area of the back of the bubble answer sheet. Bubble A for this question on the front of the answer sheet.

PROBLEM 3:

An attempt to write a correct, recursive version to solve this problem yields the code below, it does **not work as intended**.

Which of the following is the best description of the result returned by the call create(10)?

- A. It creates 10 nodes, but returns a pointer to one node which contains 10 and whose next field is null.
- **B.** It creates 10 nodes, but returns a pointer to one node which contains 1 and whose next field is null.
- **C.** It creates one node, and returns a pointer to that single node which contains 10 and whose next field is null.
- **D.** It creates 10 nodes and returns a pointer to a node containing 10, which is the first node of a linked list: 10->9->8->...->2->1->null

The *ListSum* APT was a required APT to write a method sum that adds those node values in a linked list that are greater than a threshold parameter. For example, if list = 6 -> 10 -> 7 -> 8 the call sum(list,7) returns 10+8 == 18 and the call sum(list,10) returns 0 (no values in the list are greater than 10.

In the code below, method sumIter is all green if the name is changed to sum. The method sum is a recursive method that is all green if two statements are added.

```
public class ListSum {
 1
 2
        public int sumIter(ListNode list, int thresh) {
 3
             int total = 0;
 4
             while (list != null){
 5
                 if (list.info > thresh) {
 6
                      total += list.info;
 7
                 }
 8
                 list = list.next;
 9
10
             return total;
11
12
        public int sum(ListNode list, int thresh) {
13
             if (list == null) return 0;
14
             int total = 0;
15
             if (list.info > thresh) {
16
                 // assignment statement
17
             }
18
             return expression with recursive call;
19
         3
```

PROBLEM 4:

If sum is all green, is there a linked list list, with one or more nodes all of whose node values are greater than zero, such that the value of the call sum(list,sum(list,0)) is greater than zero?

- A. Yes, there is such a list with one or more positive values
- **B.** No, the value returned by the call is zero for any linked list of one or more positive values.

PROBLEM 5:

If sum is all green, and the value of sum(list,thresh) is N > 0, what is the value of sum(list,N)

- **A.** 0
- **B.** O(N)
- C. it cannot be determined

PROBLEM 6:

Write the assignment statement for line 16, so the code is all green, in the appropriate fill-in-the blank area of the back of the bubble answer sheet. Bubble A for this question on the front of the answer sheet.

PROBLEM 7:

Write the expression that includes a recursive call for line 18, so the code is all green, in the appropriate fill-in-the blank area of the back of the bubble answer sheet. Bubble A for this question on the front of the answer sheet.

Part of the code in a working version of HashMarkovModel from the P3-Markov assignment is shown below.

```
public class HashMarkovModel extends AbstractMarkovModel {
 3
 4
        private HashMap<WordGram,List<String>> myMap;
 5
 6
        public HashMarkovModel(int order){
 7
             super(order);
             myMap = new HashMap<>();
 8
 9
        }
10
        public HashMarkovModel(){
11
             this(3);
         }
12
```

PROBLEM 8:

Execution of the statement HashMarkovModel mod = new HashMarkovModel(); results in running the code on line 11 above (because of the call of the default constructor).

Does executing line 11 and the code that line 11 then calls result in creating a HashMap<> object?

- A. No, the value of myMap will be null, perhaps until setTraining is called.
- B. Yes, myMap will be initialized to point to/reference a HashMap because of the other constructor.

PROBLEM 9:

The code below for method getRandomNextWord is the same code that runs in both BaseMarkovModel and HashMarkovModel

```
41
     public String getRandomNextWord(WordGram wgram) {
42
         List<String> follows = getFollows(wgram);
43
         if (follows.size() == 0) {
44
             return END_OF_TEXT;
         }
45
         else {
46
             int randomIndex = myRandom.nextInt(follows.size());
47
             return follows.get(randomIndex);
48
49
         }
50
     }
```

Which one of the following is true regarding getRandomNextWord when the models have been trained on a text of N words?

- A. The code has the same big-Oh complexity in both HashMarkovModel and BaseMarkovModel
- **B.** The code in BaseMarkovModel is O(N) and the code in HashMarkovModel is O(1)
- **C.** The code in HashMarkovModel is O(N) and the code in BaseMarkovModel is O(1)
- **D.** The code in BaseMarkovModel is O(N) and the code in HashMarkovModel is $O(\log N)$

PROBLEM 10:

In HashMarkovModel, which of the following best characterizes the Strings stored in the List<String> returned by the call myMap.get(wg) where wg is a WordGram? The only call is in the code shown below.

```
28 @Override
29 public List<String> getFollows(WordGram wgram) {
30 if (myMap.containsKey(wgram)) {
31 return myMap.get(wgram);
32 }
33 return new ArrayList<>();
34 }
```

- A. The Strings are those that follow an occurrence of wg in the training text for the model.
- **B.** The Strings are those that occur at indexes k, 2k, 3k, ... in the array myWords where k is the order of the model.
- C. The Strings are those that are essentially removed each time wg.shiftAdd(s) is called to create a new WordGram.

PROBLEM 11:

Both BaseMarkovModel and HashMarkovModel include the code below in the setTraining method.

myWords = text.split("\\s+");

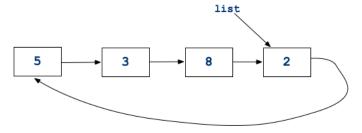
This one line of code is the *only* line in BaseMarkMode.setTraining, but is followed by code to populate a HashMap<WordGram,List<String>> object in HashMarkovModel.setTraining as shown:

15		@Override
16	•	<pre>public void setTraining(String text) {</pre>
17		<pre>myWords = text.split("\\s+");</pre>
18		myMap.clear();
19		<pre>WordGram current = new WordGram(myWords, 0, myOrder);</pre>
20		<pre>for(int k=my0rder; k < myWords.length; k++) {</pre>
21		<pre>String next = myWords[k];</pre>
22		<pre>myMap.putIfAbsent(current,new ArrayList<>());</pre>
23		<pre>myMap.get(current).add(next);</pre>
24		<pre>current = current.shiftAdd(next);</pre>
25		}

Which one statement is true?

- A. Calling setTraining(str) has the same big-Oh complexity in both HashMarkovModel and BaseMarkovModel, essentially the number of space-separated strings in str.
- **B.** Calling setTraining(str) has different big-Oh complexity in the different models: the call is *more efficient* in BaseMarkovModel.
- C. Calling setTraining(str) has different big-Oh complexity in the different models: the call is *more efficient* in HashMarkovModel.

In a *circularly-linked* list, the last node of the list points to the first node rather than pointing at null. Such lists are typically represented by a pointer to the last node as shown, so that list references the last node and list.next references the first node. The diagram below represents the list [5,3,8,2] with 2 being the last node and 5 the firstnode.



The method makeCircular below returns a pointer to the last node of a circulary-linked list so that the call ListNode list = makeCircular(new int[]{5,3,8,2}); returns a pointer to the last node of the list pictured above.

```
15
         public ListNode makeCircular(int[] a){
16
17
             ListNode first = null;
18
             ListNode last = null;
19
20
             for(int k=a.length-1; k >= 0; k--){
21
                 first = new ListNode(a[k],first);
22
                 if (last == null){
23
                     last = first;
24
                 3
25
             last.next = first;
26
             return last;
27
28
         3
```

PROBLEM 12:

In the code above, if parameter a has N values, so a.length == N, how many times is line 23 executed?

- A. never
- **B.** one time
- **C.** O(1) times, but greater than one
- **D.** O(N) times

PROBLEM 13:

What is returned by the call makeCircular(new int[]{})? Note that the parameter passed to makeCircular is an array with zero-length.

A. null

- B. nothing, a null pointer exception is thrown
- C. a list with one node storing a zero

PROBLEM 14:

The method circular2String below correctly returns a String representing its circularly-linked list parameter. The call circular2String(list) for the list at the beginning of this problem (and reproduced below) returns "5 3 8 2".

```
5
         3
                          2
                  8
      public String circular2String(ListNode list){
30
31
          if (list == null) return "";
32
33
34
          ListNode anchor = list;
35
          list = list.next;
                                    // now first node
36
37
          StringBuilder sb = new StringBuilder();
38
          sb.append(""+list.info);
39
40
          while (list != anchor){
41
              list = list.next;
42
              sb.append(" "+list.info);
43
          }
44
          return sb.toString().trim();
45
```

list

If the two lines 41 and 42 are interchanged, so that line 42 is list = list.next, what will be returned for the same call circular2String(list)?

A. "5 3 8 2"
B. "5 5 3 8 2"
C. "5 5 3 8"

PROBLEM 15:

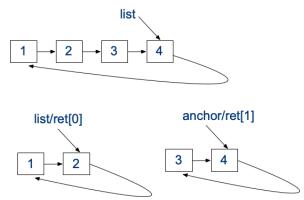
Which the following is a valid reason to use a StringBuilder with append in the code above rather than a String with + to concatenate values when parameter list has N nodes?

- A. StringBuilder results in O(N) time, but using a String would run in $O(N^2)$ time.
- **B.** StringBuilder results in O(N) time, and String also runs in O(N) time, but with a larger coefficient.

The method **split** is intended to split a circularly-linked list into two circularly-linked lists, returning pointers to the last nodes of two lists by re-arranging pointers.

To return two pointers, an array with two values is returned. If parameter list has 2N nodes, the call ListNode[] ret = split(list); should have both ret[0] and ret[1] referencing the last nodes of two lists of N nodes; ret[0] is a circulary-linked list with the first half of the nodes in list and ret[1] is a circularly-linked list with the second half of the nodes.

The diagram shows what's intended for the circularly-linked list [1,2,3,4]



The code below works *only* for lists with an even number of nodes when two statements are added after the loop.

```
47
       public ListNode[] split(ListNode list){
48
49
           ListNode anchor = list;
50
           list = list.next;
                                             // first node
51
           ListNode fast = list.next;
                                            // second node
           while (fast != anchor){
52
53
               list = list.next;
54
               fast = fast.next.next;
           }
55
           ListNode secondFirst = list.next;
56
57
           // statement 1
58
           // statement 2
59
           return new ListNode[]{list,anchor};
60
61
```

As shown in the diagram, list references the node containing 2 and anchor references the node containing 4 when the loop exits.

PROBLEM 16:

What is the complexity of split(list) when list has N nodes if the two missing statements after the loop run in O(1) time?

- A. $O(\log N)$
- **B.** O(N)
- C. $O(N^2)$
- **D.** $O(N \log N)$

```
(the code again)
```

```
47
       public ListNode[] split(ListNode list){
48
           ListNode anchor = list;
49
50
           list = list.next;
                                            // first node
                                            // second node
51
           ListNode fast = list.next;
52
           while (fast != anchor){
               list = list.next;
53
54
               fast = fast.next.next;
55
           }
56
           ListNode secondFirst = list.next;
57
           // statement 1
58
           // statement 2
59
60
           return new ListNode[]{list,anchor};
61
       }
```

PROBLEM 17:

What is the first missing statement, on line 57? It should be an assignment statement and **must reference** variable list to create the circulary-linked list whose last node will be pointed to by list. In writing code assume parameter list has an even number of nodes.

Write the answer in the fill-in-the-blank section on the back of the answer sheet. Bubble A for this question on the front of the answer sheet.

PROBLEM 18:

What is the second missing statement, on line 58? It should be an assignment statement and **must reference** variables **anchor** and **secondFirst** to create the circulary-linked list whose last node will be pointed to by **anchor**. In writing code assume parameter list has an even number of nodes.

Write the answer in the fill-in-the-blank section on the back of the answer sheet. Bubble A for this question on the front of the answer sheet.

A previous APT problem called for writing the method **filter** below which removes those nodes from its linked list parameter whose values are less than their index; where the index of the first node is zero, the index of the second is 1, and so on. Lists are guaranteed to store non-negative values, so the first node is never removed.

For example, the list 2->0->3->1->4 becomes 2->3->4 since the 0 at index 1 is removed and the 1 at index 3 is removed. The node with value 3 has index 2 and the node with value 4 has index 4, so they are not removed.

The list $3 \rightarrow 0 \rightarrow 1 \rightarrow 2$ becomes 3: all nodes but the first node are removed.

The code below is all green when two statements are added in the locations indicated.

```
2
     public ListNode filter(ListNode list) {
 3
         int index = 1;
 4
         ListNode first = list;
 5
         while (list.next != null){
 6
              if (list.next.info < index){</pre>
 7
                  // statement to remove
 8
 9
              else {
10
                  // statement to advance
11
12
              index += 1;
         }
13
         return first;
14
15
     }
```

PROBLEM 19:

Write the code for line 7 that removes a node from the linked list in the appropriate fill-in-the blank area of the back of the bubble answer sheet. Bubble A for this question on the front of the answer sheet.

PROBLEM 20:

When a node is not removed, the next node in the list must be considered. Write the code for line 10, that advances the appropriate pointer so the code is all green, in the appropriate fill-in-the blank area of the back of the bubble answer sheet. Bubble A for this question on the front of the answer sheet.

There are three classes in P3-DNA that implement the IDnaStrand interface: StringStrand, StringBuilderStrand, and LinkStrand. The append method is shown on the right for each of these:

48

49

50

}

PROBLEM 21:

If dna.length() == N, which one of these methods is O(1), the others have a big-Oh complexity that depends on N?

- **A.** StringStrand is O(1)
- **B.** StringBuilderStrand is O(1)
- C. LinkStrand is O(1)

PROBLEM 22:

In the code below, XYZ will be replaced by one of the three stand types. Note that dna.length() == 1000 because dna is aString of 1,000 'a' characters.

```
IDnaStrand strand = new XYZ();
String dna = "a".repeat(1000);
for(int k=0; k < 10000; k++) {</pre>
    strand.append(dna);
}
```

If the code strand.append(dna) executes 10,000 times where dna.length == 1000 as shown, which of strand types executes most slowly: both with big-Oh and empirically.

```
A. StringStrand
```

B. StringBuilderStrand

```
C. LinkStrand
```

StringStrand

```
61
          @Override
 62
          public IDnaStrand append(String dna) {
 63
              myInfo = myInfo + dna;
 64
              myAppends++;
 65
              return this;
 66
StringBuilderStrand
 59
         public IDnaStrand append(String dna) {
 60
             myInfo.append(dna);
              myAppends++;
 61
 62
              return this;
 63
         }
LinkStrand
          @Override
 43
 44
          public IDnaStrand append(String dna) {
 45
              myLast.next = new Node(dna,null);
 46
              myLast = myLast.next;
 47
              myAppends += 1;
```

mySize += dna.length();

return this;

The next two problems are based on the call to cutAndSplice below which has the same runtime if the 100 on line 35 is replaced by 1000, 2000, 3000.

33	<pre>public void dummy(){</pre>
34	int n = 1000;
35	<pre>String splicee = "a".repeat(100);</pre>
36	<pre>IDnaStrand strand = new LinkStrand("cgat".repeat(n));</pre>
37	<pre>IDnaStrand cut = strand.cutAndSplice("cgat",splicee);</pre>
38	}

PROBLEM 23:

Which one of the following is true about a call cut.toString() if that comes after line 37?

- A. It will take the same time regardless of whether line 35 uses 100, 1000,2000,3000
- **B.** It will take more time as the value used on line 35 increases from 100, to 1000,2000,3000

PROBLEM 24:

In the code above there are 1,000 breaks/replacements of the enzyme by the splice because of the value assigned to n on line 34 and used on line 36. If the 1000 is replaced by 2000, 3000, and so on how will the runtime of the call of cutAndSplice on line 37 change?

- A. The runtime will stay the same, it does not depend on the number of breaks.
- **B.** If there are N breaks, the runtime will be O(N)
- **C.** If there are N breaks, the runtime will be $O(N^2)$

PROBLEM 25:

In the code above, if LinkStrand is replaced by StringStrand on line 36 which one of the following is true (assume the value of 1000 on line 34 and 100 on line 35 do not change).

- A. The runtime of cutAndSplice will be about the same as when LinkStrand is used.
- B. The runtime of cutAndSplice will be much greater than when LinkStrand is used.