

SQL: Part III

CPS 116
Introduction to Database Systems

Announcements (September 22) ²

- ❖ Homework #2 due in a week
- ❖ Missing a handout and can't find it on the Web site?
 - Check the handout box outside my office (D327)
- ❖ Midterm exam in class in two weeks
- ❖ Project Milestone #1 due in three weeks

"Active" data ³

- ❖ Constraint enforcement: When an operation violates a constraint, abort the operation or try to "fix" the data
 - Example: enforcing referential integrity constraints
 - Generalize to arbitrary constraints?
- ❖ Data monitoring: When something happens to the data, automatically execute some action
 - Example: When price rises above \$20 per share, sell
 - Example: When enrollment is at the limit and more students try to register, email the instructor

Triggers

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- ❖ A trigger is an event-condition-action (ECA) rule
 - When event occurs, test condition; if condition is satisfied, execute action
- ❖ Example:
 - Event: whenever there comes a new student...
 - Condition: with GPA higher than 3.0...
 - Action: then make him/her take CPS116!

Trigger example

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```
CREATE TRIGGER CPS116AutoRecruit
AFTER INSERT ON Student → Event
REFERENCING NEW ROW AS newStudent
FOR EACH ROW
WHEN (newStudent.GPA > 3.0) → Condition
INSERT INTO Enroll
  VALUES(newStudent.SID, 'CPS116');
  ↓
  Action
```

Trigger options

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- ❖ Possible events include:
 - INSERT ON *table*
 - DELETE ON *table*
 - UPDATE [OF *column*] ON *table*
- ❖ Granularity—trigger can be activated:
 - FOR EACH ROW modified
 - FOR EACH STATEMENT that performs modification
- ❖ Timing—action can be executed:
 - AFTER or BEFORE the triggering event

Transition variables

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- ❖ OLD ROW: the modified row before the triggering event
- ❖ NEW ROW: the modified row after the triggering event
- ❖ OLD TABLE: a hypothetical read-only table containing all modified rows before the triggering event
- ❖ NEW TABLE: a hypothetical table containing all modified rows after the triggering event
- ☞ Not all of them make sense all the time, e.g.
 - AFTER INSERT statement-level triggers
 - Can use only NEW TABLE
 - BEFORE DELETE row-level triggers
 - etc.

Statement-level trigger example

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```
CREATE TRIGGER CPS116AutoRecruit
AFTER INSERT ON Student
REFERENCING NEW TABLE AS newStudents
FOR EACH STATEMENT
INSERT INTO Enroll
(SELECT SID, 'CPS116'
FROM newStudents
WHERE GPA > 3.0);
```

BEFORE trigger example

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- ❖ Never give faculty more than 50% raise in one update
- ```
CREATE TRIGGER NotTooGreedy
BEFORE UPDATE OF salary ON Faculty
REFERENCING OLD ROW AS o, NEW ROW AS n
FOR EACH ROW
WHEN (n.salary > 1.5 * o.salary)
SET n.salary = 1.5 * o.salary;
```
- ☞ BEFORE triggers are often used to “condition” data
  - ☞ Another option is to raise an error in the trigger body to abort the transaction that caused the trigger to fire

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## Statement- vs. row-level triggers 10

Why are both needed?

- ❖ Certain triggers are only possible at statement level
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- ❖ Simple row-level triggers are easier to implement and may be more efficient
  - Statement-level triggers require significant amount of state to be maintained in OLD TABLE and NEW TABLE
  - However, a row-level trigger does get fired for each row, so complex row-level triggers may be inefficient for statements that generate lots of modifications

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## Another statement-level trigger 11

- ❖ Give faculty a raise if GPA's in one update statement are all increasing

```
CREATE TRIGGER AutoRaise
AFTER UPDATE OF GPA ON Student
REFERENCING OLD TABLE AS o, NEW TABLE AS n
FOR EACH STATEMENT
```

```
UPDATE Faculty SET salary = salary + 1000;
```

- ☞ A row-level trigger would be difficult to write in this case

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## System issues 12

- ❖ Recursive firing of triggers
  - Action of one trigger causes another trigger to fire
  - Can get into an infinite loop
    - Some DBMS restrict trigger actions
    - Most DBMS set a maximum level of recursion (16 in DB2)
- ❖ Interaction with constraints (very tricky to get right!)
  - When do we check if a triggering event violates constraints?
    - After a BEFORE trigger (so the trigger can fix a potential violation)
    - Before an AFTER trigger
  - AFTER triggers also see the effects of, say, cascaded deletes caused by referential integrity constraint violations  
(Based on DB2; other DBMS may implement a different policy)

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## Views

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- ❖ A view is like a “virtual” table
  - Defined by a query, which describes how to compute the view contents on the fly
  - DBMS stores the view definition query instead of view contents
  - Can be used in queries just like a regular table

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## Creating and dropping views

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- ❖ Example: CPS116 roster
  - ```
CREATE VIEW CPS116Roster AS
  SELECT SID, name, age, GPA
  FROM Student
  WHERE SID IN (SELECT SID FROM Enroll
                WHERE CID = 'CPS116');
```

Called “base tables”
- ❖ To drop a view
 - ```
DROP VIEW view_name;
```

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## Using views in queries

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- ❖ Example: find the average GPA of CPS116 students
  - ```
SELECT AVG(GPA) FROM CPS116Roster;
```
 - To process the query, replace the reference to the view by its definition
 - ```
SELECT AVG(GPA)
 FROM (SELECT SID, name, age, GPA
 FROM Student
 WHERE SID IN (SELECT SID
 FROM Enroll
 WHERE CID = 'CPS116'));
```

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## Why use views?

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- ❖ To hide data from users
- ❖ To hide complexity from users
- ❖ Logical data independence
  - If applications deal with views, we can change the underlying schema without affecting applications
  - Recall physical data independence: change the physical organization of data without affecting applications
- ❖ To provide a uniform interface for different implementations or sources
- ☞ Real database applications use tons of views

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## Modifying views

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- ❖ Does not seem to make sense since views are virtual
- ❖ But does make sense if that is how users see the database
- ❖ Goal: modify the base tables such that the modification would appear to have been accomplished on the view

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## A simple case

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```
CREATE VIEW StudentGPA AS
 SELECT SID, GPA FROM Student;
DELETE FROM StudentGPA WHERE SID = 123;
translates to:
DELETE FROM Student WHERE SID = 123;
```

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## An impossible case

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```
CREATE VIEW HighGPASStudent AS
 SELECT SID, GPA FROM Student
 WHERE GPA > 3.7;
```

```
INSERT INTO HighGPASStudent
 VALUES(987, 2.5);
```

❖ No matter what you do on *Student*, the inserted row will not be in *HighGPASStudent*

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## A case with too many possibilities

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```
CREATE VIEW AverageGPA(GPA) AS
 SELECT AVG(GPA) FROM Student;
```

- Note that you can rename columns in view definition

```
UPDATE AverageGPA SET GPA = 2.5;
```

- ❖ Set everybody's GPA to 2.5?
- ❖ Adjust everybody's GPA by the same amount?
- ❖ Just lower Bart's GPA?

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## SQL92 updateable views

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- ❖ Single-table SFW

- No aggregation
- No subqueries

- ❖ Overly restrictive

- ❖ Still might get it wrong in some cases
  - See the slide titled "An impossible case"

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## Indexes

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- ❖ An index is an auxiliary persistent data structure
  - Search tree (e.g., B<sup>+</sup>-tree), lookup table (e.g., hash table), etc.
- ☞ More on indexes in the second half of this course!
- ❖ An index on  $R.A$  can speed up accesses of the form
  - $R.A = value$
  - $R.A > value$  (sometimes; depending on the index type)
- ❖ An index on  $(R.A_1, \dots, R.A_n)$  can speed up
  - $R.A_1 = value_1 \wedge \dots \wedge R.A_n = value_n$
  - $(R.A_1, \dots, R.A_n) > (value_1, \dots, value_n)$  (again depends)
- ☞ Is an index on  $(R.A, R.B)$  equivalent to one on  $(R.B, R.A)$ ?
- ☞ How about an index on  $R.A$  plus another index on  $R.B$ ?

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## Examples of using indexes

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- ❖ `SELECT * FROM Student WHERE name = 'Bart'`
  - Without an index on `Student.name`: must scan the entire table if we store `Student` as a flat file of unordered rows
  - With index: go “directly” to rows with `name = 'Bart'`
- ❖ `SELECT * FROM Student, Enroll WHERE Student.SID = Enroll.SID;`
  - Without any index: for each `Student` row, scan the entire `Enroll` table for matching `SID`
    - Sorting could help
  - With an index on `Enroll.SID`: for each `Student` row, directly look up `Enroll` rows with matching `SID`

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## Creating and dropping indexes in SQL

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- ❖ `CREATE [UNIQUE] INDEX index_name ON table_name(column_name1, ..., column_namen);`
  - With `UNIQUE`, the DBMS will also enforce that  $\{column\_name_1, \dots, column\_name_n\}$  is a key of `table_name`
- ❖ `DROP INDEX index_name;`
- ❖ Typically, the DBMS will automatically create indexes for `PRIMARY KEY` and `UNIQUE` constraint declarations

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## Choosing indexes to create

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More indexes = better performance?

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## Summary of SQL features covered so far

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- ❖ Query
- ❖ Modification
- ❖ Constraints
- ❖ Triggers
- ❖ Views
- ❖ Indexes

☞ Next: transactions

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