

Relational Database Design Part I

CPS 116
Introduction to Database Systems

Announcements (September 5)

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- ❖ rack040 accounts created; change your password!
 - Let me know if you have NOT received the email
- ❖ Homework #1 isn't quite ready yet
 - Will be handed out on Thursday
- ❖ Book value pack order fixed
 - Will probably arrive early next week
- ❖ Make use of office hours

Relational model: review

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- ❖ A database is a collection of relations (or tables)
- ❖ Each relation has a list of attributes (or columns)
- ❖ Each attribute has a domain (or type)
- ❖ Each relation contains a set of tuples (or rows)

Keys

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- ❖ A set of attributes K is a key for a relation R if
 - In no instance of R will two different tuples agree on all attributes of K
 - That is, K is a “tuple identifier”
 - No proper subset of K satisfies the above condition
 - That is, K is minimal
- ❖ Example: *Student* (SID , $name$, age , GPA)
 - SID is a key of *Student*
 - age is not a key (not an identifier)
 - $\{SID, name\}$ is not a key (not minimal)

Schema vs. data

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Student

<i>SID</i>	<i>name</i>	<i>age</i>	<i>GPA</i>
142	Bart	10	2.3
123	Milhouse	10	3.1
857	Lisa	8	4.3
456	Ralph	8	2.3

- ❖ Is $name$ a key of *Student*?
- ❖ Key declarations are part of the schema

More examples of keys

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- ❖ *Enroll* (SID , CID)
- ❖ *Address* ($street_address$, $city$, $state$, zip)
 - We typically pick one as the “primary” key, and underline all its attributes, e.g.,

Usage of keys

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- ❖ More constraints on data, fewer mistakes
- ❖ Look up a row by its key value
 - Many selection conditions are “key = value”
- ❖ “Pointers”
 - Example: *Enroll* (*SID*, *CID*)
 - *SID* is a key of *Student*
 - *CID* is a key of *Course*
 - An *Enroll* tuple “links” a *Student* tuple with a *Course* tuple
 - Many join conditions are “key = key value stored in another table”

Database design

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- ❖ Understand the real-world domain being modeled
- ❖ Specify it using a database design model
 - More intuitive and convenient for schema design
 - But not necessarily implemented by DBMS
 - A few popular ones:
 - Entity/Relationship (E/R) model
 - Object Definition Language (ODL)
 - UML (Unified Modeling Language)
- ❖ Translate specification to the data model of DBMS
 - Relational, XML, object-oriented, etc.
- ❖ Create DBMS schema

Entity-relationship (E/R) model

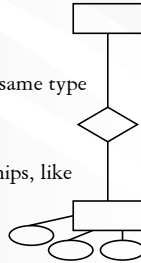
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- ❖ Historically and still very popular
- ❖ Can think of as a “watered-down” object-oriented design model
- ❖ Primarily a design model—not directly implemented by DBMS
- ❖ Designs represented by E/R diagrams
 - We use the style of E/R diagram covered by GMUW; there are other styles/extensions
 - Very similar to UML diagrams

E/R basics

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- ❖ Entity: a “thing,” like an object
- ❖ Entity set: a collection of things of the same type, like a relation of tuples or a class of objects
 - Represented as a rectangle
- ❖ Relationship: an association among entities
- ❖ Relationship set: a set of relationships of the same type (among the same entity sets)
 - Represented as a diamond
- ❖ Attributes: properties of entities or relationships, like attributes of tuples or objects
 - Represented as ovals



An example E/R diagram

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- ❖ Students enroll in courses

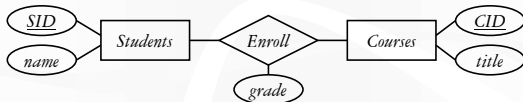


- ❖ A key of an entity set is represented by underlining all attributes in the key
 - A key is a set of attributes whose values can belong to at most one entity in an entity set—like a key of a relation

Attributes of relationships

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- ❖ Example: students take courses and receive grades


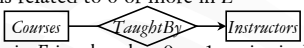



- ❖ Where do the grades go?

More on relationships

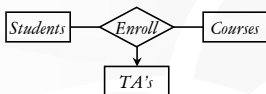
- ❖ There could be multiple relationship sets between the same entity sets
 - Example: *Students Enroll Courses*; *Students TA Courses*
- ❖ In a relationship set, each relationship is uniquely identified by the entities it connects
 - Example: Between Bart and CPS116, there can be at most one *Enroll* relationship and at most one *TA* relationship
 - ☞ What if Bart took CPS116 twice and got two different grades?

Multiplicity of relationships

- ❖ *E* and *F*: entity sets
- ❖ Many-many: Each entity in *E* is related to 0 or more entities in *F* and vice versa
 - Example: 
- ❖ Many-one: Each entity in *E* is related to 0 or 1 entity in *F*, but each entity in *F* is related to 0 or more in *E*
 - Example: 
- ❖ One-one: Each entity in *E* is related to 0 or 1 entity in *F* and vice versa
 - Example: 
- ❖ "One" (0 or 1) is represented by an arrow \longrightarrow
- ❖ "Exactly one" is represented by a rounded arrow $\longrightarrow\curvearrowright$

N-ary relationships

- ❖ Example: Each course has multiple TA's; each student is assigned to one TA

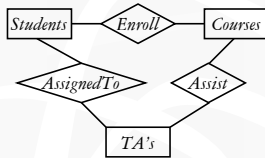


- ❖ Meaning of an arrow into *E*: Pick one entity from each of the other entity sets; together they must be related to either 0 or 1 entity in *E*

N-ary versus binary relationships

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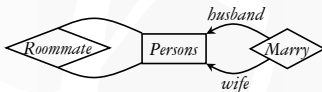
- ❖ Can we model n -ary relationships using just binary relationships?



Roles in relationships

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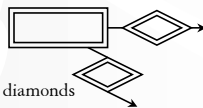
- ❖ An entity set may participate more than once in a relationship set
- ☞ May need to label edges to distinguish roles
- ❖ Examples
 - People are married as husband and wife; label needed
 - People are roommates of each other; label not needed



Weak entity sets

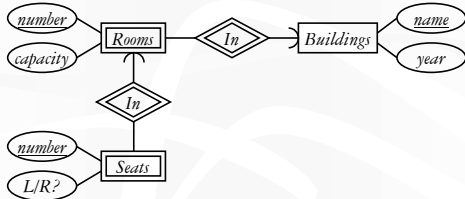
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- ❖ Sometimes, the key of an entity set E comes not completely from its own attributes, but from the keys of other (one or more) entity sets to which E is linked by many-one (or one-one) relationship sets
- Example: *Rooms* inside *Buildings* are partly identified by *Buildings*' name
- E is called a weak entity set
 - Denoted by double rectangle
 - The relationship sets through which E obtains its key are drawn as double diamonds



Weak entity set examples

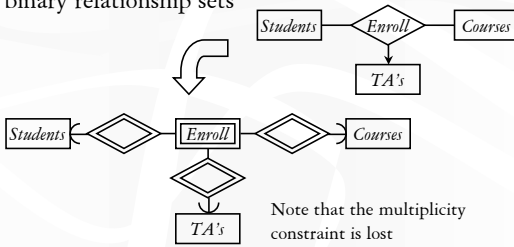
❖ Seats in rooms in buildings



❖ Why must double diamonds be many-one/one-one?

Modeling *n*-ary relationships

❖ An *n*-ary relationship set can be replaced by a weak entity set (called a connecting entity set) and *n* binary relationship sets

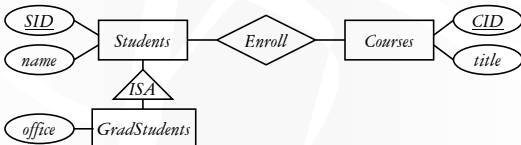


ISA relationships

❖ Similar to the idea of subclasses in object-oriented programming: subclass = special case, fewer entities, and possibly more properties

- Represented as a triangle (direction is important)

❖ Example: Graduate students are students, but they also have offices



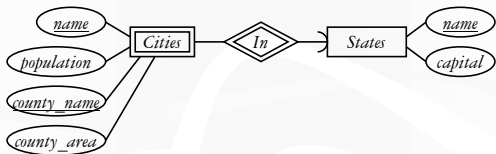
Summary of E/R concepts

- ❖ Entity sets
 - Keys
 - Weak entity sets
- ❖ Relationship sets
 - Attributes of relationships
 - Multiplicity
 - Roles
 - Binary versus *N*-ary relationships
 - Modeling *N*-ary relationships with weak entity sets and binary relationships
 - ISA relationships

Case study 1

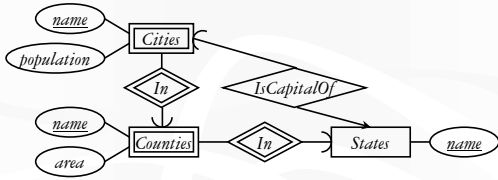
- ❖ Design a database representing cities, counties, and states
 - For states, record name and capital (city)
 - For counties, record name, area, and location (state)
 - For cities, record name, population, and location (county and state)
- ❖ Assume the following:
 - Names of states are unique
 - Names of counties are only unique within a state
 - Names of cities are only unique within a county
 - A city is always located in a single county
 - A county is always located in a single state

Case study 1: first design



Case study 1: second design

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❖ Technically, nothing in this design could prevent a city in state X from being the capital of another state Y, but oh well...

Case study 2

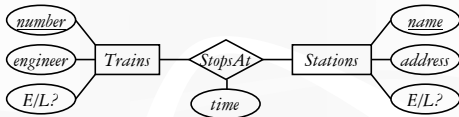
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❖ Design a database consistent with the following:

- A station has a unique name and an address, and is either an express station or a local station
- A train has a unique number and an engineer, and is either an express train or a local train
- A local train can stop at any station
- An express train only stops at express stations
- A train can stop at a station for any number of times during a day
- Train schedules are the same everyday

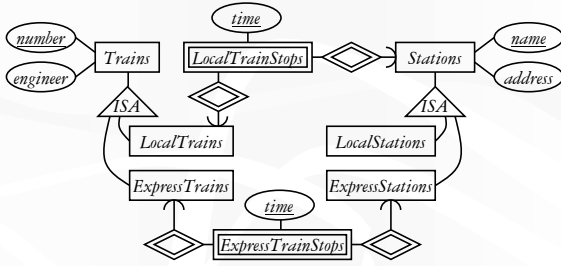
Case study 2: first design

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Case study 2: second design

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Is the extra complexity worth it?
