Relational Database Design
Part I

CPS 116
Introduction to Database Systems

Announcements (September 5)

- rack040 accounts created; change your password!
  - Let me know if you have NOT received the email
- Homework #1 isn't quite ready yet
  - Will be handed out on Thursday
- Book value pack order fixed
  - Will probably arrive early next week
- Make use of office hours

Relational model: review

- A database is a collection of relations (or tables)
- Each relation has a list of attributes (or columns)
- Each attribute has a domain (or type)
- Each relation contains a set of tuples (or rows)

Keys

- A set of attributes *K* is a key for a relation *R* if
  - In no instance of *R* will two different tuples agree on all attributes of *K*
    - That is, *K* is a "tuple identifier"
  - No proper subset of *K* satisfies the above condition
    - That is, *K* is minimal
- Example: *Student* (*SID*, *name*, *age*, *GPA*)
  - *SID* is a key of *Student*
  - *age* is not a key (not an identifier)
  - {*SID*, *name*} is not a key (not minimal)

Schema vs. data

Student

SID	name	age	GPA
142	Bart	10	2.3
123	Milhouse	10	3.1
857	Lisa	8	4.3
456	Ralph	8	2.3

- Is *name* a key of *Student*?
  - Yes? Seems reasonable for this instance
  - No! Student names are not unique in general
- Key declarations are part of the schema

More examples of keys

- *Enroll* (*SID*, *CID*)
  - {*SID*, *CID*}
  - ☞ A key can contain multiple attributes!
- *Address* (*street_address*, *city*, *state*, *zip*)
  - {*street_address*, *city*, *state*}
  - {*street_address*, *zip*}
  - A relation can have multiple keys!
    - We typically pick one as the "primary" key, and underline all its attributes, e.g., *Address* (<u>*street_address*</u>, *city*, *state*, <u>*zip*</u>)

Usage of keys

- More constraints on data, fewer mistakes
- Look up a row by its key value
  - Many selection conditions are "key = value"
- "Pointers"
  - Example: *Enroll* (*SID*, *CID*)
    - *SID* is a key of *Student*
    - *CID* is a key of *Course*
    - An *Enroll* tuple "links" a *Student* tuple with a *Course* tuple
  - Many join conditions are "key = key value stored in another table"

Database design

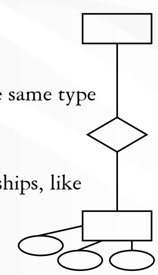
- Understand the real-world domain being modeled
- Specify it using a database design model
  - More intuitive and convenient for schema design
  - But not necessarily implemented by DBMS
  - A few popular ones:
    - Entity/Relationship (E/R) model
    - Object Definition Language (ODL)
    - UML (Unified Modeling Language)
- Translate specification to the data model of DBMS
  - Relational, XML, object-oriented, etc.
- Create DBMS schema

Entity-relationship (E/R) model

- Historically and still very popular
- Can think of as a "watered-down" object-oriented design model
- Primarily a design model—not directly implemented by DBMS
- Designs represented by E/R diagrams
  - We use the style of E/R diagram covered by GMUW; there are other styles/extensions
  - Very similar to UML diagrams

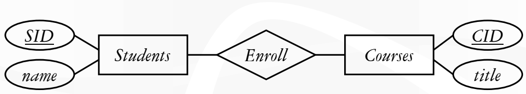
E/R basics

- Entity: a "thing," like an object
- Entity set: a collection of things of the same type, like a relation of tuples or a class of objects
  - Represented as a rectangle
- Relationship: an association among entities
- Relationship set: a set of relationships of the same type (among same entity sets)
  - Represented as a diamond
- Attributes: properties of entities or relationships, like attributes of tuples or objects
  - Represented as ovals



An example E/R diagram

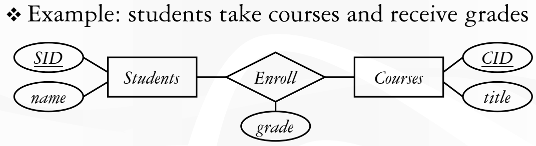
- Students enroll in courses



- A key of an entity set is represented by underlining all attributes in the key
  - A key is a set of attributes whose values can belong to at most one entity in an entity set—like a key of a relation

Attributes of relationships

- Example: students take courses and receive grades



- Where do the grades go?
  - With *Students*?
    - But a student can have different grades for multiple courses
  - With *Courses*?
    - But a course can assign different grades for multiple students
  - With *Enroll*!

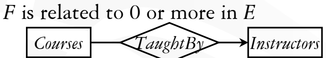
More on relationships

- There could be multiple relationship sets between the same entity sets
  - Example: *Students Enroll Courses*; *Students TA Courses*
- In a relationship set, each relationship is uniquely identified by the entities it connects
  - Example: Between Bart and CPS116, there can be at most one *Enroll* relationship and at most one *TA* relationship
  - ☞What if Bart took CPS116 twice and got two different grades?

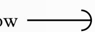
Multiplicity of relationships

- *E* and *F*: entity sets
- Many-many: Each entity in *E* is related to 0 or more entities in *F* and vice versa
  - Example:

- Many-one: Each entity in *E* is related to 0 or 1 entity in *F*, but each entity in *F* is related to 0 or more in *E*
  - Example:

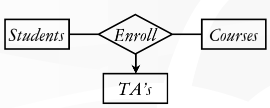


- One-one: Each entity in *E* is related to 0 or 1 entity in *F* and vice versa
  - Example:

- "One" (0 or 1) is represented by an arrow ⟶
- "Exactly one" is represented by a rounded arrow 

N-ary relationships

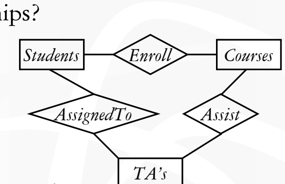
- Example: Each course has multiple TA's; each student is assigned to one TA



- Meaning of an arrow into *E*: Pick one entity from each of the other entity sets; together they must be related to either 0 or 1 entity in *E*

N-ary versus binary relationships

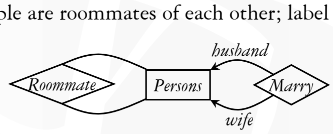
- Can we model *n*-ary relationships using just binary relationships?



- No; for example:
  - Bart takes CPS116 and CPS114
  - Lisa TA's CPS116 and CPS114
  - Bart is assigned to Lisa in CPS116, but not in CPS114

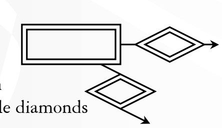
Roles in relationships

- An entity set may participate more than once in a relationship set
- ☞ May need to label edges to distinguish roles
- Examples
  - People are married as husband and wife; label needed
  - People are roommates of each other; label not needed



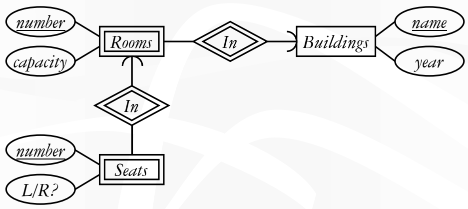
Weak entity sets

- Sometimes, the key of an entity set *E* comes not completely from its own attributes, but from the keys of other (one or more) entity sets to which *E* is linked by many-one (or one-one) relationship sets
  - Example: *Rooms* inside *Buildings* are partly identified by *Buildings*' name
  - *E* is called a weak entity set
    - Denoted by double rectangle
    - The relationship sets through which *E* obtains its key are drawn as double diamonds



Weak entity set examples

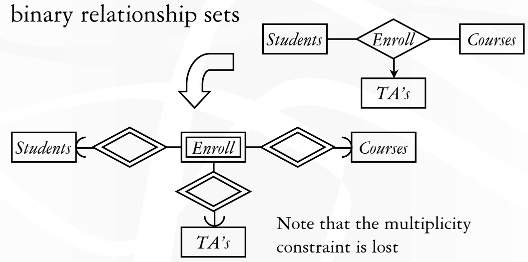
- Seats in rooms in buildings



- Why must double diamonds be many-one/one-one?
  - With many-many, we would not know which entity provides the key value!

Modeling *n*-ary relationships

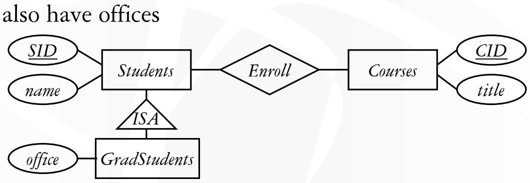
- An *n*-ary relationship set can be replaced by a weak entity set (called a connecting entity set) and *n* binary relationship sets



Note that the multiplicity constraint is lost

ISA relationships

- Similar to the idea of subclasses in object-oriented programming: subclass = special case, fewer entities, and possibly more properties
  - Represented as a triangle (direction is important)
- Example: Graduate students are students, but they also have offices



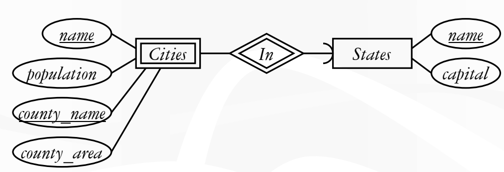
Summary of E/R concepts

- Entity sets
  - Keys
  - Weak entity sets
- Relationship sets
  - Attributes of relationships
  - Multiplicity
  - Roles
  - Binary versus *N*-ary relationships
    - Modeling *N*-ary relationships with weak entity sets and binary relationships
  - ISA relationships

Case study 1

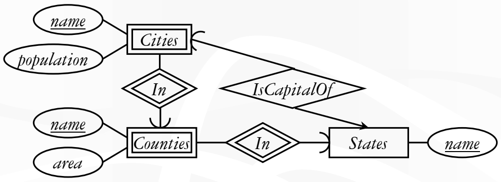
- Design a database representing cities, counties, and states
  - For states, record name and capital (city)
  - For counties, record name, area, and location (state)
  - For cities, record name, population, and location (county and state)
- Assume the following:
  - Names of states are unique
  - Names of counties are only unique within a state
  - Names of cities are only unique within a county
  - A city is always located in a single county
  - A county is always located in a single state

Case study 1: first design



- County area information is repeated for every city in the county
  - ☞Redundancy is bad (why?)
- State capital should really be a city
  - ☞Should "reference" entities through explicit relationships

Case study 1: second design



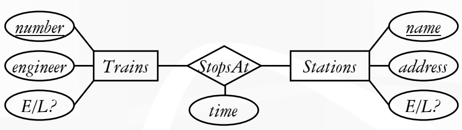
❖ Technically, nothing in this design could prevent a city in state *X* from being the capital of another state *Y*, but oh well…

Case study 2

❖ Design a database consistent with the following:

- A station has a unique name and an address, and is either an express station or a local station
- A train has a unique number and an engineer, and is either an express train or a local train
- A local train can stop at any station
- An express train only stops at express stations
- A train can stop at a station for any number of times during a day
- Train schedules are the same everyday

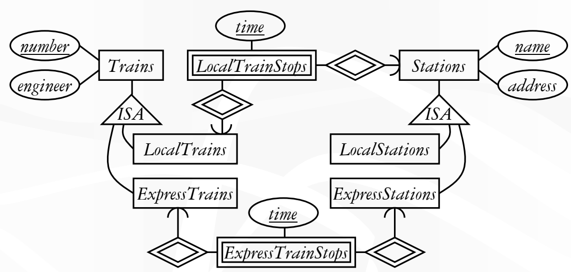
Case study 2: first design



❖ Nothing in this design prevents express trains from stopping at local stations
  - ☞Should capture as many constraints as possible

❖ A train can stop at a station only once during a day
  - ☞Should not introduce constraints

Case study 2: second design



Is the extra complexity worth it?