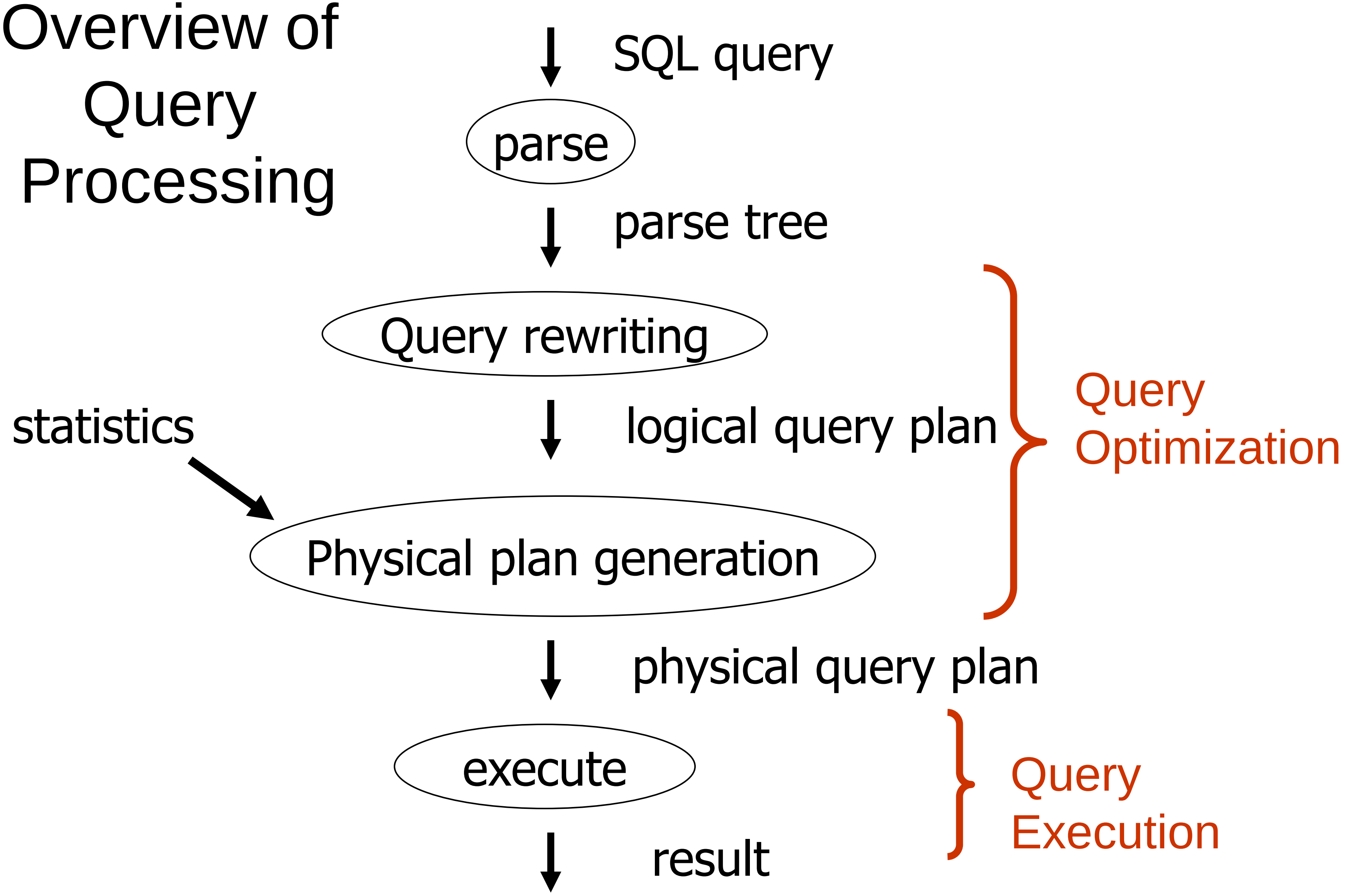
CPS216: Data-Intensive Computing Systems

Query Processing (contd.)

Shivnath Babu

Overview of Query Processing



SQL query

parse

parse tree

Query rewriting

logical query plan

statistics

Physical plan generation

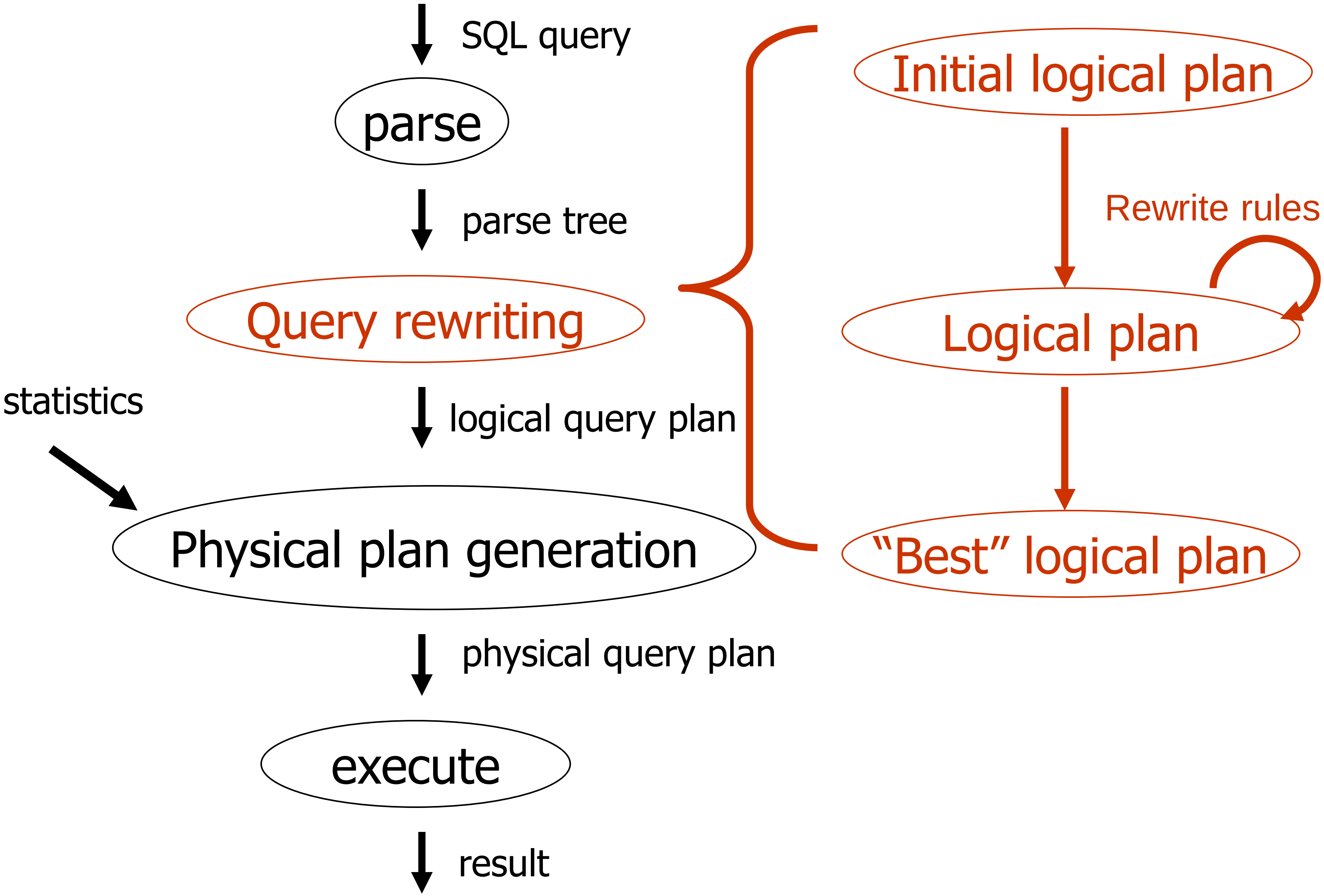
physical query plan

execute

result

Query Optimization

Query Execution



SQL query

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physical query plan

execute

result

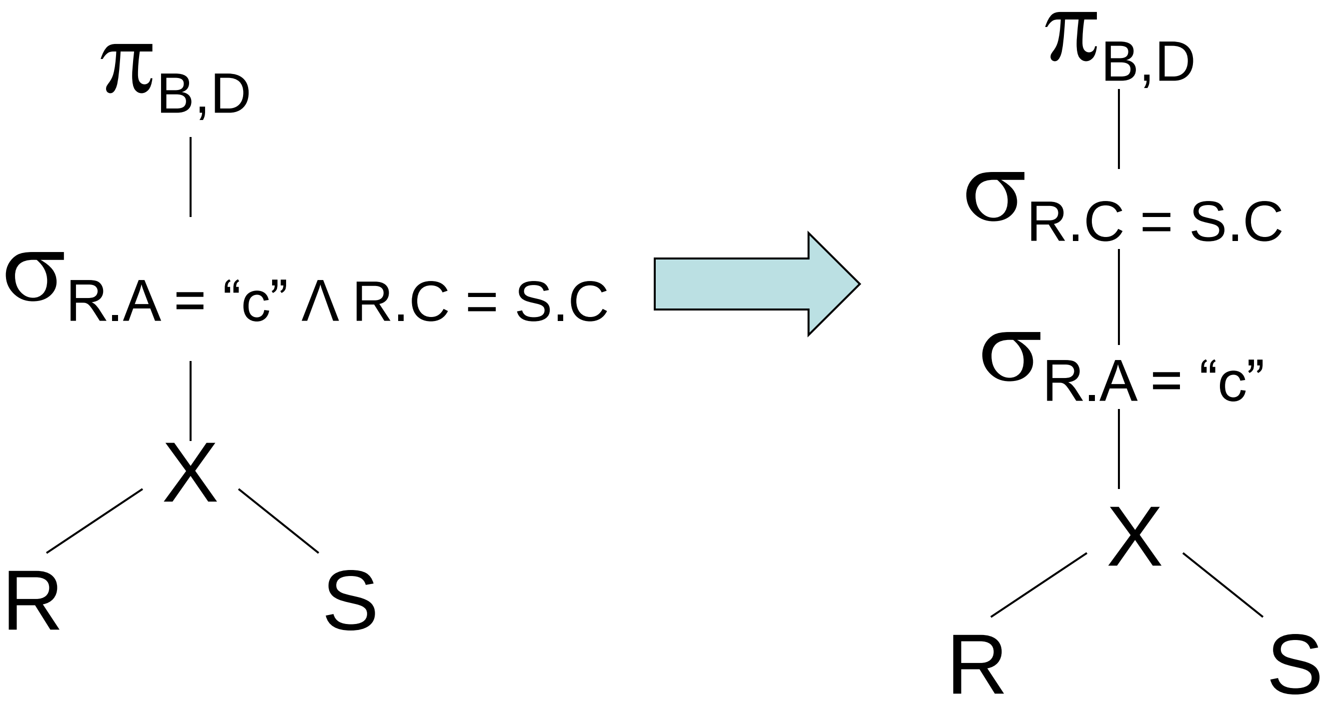
Initial logical plan

Rewrite rules

Logical plan

"Best" logical plan

Query Rewriting



$\pi_{B,D}$

$\sigma_{R.A = \text{“c”} \land R.C = S.C}$

X

R S

$\pi_{B,D}$

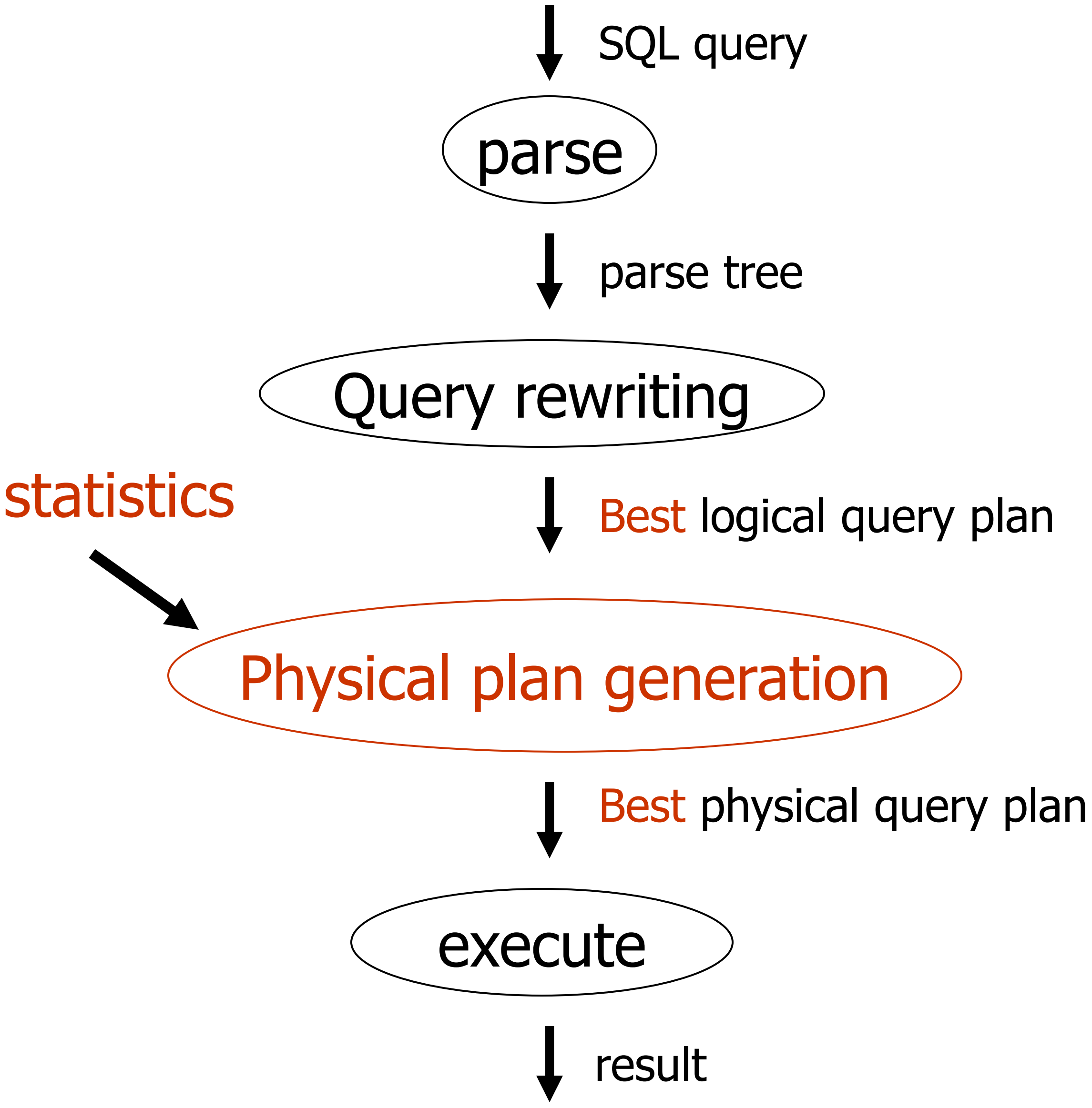
$\sigma_{R.C = S.C}$

$\sigma_{R.A = \text{“c”}}$

X

R S

We will revisit it towards the end of this lecture



SQL query

parse

parse tree

Query rewriting

statistics

Best logical query plan

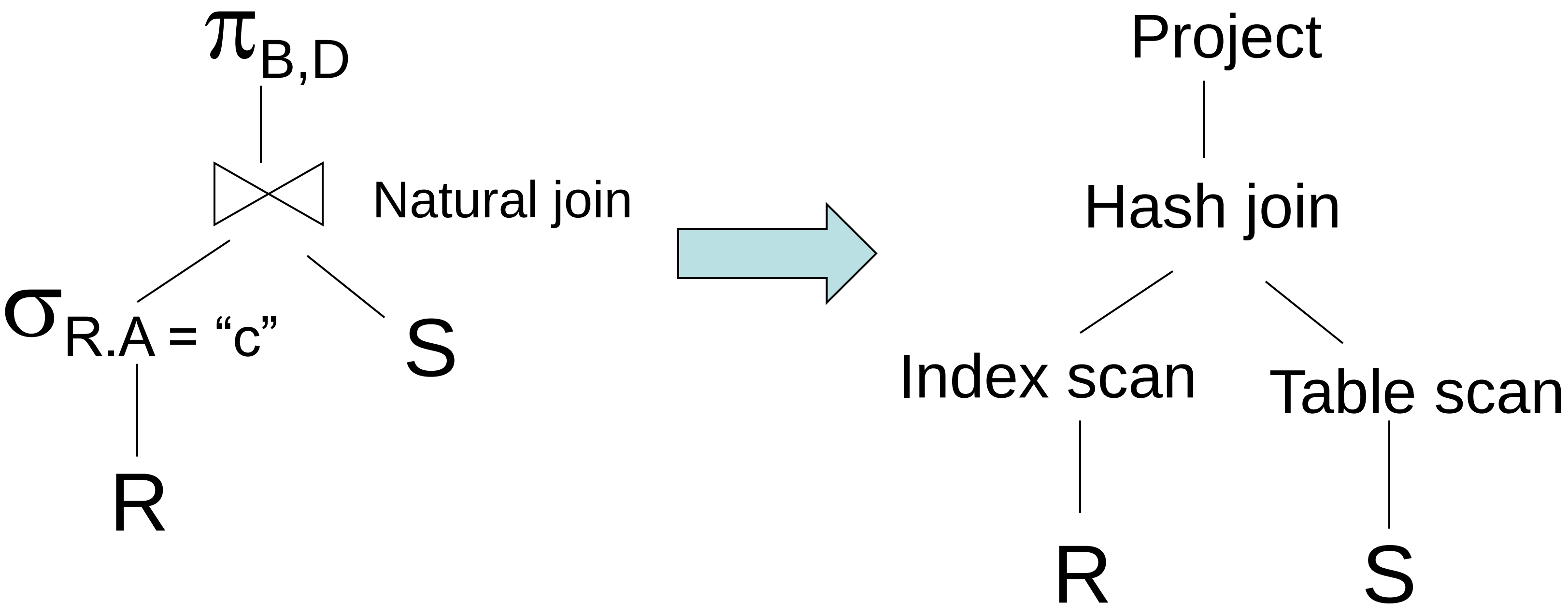
Physical plan generation

Best physical query plan

execute

result

Physical Plan Generation



$\pi_{B,D}$

$\bowtie$ Natural join

$\sigma_{R.A = \text{“c”}}$

S

R

Best logical plan

Project

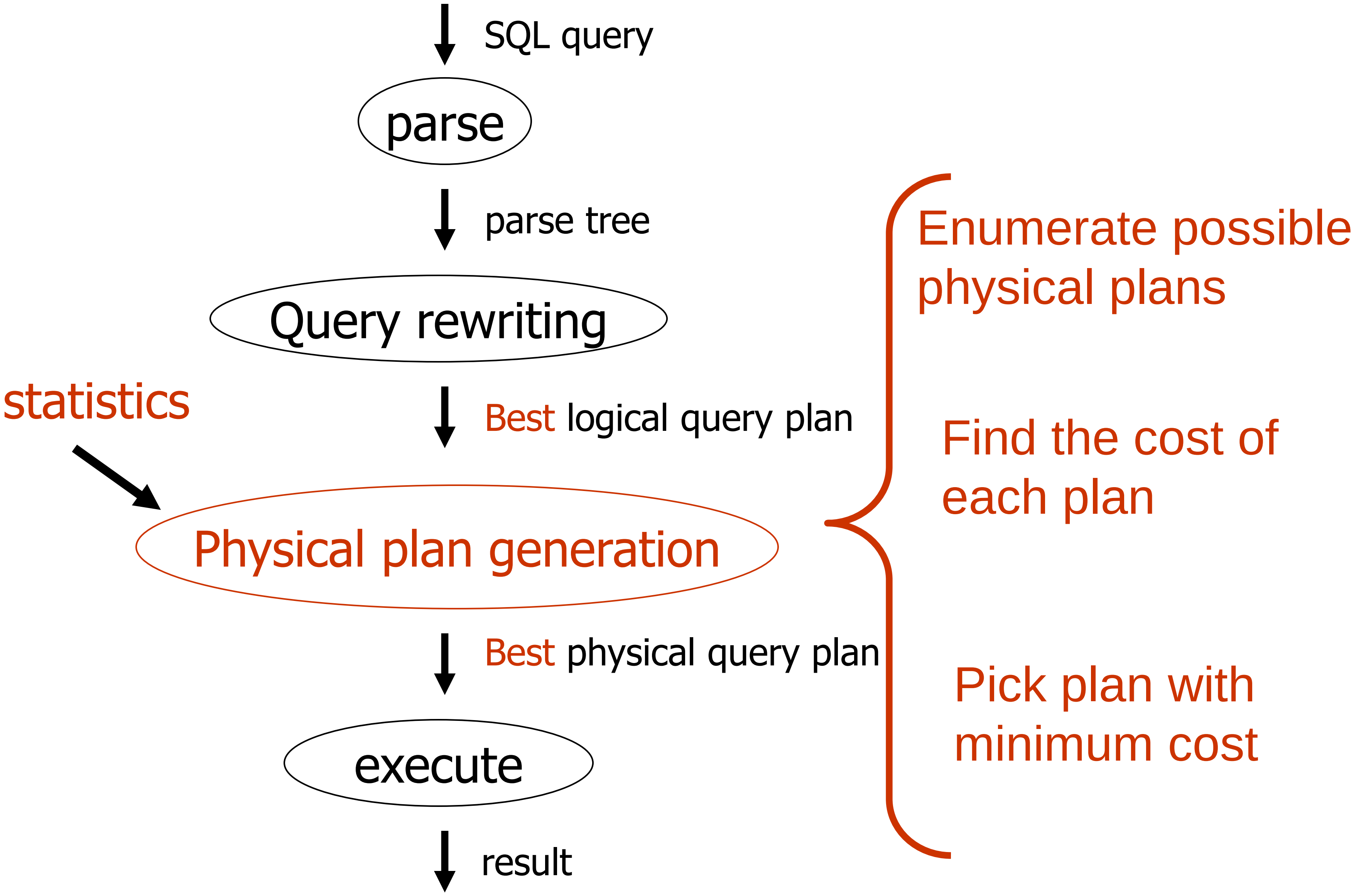
Hash join

Index scan

Table scan

R

S



SQL query

parse

parse tree

Query rewriting

Best logical query plan

statistics

Physical plan generation

Best physical query plan

execute

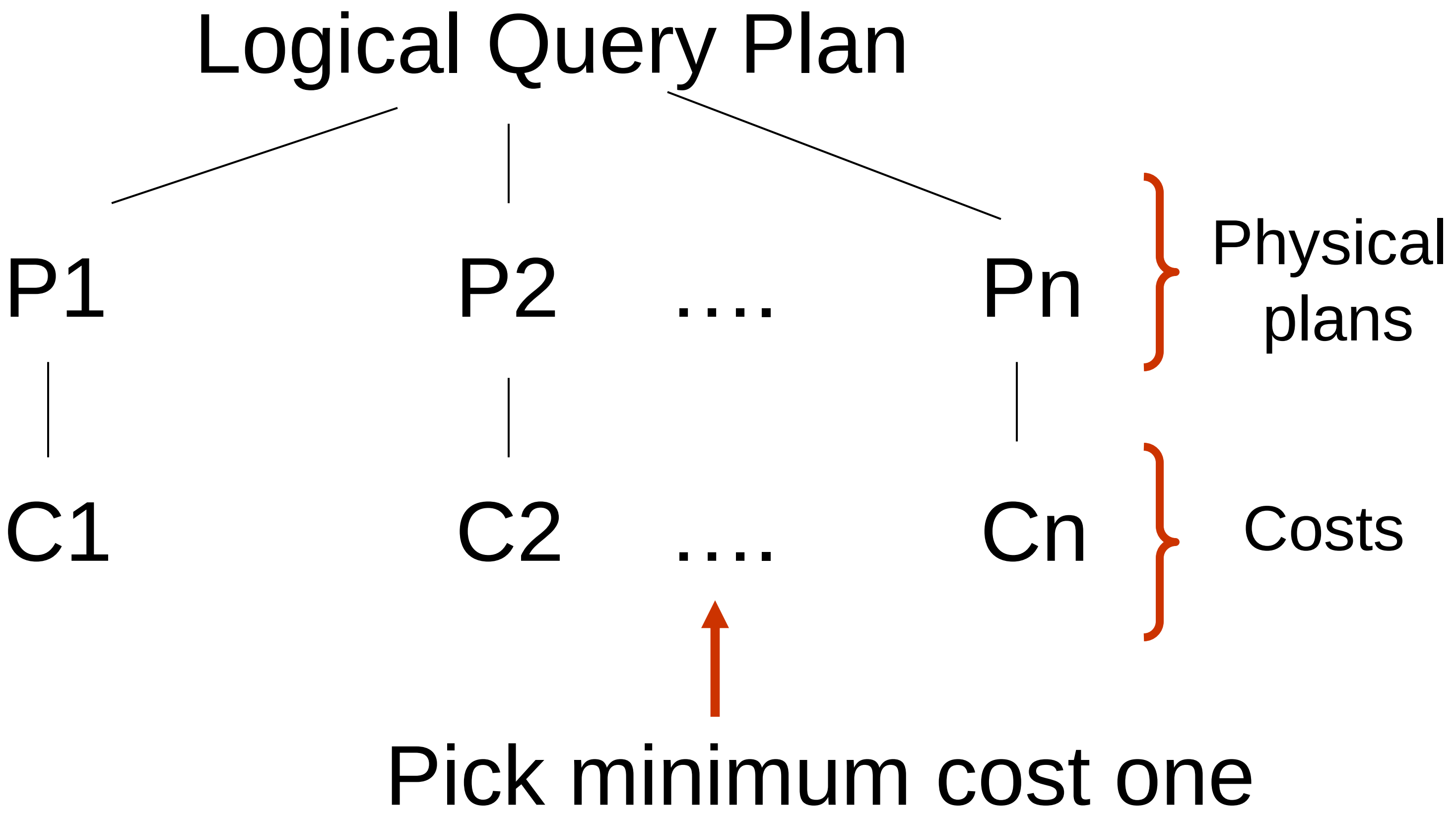
result

Enumerate possible physical plans

Find the cost of each plan

Pick plan with minimum cost

Physical Plan Generation



Logical Query Plan

P1 P2 …. Pn

Physical plans

C1 C2 …. Cn

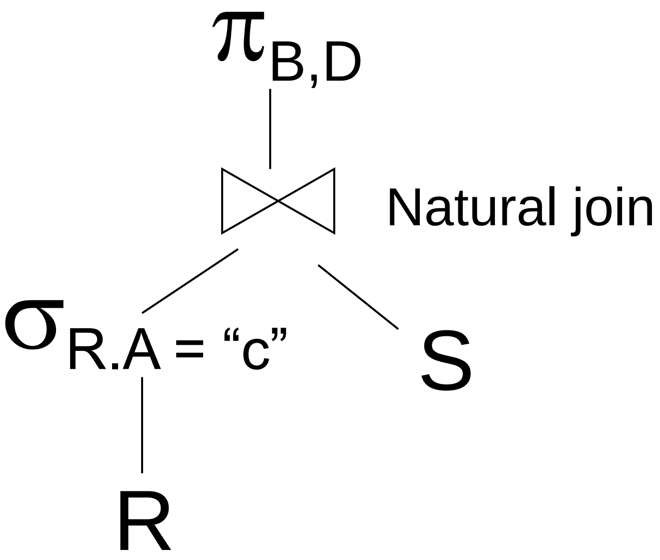
Costs

Pick minimum cost one

Plans for Query Execution

- Roadmap
  - Path of a SQL query
  - Operator trees
  - Physical Vs Logical plans
  - Plumbing: Materialization Vs pipelining

Logical Plans Vs. Physical Plans



$\pi_{B,D}$

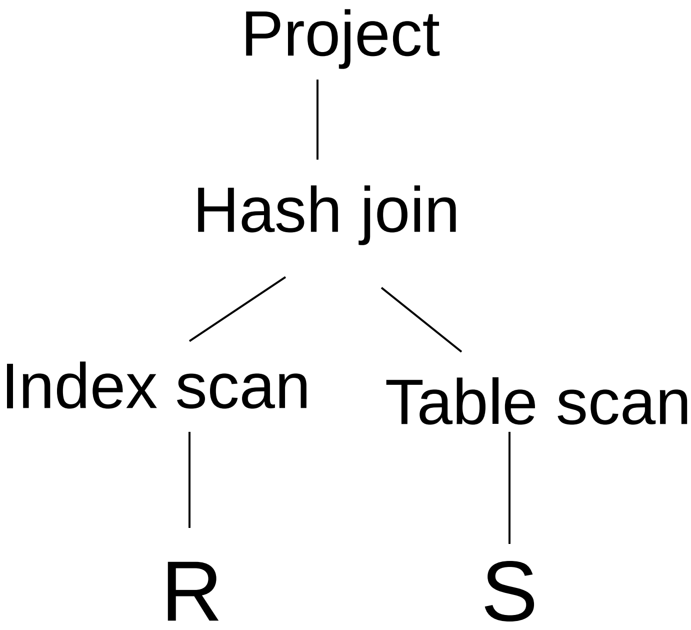
$\bowtie$ Natural join

$\sigma_{R.A = \text{“c”}}$

S

R

Best logical plan



Project

Hash join

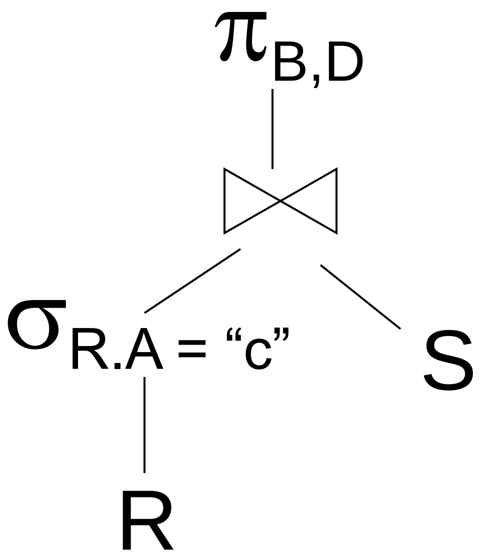
Index scan

Table scan

R

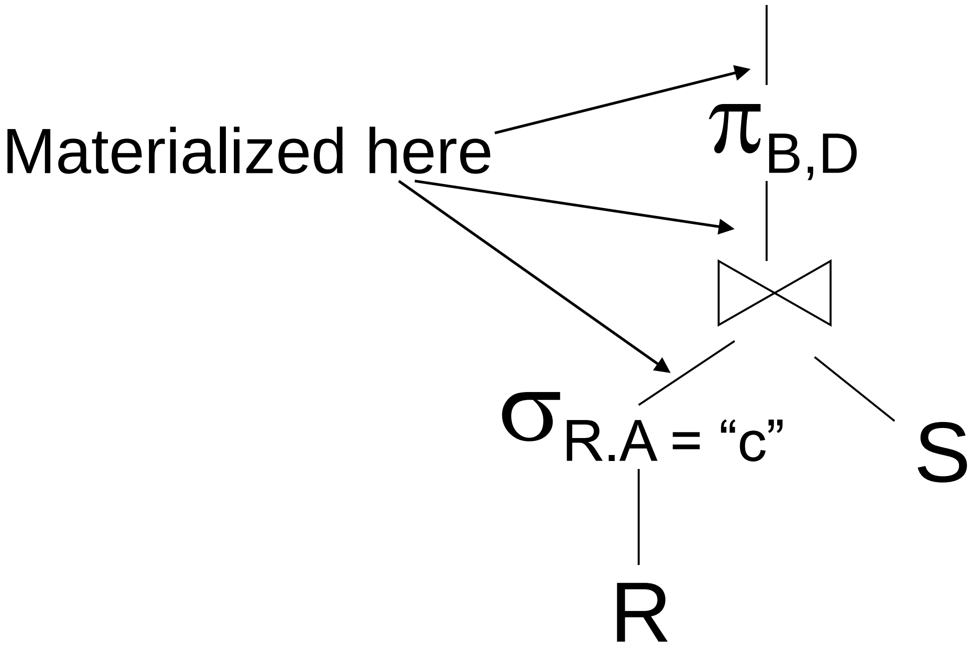
S

Operator Plumbing



- Materialization: output of one operator written to disk, next operator reads from the disk
- Pipelining: output of one operator directly fed to next operator

Materialization



Materialized here

$\pi_{B,D}$

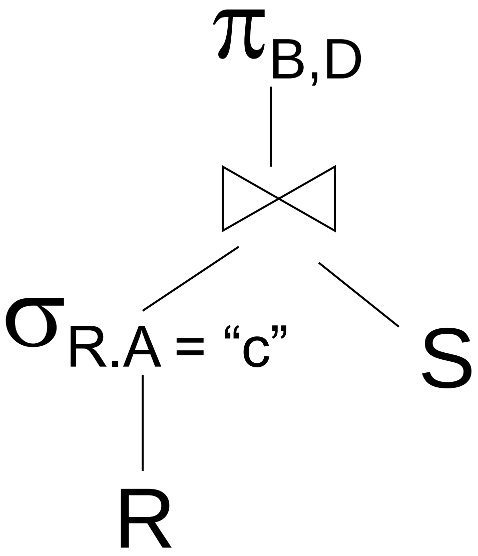
$\bowtie$

$\sigma_{R.A = \text{“c”}}$

S

R

Iterators: Pipelining



➔ Each operator supports:

- Open()
- GetNext()
- Close()

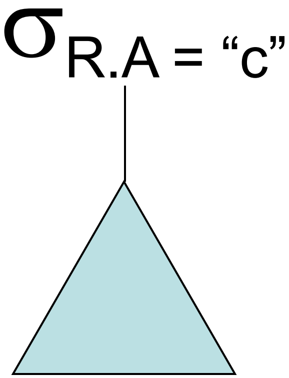
Iterator for Table Scan (R)

```
Open() {
  /** initialize variables */
  b = first block of R;
  t = first tuple in block b;
}
```

```
GetNext() {
  IF (t is past last tuple in block b) {
      set b to next block;
      IF (there is no next block)
          /** no more tuples */
          RETURN EOT;
      ELSE t = first tuple in b;
  }
  /** return current tuple */
  oldt = t;
  set t to next tuple in block b;
  RETURN oldt;
}
```

```
Close() {
  /** nothing to be done */
}
```

Iterator for Select

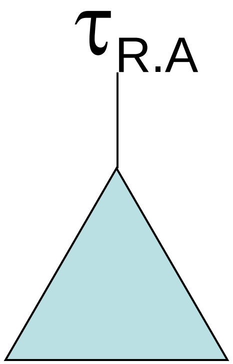


```
Open() {
  /** initialize child */
  Child.Open();
}

Close() {
  /** inform child */
  Child.Close();
}
```

```
GetNext() {
  LOOP:
      t = Child.GetNext();
      IF (t == EOT) {
        /** no more tuples */
        RETURN EOT;
      }
      ELSE IF (t.A == “c”)
        RETURN t;
  ENDLOOP:
}
```

Iterator for Sort

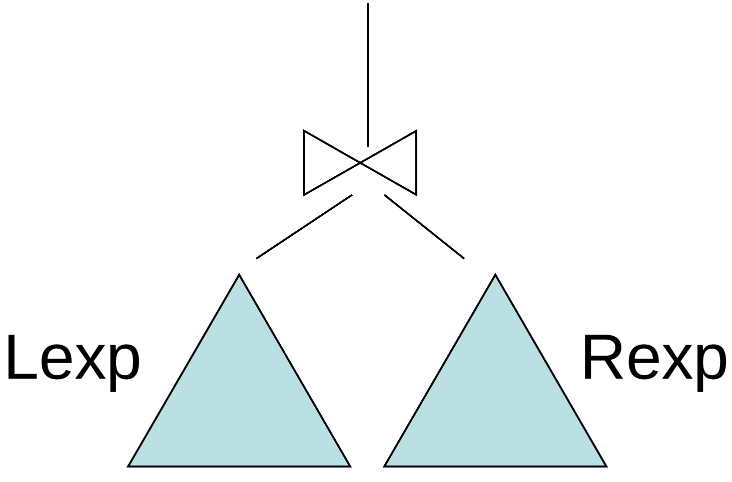


```
GetNext() {
   IF (more tuples)
      RETURN next tuple in order;
   ELSE RETURN EOT;
}
```

```
Open() {
  /** Bulk of the work is here */
  Child.Open();
  Read all tuples from Child
     and sort them
}
```

```
Close() {
  /** inform child */
  Child.Close();
}
```

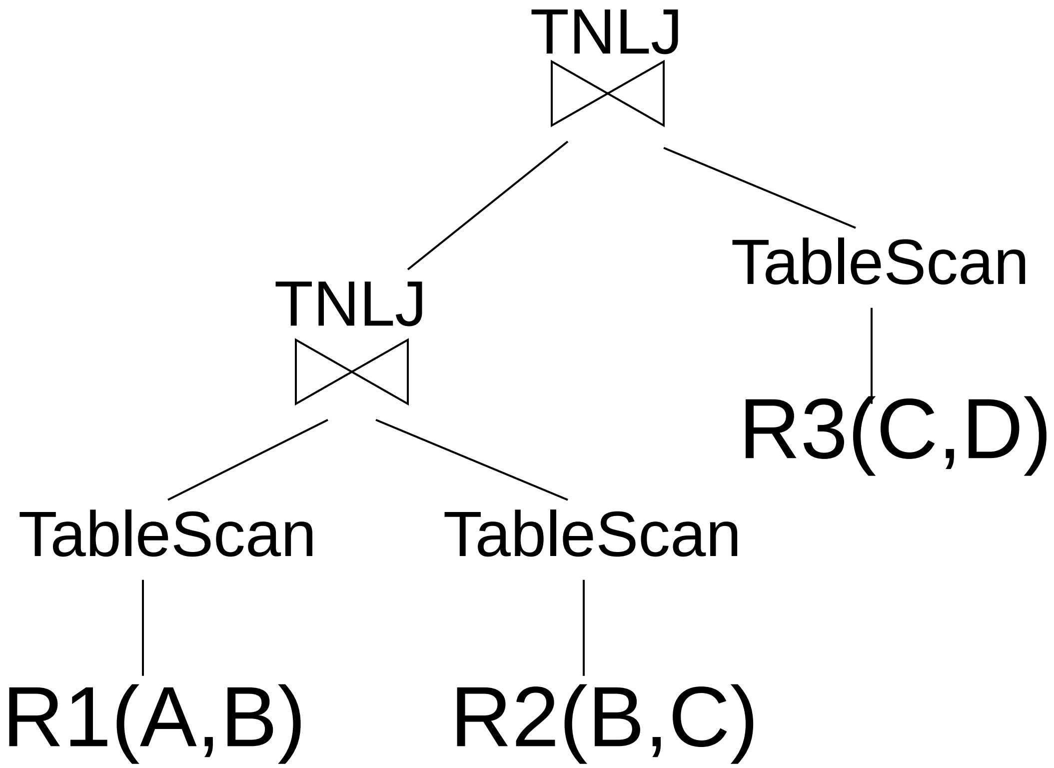
Iterator for Tuple Nested Loop Join



- TNLJ (conceptually)

```
for each r $\in$ Lexp do
    for each s $\in$ Rexp do
        if Lexp.C = Rexp.C, output r,s
```

Example 1: Left-Deep Plan



TNLJ ⋈

TNLJ ⋈

TableScan

R3(C,D)

TableScan

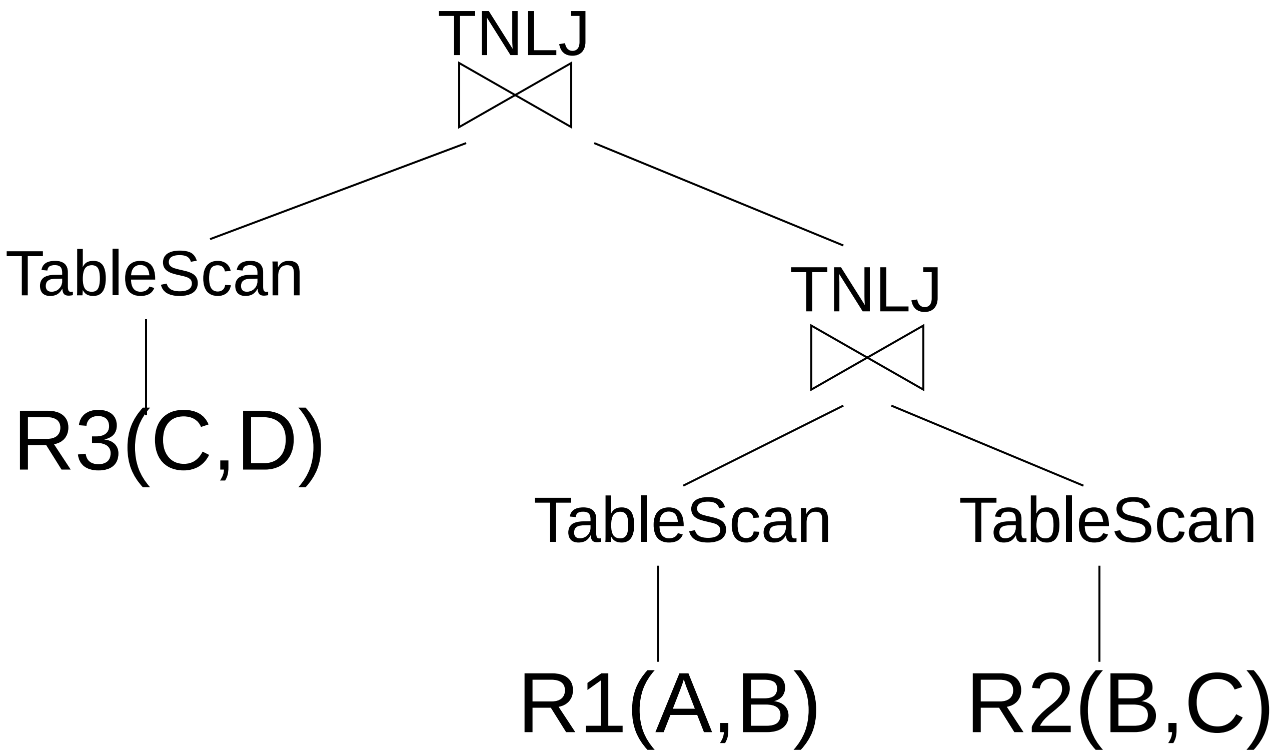
TableScan

R1(A,B)

R2(B,C)

Question: What is the sequence of getNext() calls?

Example 2: Right-Deep Plan



TNLJ

TableScan

R3(C,D)

TNLJ

TableScan

TableScan

R1(A,B)

R2(B,C)

Question: What is the sequence of getNext() calls?

Cost Measure for a Physical Plan

- There are many cost measures
  - Time to completion
  - Number of I/Os (we will see a lot of this)
  - Number of getNext() calls
- Tradeoff: Simplicity of estimation Vs. Accurate estimation of performance as seen by user

Why do we need Query Rewriting?

- Pruning the HUGE space of physical plans
  - Eliminating redundant conditions/operators
  - Rules that will improve performance with very high probability
- Preprocessing
  - Getting queries into a form that we know how to handle best

➔ Reduces optimization time drastically without noticeably affecting quality

Some Query Rewrite Rules

- Transform one logical plan into another
  - Do not use statistics
- Equivalences in relational algebra
- Push-down predicates
- Do projects early
- Avoid cross-products if possible

Equivalences in Relational Algebra

$R \bowtie S = S \bowtie R$ Commutativity

$(R \bowtie S) \bowtie T = R \bowtie (S \bowtie T)$ Associativity

Also holds for: Cross Products, Union, Intersection

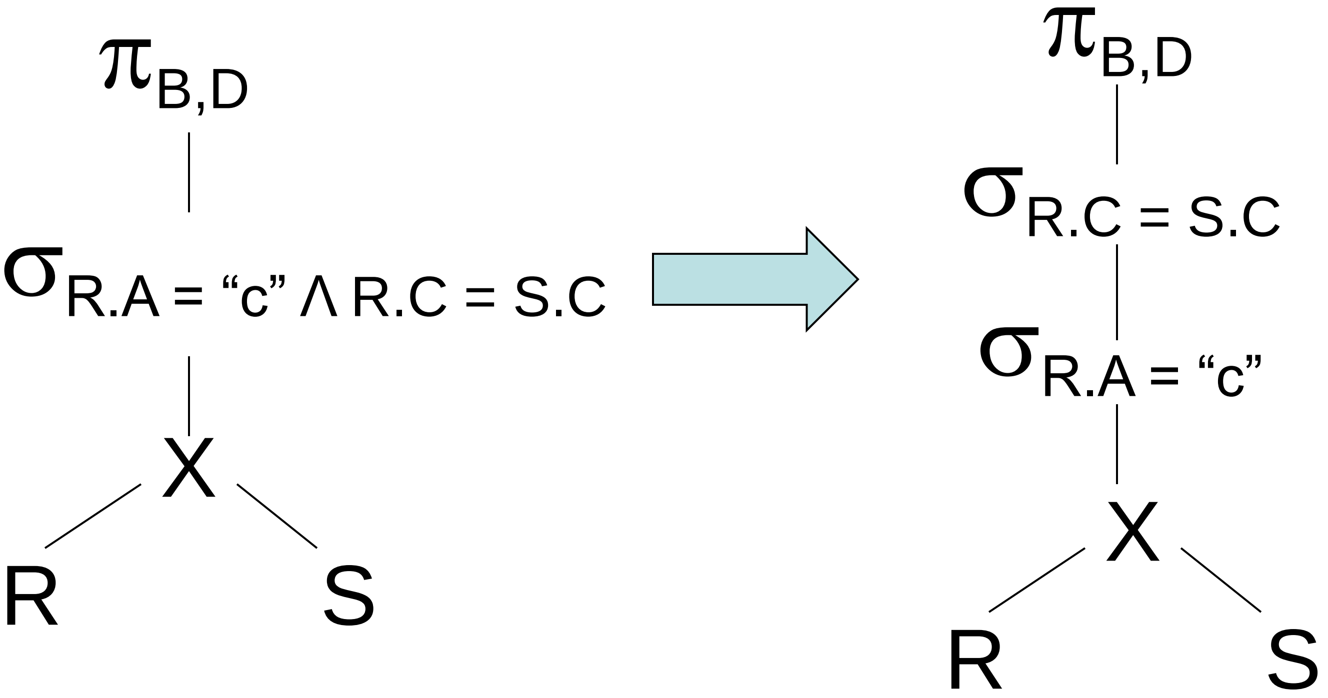
R x S = S x R

(R x S) x T = R x (S x T)

R U S = S U R

R U (S U T) = (R U S) U T

Apply Rewrite Rule (1)



$\pi_{B,D}$

$\sigma_{R.A = “c” \wedge R.C = S.C}$

X

R

S

$\pi_{B,D}$

$\sigma_{R.C = S.C}$

$\sigma_{R.A = “c”}$

X

R

S

$$\Pi_{B,D}[\sigma_{R.C=S.C}[\sigma_{R.A=“c”}(R \text{ X } S)]]$$

Rules: Project

Let: X = set of attributes

Y = set of attributes

XY = X U Y

$$\pi_{xy}(R) = \xcancel{\pi_x[\pi_y(R)]}$$

Rules: $\sigma + \bowtie$ combined

Let p = predicate with only R attribs

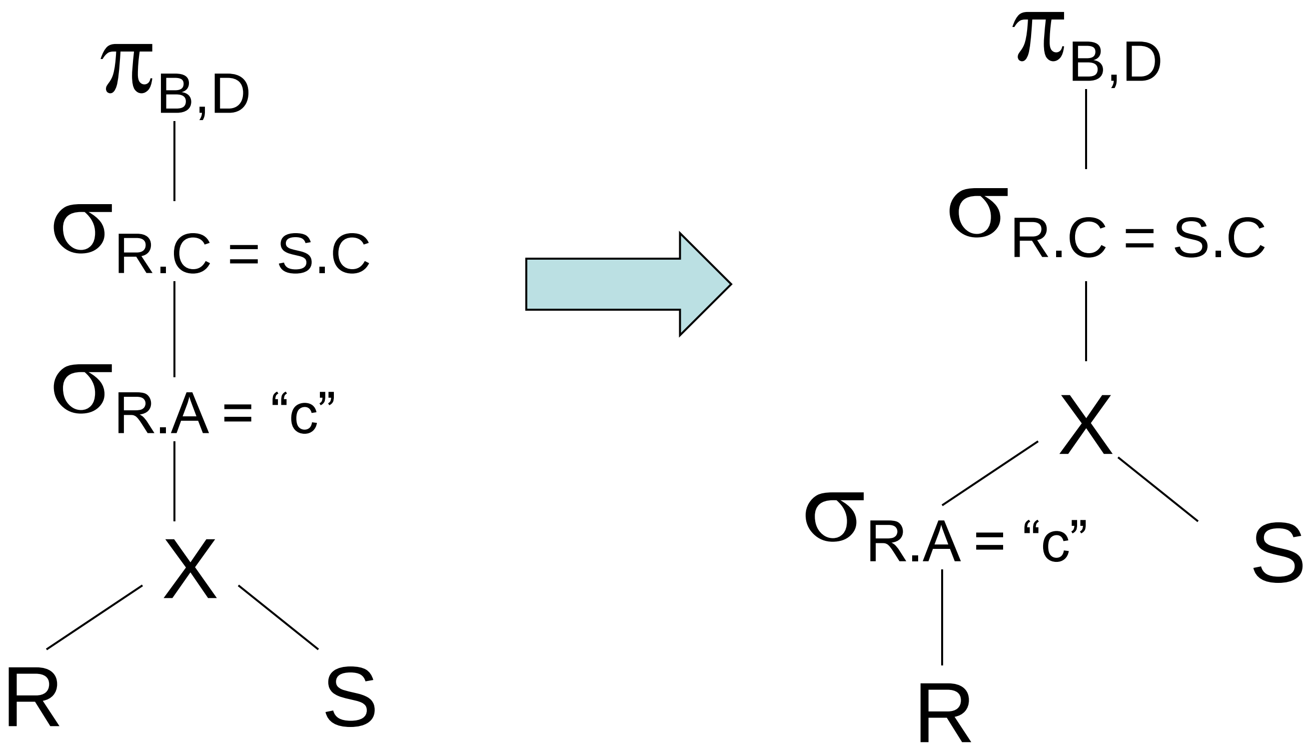
q = predicate with only S attribs

m = predicate with only R,S attribs

$$\sigma_p (R \bowtie S) = [\sigma_p (R)] \bowtie S$$

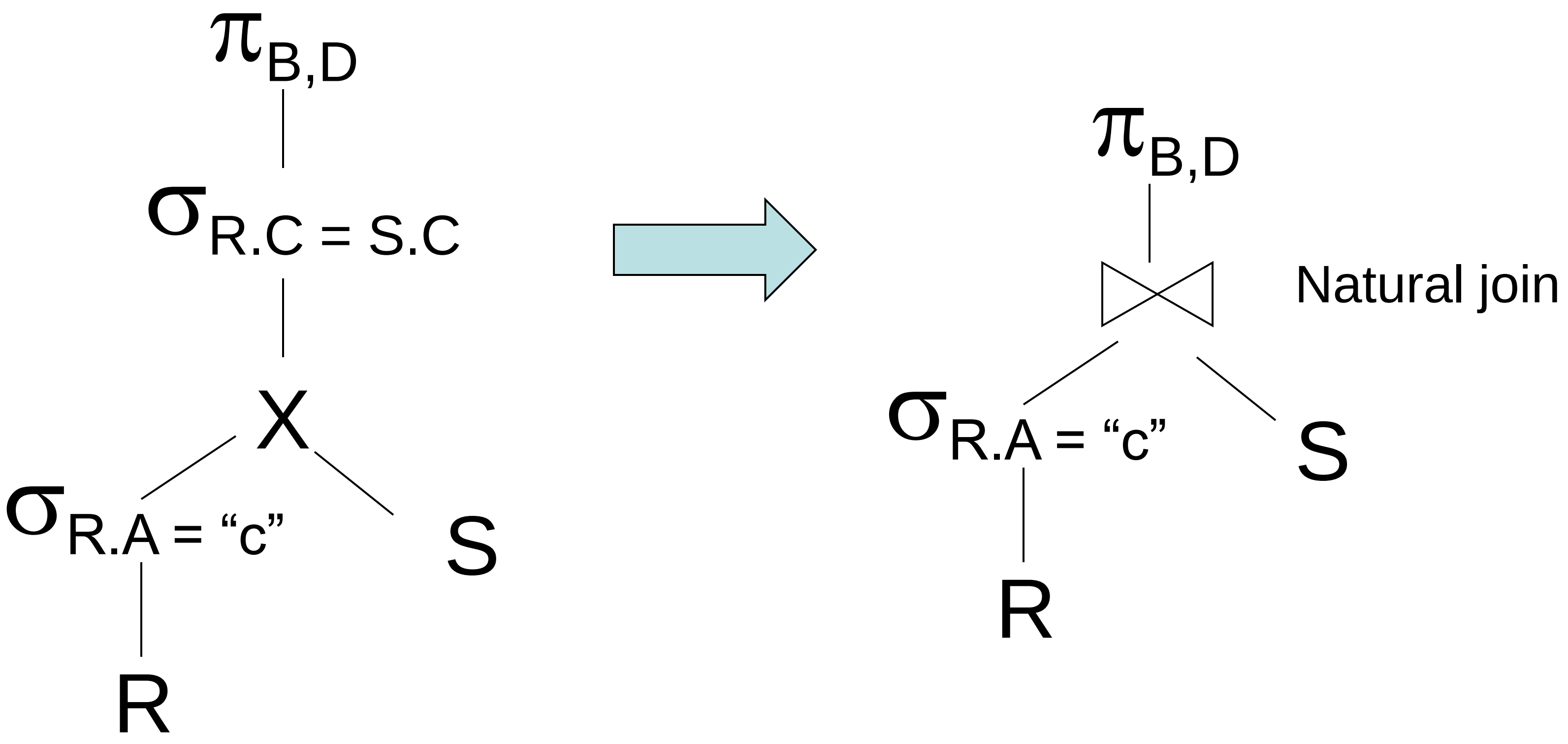
$$\sigma_q (R \bowtie S) = R \bowtie [\sigma_q (S)]$$

Apply Rewrite Rule (2)



$$\Pi_{B,D}\,[\,\sigma_{R.C=S.C}\,[\sigma_{R.A=“c”}(R)]\ X\ S]$$

Apply Rewrite Rule (3)



$$\Pi_{B,D}\,[[\sigma_{R.A=\text{“c”}}(R)] \bowtie S]$$

Rules: $\sigma + \bowtie$ combined (continued)

$$\sigma_{p \wedge q} (R \bowtie S) = [\sigma_p (R)] \bowtie [\sigma_q (S)]$$

$$\sigma_{p \wedge q \wedge m} (R \bowtie S) =$$

$$\sigma_m \left[(\sigma_p R) \bowtie (\sigma_q S) \right]$$

$$\sigma_{p \vee q} (R \bowtie S) =$$

$$\left[(\sigma_p R) \bowtie S \right] \cup \left[R \bowtie (\sigma_q S) \right]$$

Which are “good” transformations?

☐ $\sigma_{p1 \wedge p2} (R) \rightarrow \sigma_{p1} [\sigma_{p2} (R)]$

☐ $\sigma_{p} (R \bowtie S) \rightarrow [\sigma_{p} (R)] \bowtie S$

☐ $R \bowtie S \rightarrow S \bowtie R$

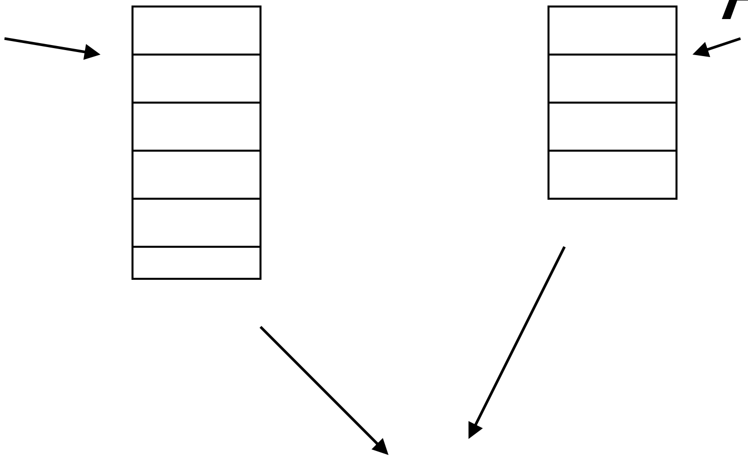
☐ $\pi_{x} [\sigma_{p} (R)] \rightarrow \pi_{x} \{\sigma_{p} [\pi_{xz} (R)]\}$

Conventional wisdom: do projects early

Example: R(A,B,C,D,E)
P: (A=3) ∧ (B="cat")

$\pi_E \{\sigma_p (R)\}$ vs. $\pi_E \left\{\sigma_p \{\pi_{ABE}(R)\}\right\}$

But: What if we have A, B indexes?



B = “cat”

A=3

Intersect pointers to get

pointers to matching tuples

Bottom line:

- No transformation is always good
- Some are usually good:
  - Push selections down
  - Avoid cross-products if possible
  - Subqueries → Joins

Avoid Cross Products (if possible)

Select B,D
From R,S,T,U
Where R.A = S.B $\wedge$
R.C=T.C $\wedge$ R.D = U.D

- Which join trees avoid cross-products?
- If you can't avoid cross products, perform them as late as possible

More Query Rewrite Rules

- Transform one logical plan into another
  - Do not use statistics
- Equivalences in relational algebra
- Push-down predicates
- Do projects early
- Avoid cross-products if possible
- Use left-deep trees
- Subqueries → Joins
- Use of constraints, e.g., uniqueness