Query Optimization

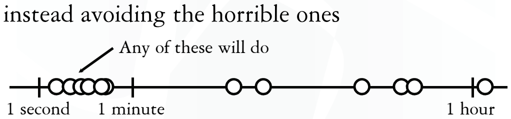
CompSci 316
Introduction to Database Systems

Announcements (Tue. Nov. 27)

- Homework #4 due in a week
- Sign up (via email) for a 30-minute slot in the project demo period, Dec. 10-12
  - "Public" demo slots available on Dec. 6
- Final exam 2-5pm Dec. 12
  - Open book, open notes
  - Focus on the second half of the course
  - Sample final available soon

Query optimization

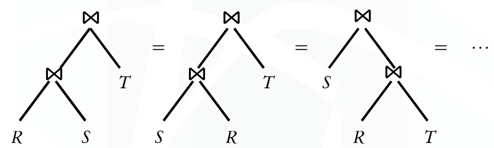
- One logical plan → "best" physical plan
- Questions
  - How to enumerate possible plans
  - How to estimate costs
  - How to pick the "best" one
- Often the goal is not getting the optimum plan, but instead avoiding the horrible ones



Plan enumeration in relational algebra

4

- Apply relational algebra equivalences
- ☞ Join reordering: × and ⋈ are associative and commutative (except column ordering, but that is unimportant)



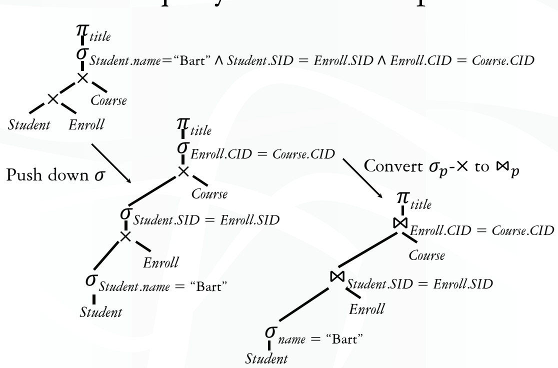
More relational algebra equivalences

5

- Convert σ_p-× to/from $\bowtie_p$: $\sigma_p(R \times S) = R \bowtie_p S$
- Merge/split σ's: $\sigma_{p_1}(\sigma_{p_2} R) = \sigma_{p_1 \wedge p_2} R$
- Merge/split π's: $\pi_{L_1}(\pi_{L_2} R) = \pi_{L_1} R$, where $L_1 \subseteq L_2$
- Push down/pull up σ:
  $\sigma_{p \wedge p_r \wedge p_s}(R \bowtie_{p'} S) = (\sigma_{p_r} R) \bowtie_{p \wedge p'} (\sigma_{p_s} S)$, where
  - p_r is a predicate involving only R columns
  - p_s is a predicate involving only S columns
  - p and p' are predicates involving both R and S columns
- Push down π: $\pi_L(\sigma_p R) = \pi_L(\sigma_p(\pi_{LL'} R))$, where
  - L' is the set of columns referenced by p that are not in L
- Many more (seemingly trivial) equivalences…
  - Can be systematically used to transform a plan to new ones

Relational query rewrite example

6



Heuristics-based query optimization

7

- Start with a logical plan
- Push selections/projections down as much as possible
  - Why?
  - Why not?
- Join smaller relations first, and avoid cross product
  - Why?
  - Why not?
- Convert the transformed logical plan to a physical plan (by choosing appropriate physical operators)

SQL query rewrite

8

- More complicated—subqueries and views divide a query into nested "blocks"
  - Processing each block separately forces particular join methods and join order
  - Even if the plan is optimal for each block, it may not be optimal for the entire query
- Unnest query: convert subqueries/views to joins
- ☞ We can just deal with select-project-join queries
  - Where the clean rules of relational algebra apply

SQL query rewrite example

9

- ```
  SELECT name
  FROM Student
  WHERE SID = ANY (SELECT SID FROM Enroll);
  ```
- ```
  SELECT name
  FROM Student, Enroll
  WHERE Student.SID = Enroll.SID;
  ```
  - Wrong—

Dealing with correlated subqueries

10

- ```
  SELECT CID FROM Course
  WHERE title LIKE 'CPS%'
  AND min_enroll > (SELECT COUNT(*) FROM Enroll
                    WHERE Enroll.CID = Course.CID);
  ```
- ```
  SELECT CID
  FROM Course, (SELECT CID, COUNT(*) AS cnt
                FROM Enroll GROUP BY CID) t
  WHERE t.CID = Course.CID AND min_enroll > t.cnt
  AND title LIKE 'CPS%';
  ```

"Magic" decorrelation

11

- ```
  SELECT CID FROM Course
  WHERE title LIKE 'CPS%'
  AND min_enroll > (SELECT COUNT(*) FROM Enroll
                    WHERE Enroll.CID = Course.CID);
  ```
- ```
  CREATE VIEW Supp_Course AS
  SELECT * FROM Course WHERE title LIKE 'CPS%';
  ```
  Process the outer query without the subquery

  ```
  CREATE VIEW Magic AS
  SELECT DISTINCT CID FROM Supp_Course;
  ```
  Collect bindings

  ```
  CREATE VIEW DS AS
  (SELECT Enroll.CID, COUNT(*) AS cnt
   FROM Magic, Enroll WHERE Magic.CID = Enroll.CID
   GROUP BY Enroll.CID) UNION
  (SELECT Magic.CID, 0 AS cnt FROM Magic
   WHERE Magic.CID NOT IN (SELECT CID FROM Enroll));
  ```
  Evaluate the subquery with bindings

  ```
  SELECT Supp_Course.CID FROM Supp_Course, DS
  WHERE Supp_Course.CID = DS.CID
  AND min_enroll > DS.cnt;
  ```
  Finally, refine the outer query

Heuristics- vs. cost-based optimization

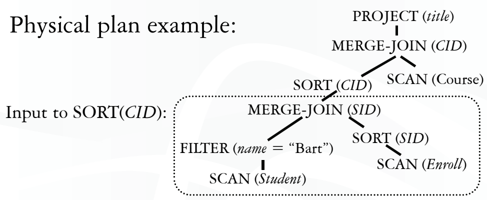
12

- Heuristics-based optimization
  - Apply heuristics to rewrite plans into cheaper ones
- Cost-based optimization
  - Rewrite logical plan to combine "blocks" as much as possible
  - Optimize query block by block
    - Enumerate logical plans (already covered)
    - Estimate the cost of plans
    - Pick a plan with acceptable cost
  - Focus: select-project-join blocks

Cost estimation

13

Physical plan example:



- We have: cost estimation for each operator
  - Example: SORT(CID) takes $O(B(\text{input}) \times \log_M B(\text{input}))$
    - But what is $B(\text{input})$?
- We need: size of intermediate results

Selections with equality predicates

14

- Q: $\sigma_{A=v} R$
- Suppose the following information is available
  - Size of R: $|R|$
  - Number of distinct A values in R: $|\pi_A R|$
- Assumptions
  - Values of A are uniformly distributed in R
  - Values of v in Q are uniformly distributed over all $R.A$ values
- $|Q| \approx {}^{|R|}/_{|\pi_A R|}$
  - Selectivity factor of $(A = v)$ is ${}^{1}/_{|\pi_A R|}$

Conjunctive predicates

15

- Q: $\sigma_{A=u \wedge B=v} R$
- Additional assumptions
  - $(A = u)$ and $(B = v)$ are independent
    - Counterexample: major and advisor
  - No "over"-selection
    - Counterexample: A is the key

Negated and disjunctive predicates

16

- Q: $\sigma_{A \neq v} R$
  - $|Q| \approx |R| \cdot (1 - 1/|\pi_A R|)$
    - Selectivity factor of $\neg p$ is $(1 -$ selectivity factor of $p)$
- Q: $\sigma_{A=u \vee B=v} R$
  - $|Q| \approx |R| \cdot (1/|\pi_A R| + 1/|\pi_B R|)$?

Range predicates

17

- Q: $\sigma_{A>v} R$
- Not enough information!
  - Just pick, say, $|Q| \approx |R| \cdot 1/3$
- With more information
  - Largest $R.A$ value: $\text{high}(R.A)$
  - Smallest $R.A$ value: $\text{low}(R.A)$
  - $|Q| \approx |R| \cdot \frac{\text{high}(R.A) - v}{\text{high}(R.A) - \text{low}(R.A)}$
  - In practice: sometimes the second highest and lowest are used instead

Two-way equi-join

18

- Q: $R(A, B) \bowtie S(A, C)$
- Assumption: containment of value sets
  - Every tuple in the "smaller" relation (one with fewer distinct values for the join attribute) joins with some tuple in the other relation
  - That is, if $|\pi_A R| \leq |\pi_A S|$ then $\pi_A R \subseteq \pi_A S$
  - Certainly not true in general
  - But holds in the common case of foreign key joins
- $|Q| \approx \frac{|R| \cdot |S|}{\max(|\pi_A R|, |\pi_A S|)}$
  - Selectivity factor of $R.A = S.A$ is $1/\max(|\pi_A R|, |\pi_A S|)$

Multiway equi-join

19

- Q: $R(A,B) \bowtie S(B,C) \bowtie T(C,D)$
- What is the number of distinct C values in the join of R and S?
- Assumption: preservation of value sets
  - A non-join attribute does not lose values from its set of possible values
  - That is, if A is in R but not S, then $\pi_A(R \bowtie S) = \pi_A R$
  - Certainly not true in general
  - But holds in the common case of foreign key joins (for value sets from the referencing table)

Multiway equi-join (cont'd)

20

- Q: $R(A,B) \bowtie S(B,C) \bowtie T(C,D)$
- Start with the product of relation sizes
  - $|R| \cdot |S| \cdot |T|$
- Reduce the total size by the selectivity factor of each join predicate
  - $R.B = S.B$: ${}^{1}\!/_{\max(|\pi_B R|, |\pi_B S|)}$
  - $S.C = T.C$: ${}^{1}\!/_{\max(|\pi_C S|, |\pi_C T|)}$
  - $|Q| \approx \frac{|R| \cdot |S| \cdot |T|}{\max(|\pi_B R|, |\pi_B S|) \cdot \max(|\pi_C S|, |\pi_C T|)}$

Cost estimation: summary

21

- Using similar ideas, we can estimate the size of projection, duplicate elimination, union, difference, aggregation (with grouping)
- Lots of assumptions and very rough estimation
  - Accurate estimate is not needed
  - Maybe okay if we overestimate or underestimate consistently
  - May lead to very nasty optimizer "hints"

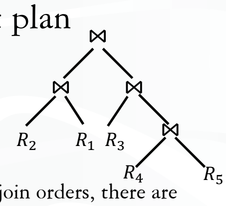
```
SELECT * FROM Student WHERE GPA > 3.9;
SELECT * FROM Student WHERE GPA > 3.9 AND GPA > 3.9;
```

- Not covered: better estimation using histograms

22

Search for the best plan

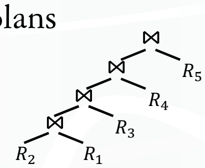
- ❖ Huge search space
- ❖ "Bushy" plan example:



- ❖ Just considering different join orders, there are $\frac{(2n-2)!}{(n-1)!}$ bushy plans for $R_1 \bowtie \cdots \bowtie R_n$
  - ▪ 30240 for $n = 6$
- ❖ And there are more if we consider:
  - ▪ Multiway joins
  - ▪ Different join methods
  - ▪ Placement of selection and projection operators

23

Left-deep plans

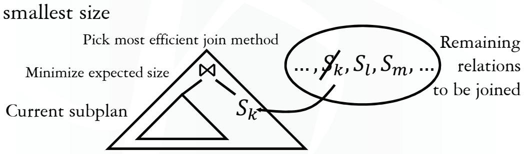


- ❖ Heuristic: consider only "left-deep" plans, in which only the left child can be a join
  - ▪ Tend to be better than plans of other shapes, because many join algorithms scan inner (right) relation multiple times—you will not want it to be a complex subtree
- ❖ How many left-deep plans are there for $R_1 \bowtie \cdots \bowtie R_n$?

24

A greedy algorithm

- ❖ $S_1, \ldots, S_n$
  - ▪ Say selections have been pushed down; i.e., $S_i = \sigma_p(R_i)$
- ❖ Start with the pair S_i, S_j with the smallest estimated size for $S_i \bowtie S_j$
- ❖ Repeat until no relation is left:
  Pick S_k from the remaining relations such that the join of S_k and the current result yields an intermediate result of the smallest size



A dynamic programming approach

25

- Generate optimal plans bottom-up
  - Pass 1: Find the best single-table plans (for each table)
  - Pass 2: Find the best two-table plans (for each pair of tables) by combining best single-table plans
  - …
  - Pass k: Find the best k-table plans (for each combination of k tables) by combining two smaller best plans found in previous passes
  - …
- Rationale: Any subplan of an optimal plan must also be optimal (otherwise, just replace the subplan to get a better overall plan)
- ☞ Well, not quite…

The need for "interesting order"

26

- Example: $R(A, B) \bowtie S(A, C) \bowtie T(A, D)$
- Best plan for $R \bowtie S$: hash join (beats sort-merge join)
- Best overall plan: sort-merge join R and S, and then sort-merge join with T
  - Subplan of the optimal plan is not optimal!
- Why?
  - The result of the sort-merge join of R and S is sorted on A
  - This is an interesting order that can be exploited by later processing (e.g., join, dup elimination, `GROUP BY`, `ORDER BY`, etc.)!

Dealing with interesting orders

27

- When picking the best plan
  - Comparing their costs is not enough
    - Plans are not totally ordered by cost anymore
  - Comparing interesting orders is also needed
    - Plans are now partially ordered
    - Plan X is better than plan Y if
      - Cost of X is lower than Y, and
      - Interesting orders produced by X "subsume" those produced by Y
- Need to keep a set of optimal plans for joining every combination of k tables
  - At most one for each interesting order

28

Summary

- Relational algebra equivalence
- SQL rewrite tricks
- Heuristics-based optimization
- Cost-based optimization
  - Need statistics to estimate sizes of intermediate results
  - Greedy approach
  - Dynamic programming approach