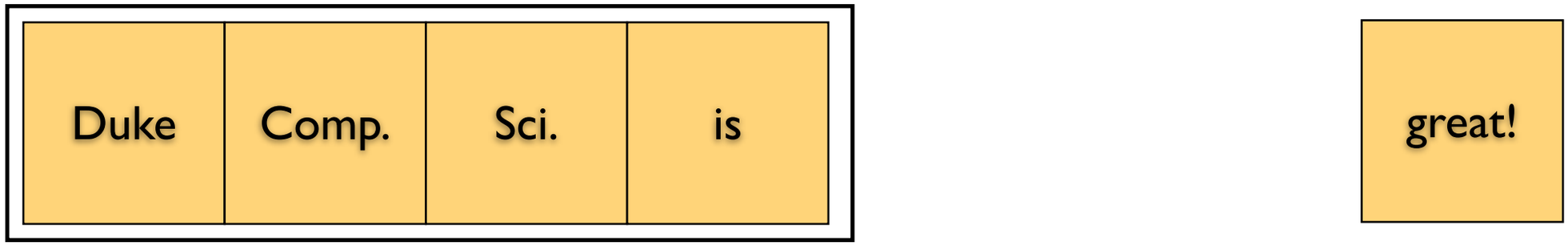
Amortized Analysis

(or: How ArrayLists work)

Recall ExpandingArray



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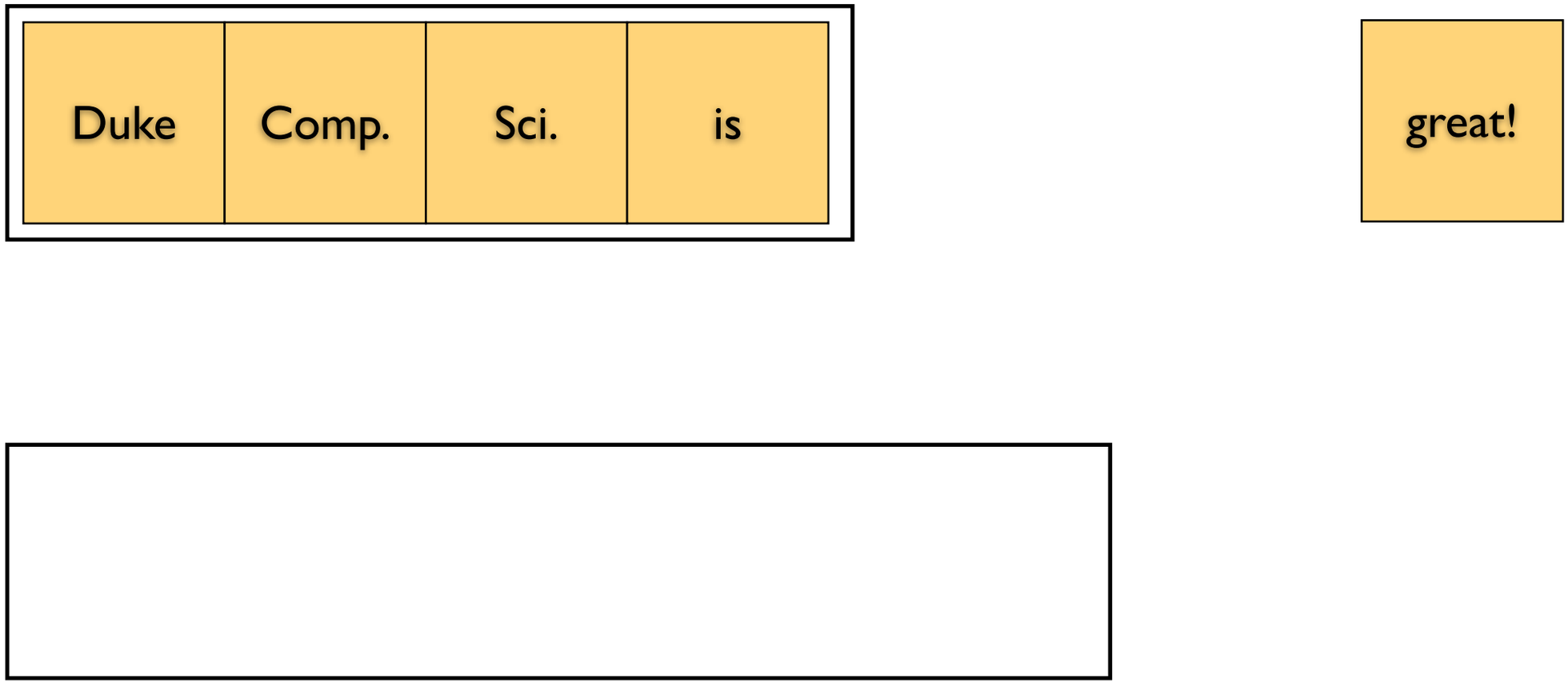
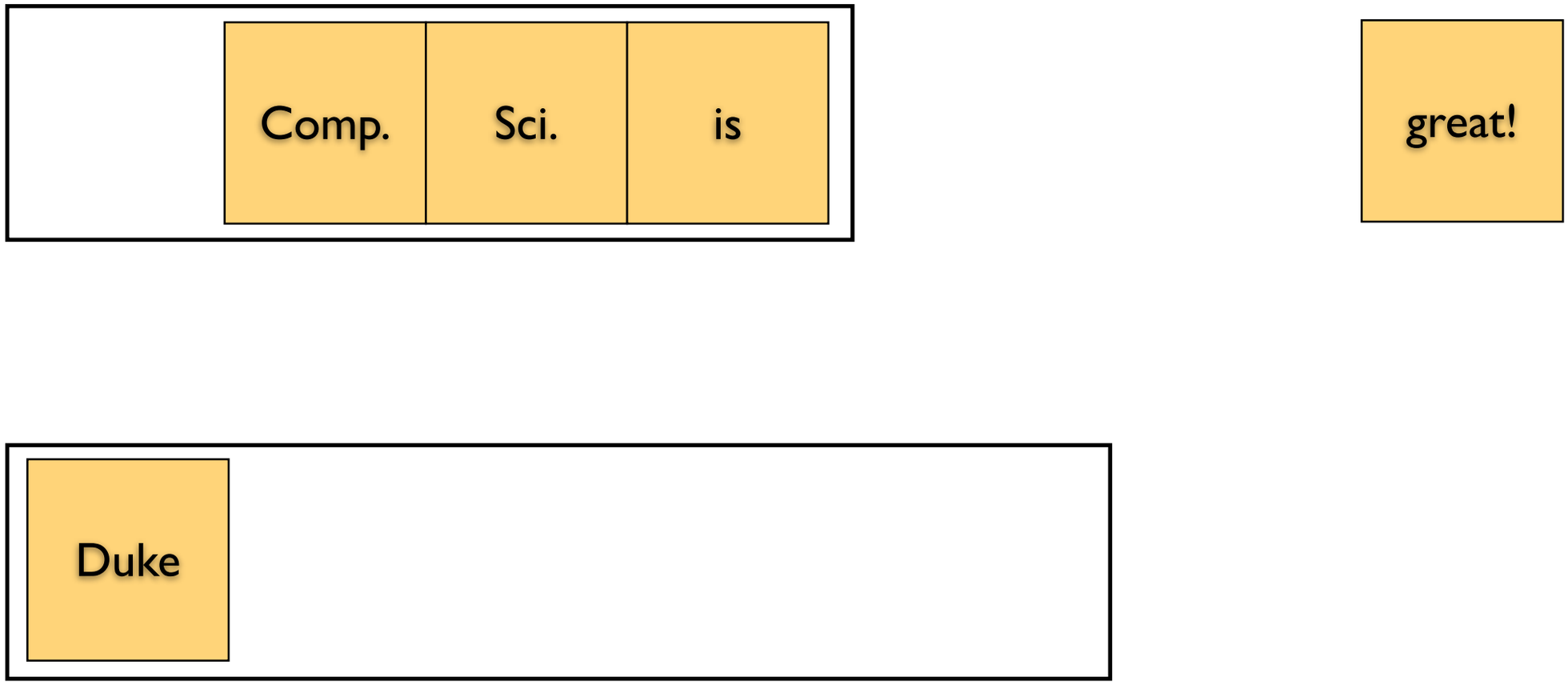
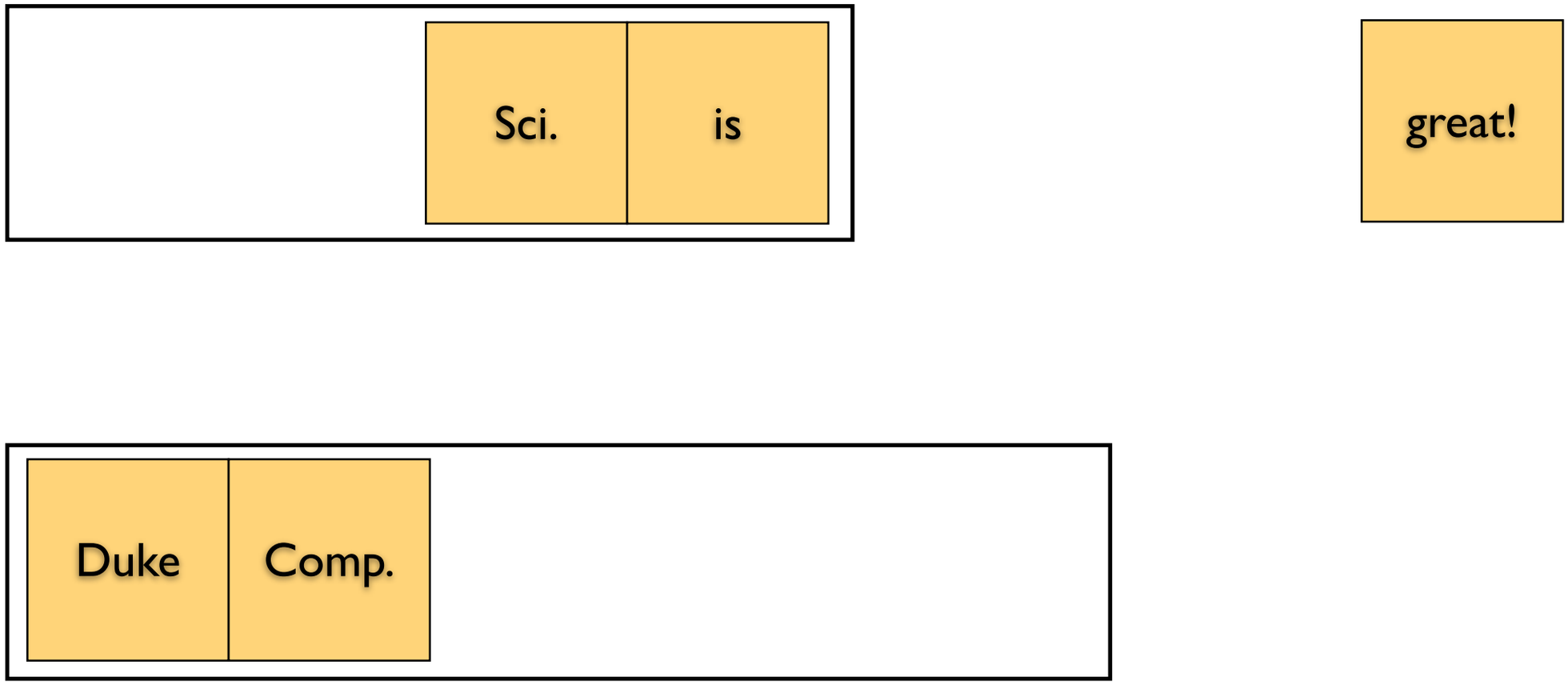


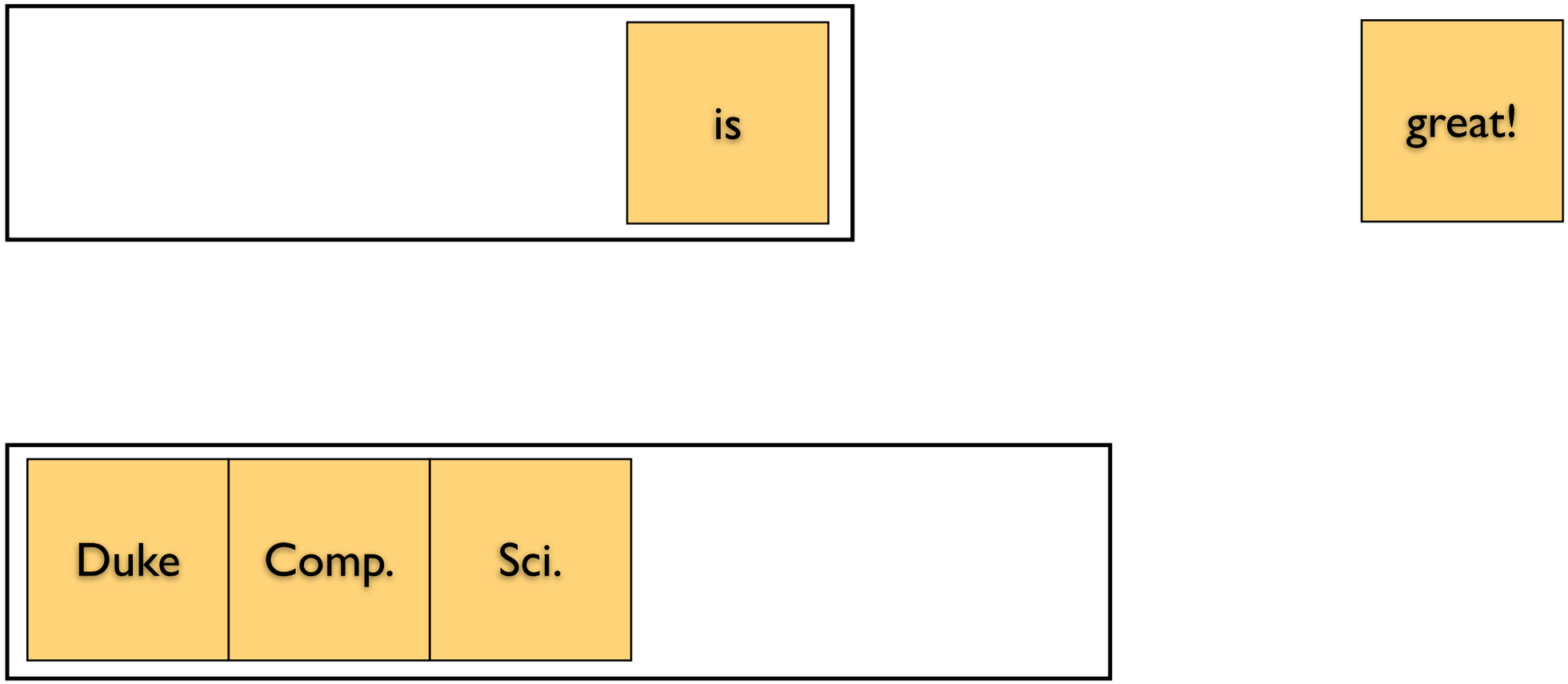
Recall ExpandingArray



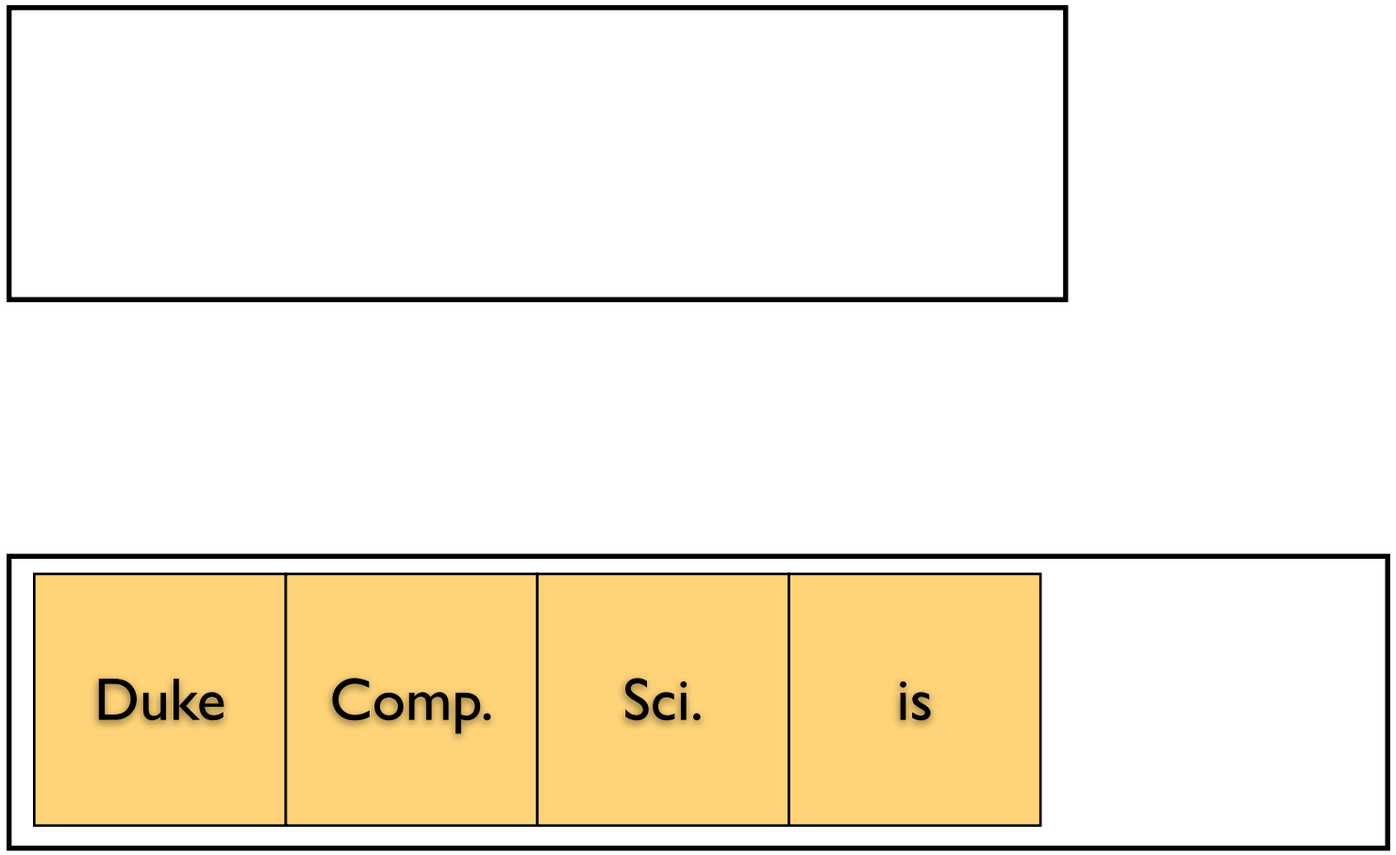
Recall ExpandingArray



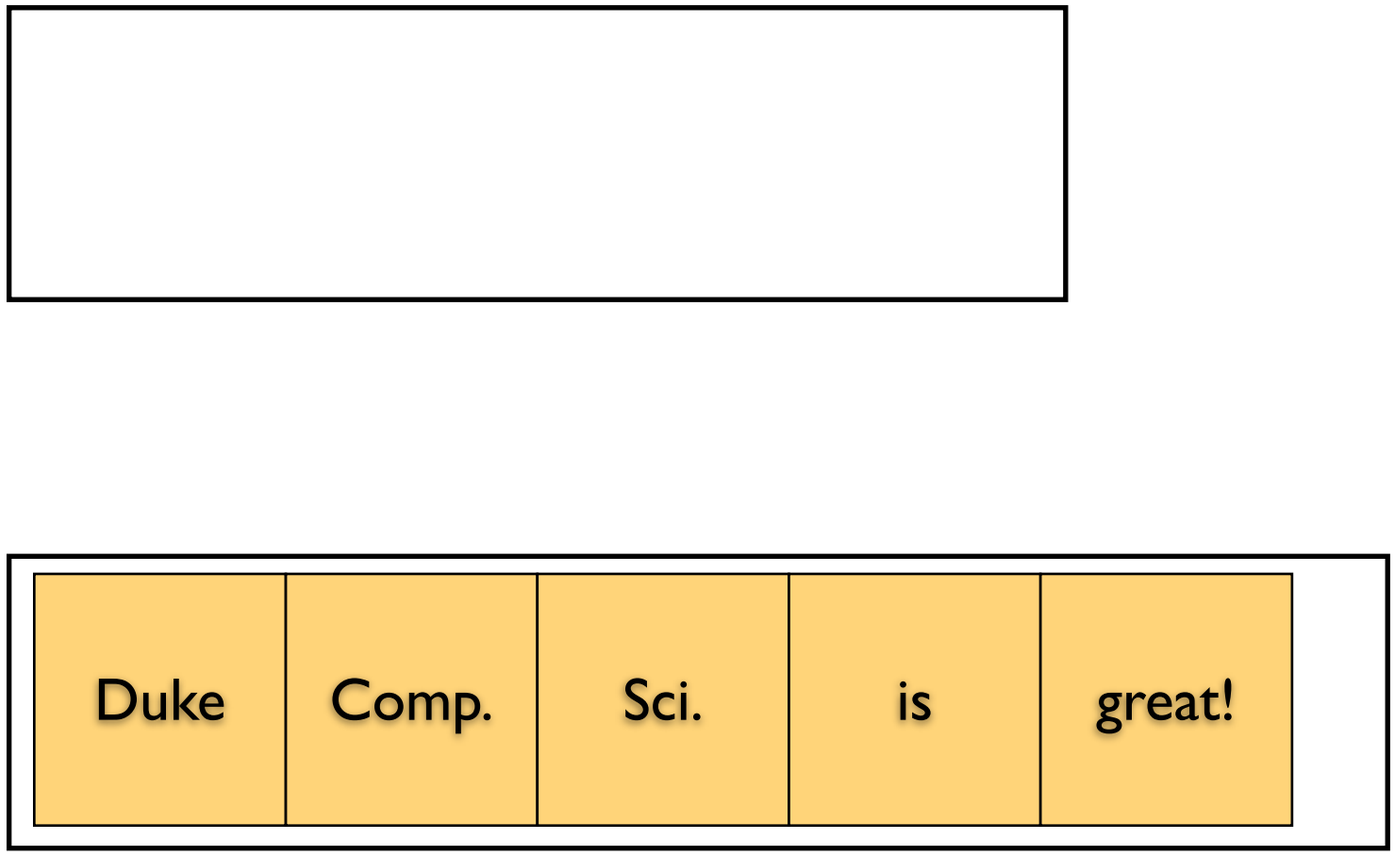
Recall ExpandingArray



Recall ExpandingArray



Recall ExpandingArray



Adding one each time leads to $\frac{n(n+1)}{2} \in O(n^2)$

Why not Linked Lists?

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`ExpandingArray`

.get(): *O(1)* *You couldn't hope for better!*

.add(): *O(n)* *Which means O(n^2) for n operations...*

Linked List *Re: DNA: Good at splicing, too!*

.get(): *O(n)* *Which means O(n^2) for n operations...*

.add(): *O(1)* *Best it can be!*

What we want:

.get(): *O(1)* *Best it can be!*

.add(): *O(1)* *Best it can be!*

It can be done!

ArrayList & StringBuilder, for example

Also: StringBuffer, C++'s vector, and Python's list. *Not* Matlab's array.

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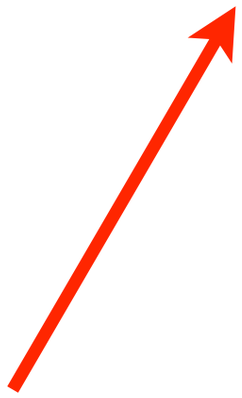
ArrayList & StringBuilder, for example



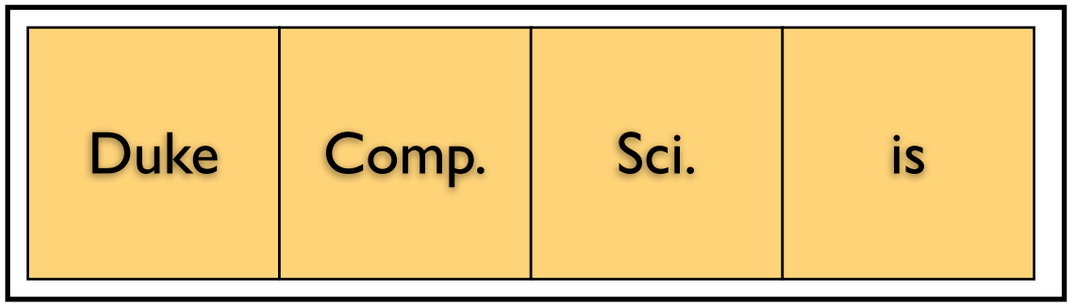
What we want:

.get(): *O(1)* *Best it can be!*

.add(): *O(1)* *Best it can be!*

Also: StringBuffer, C++'s vector, and Python's list. *Not* Matlab's array.

Backed by an array!

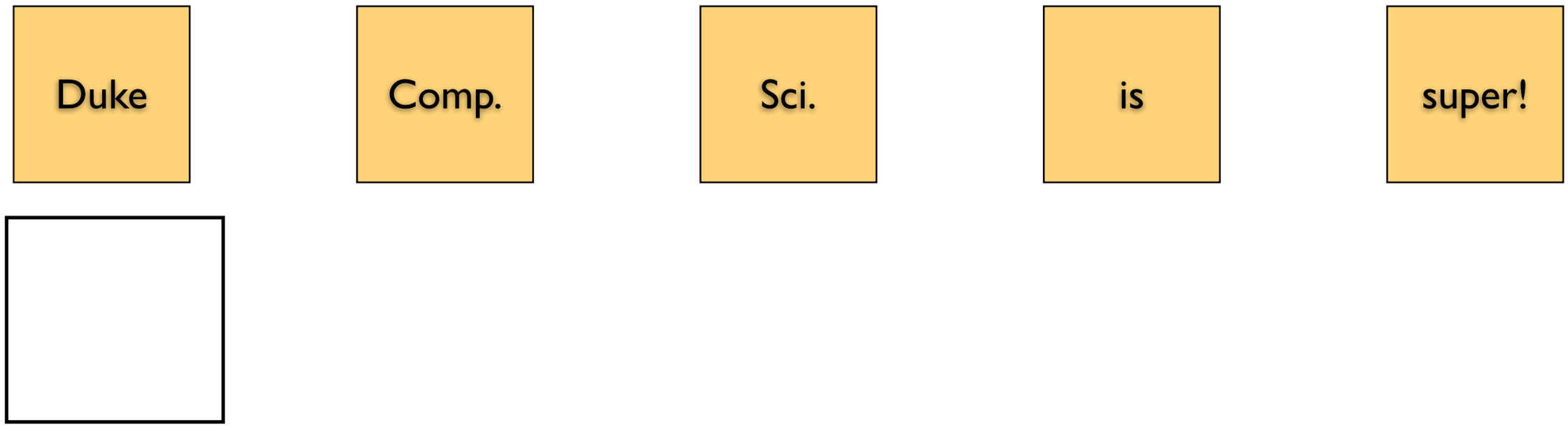


What we want:

.get(): *O(1)* *Best it can be!*

.add(): *O(...)* *...?*

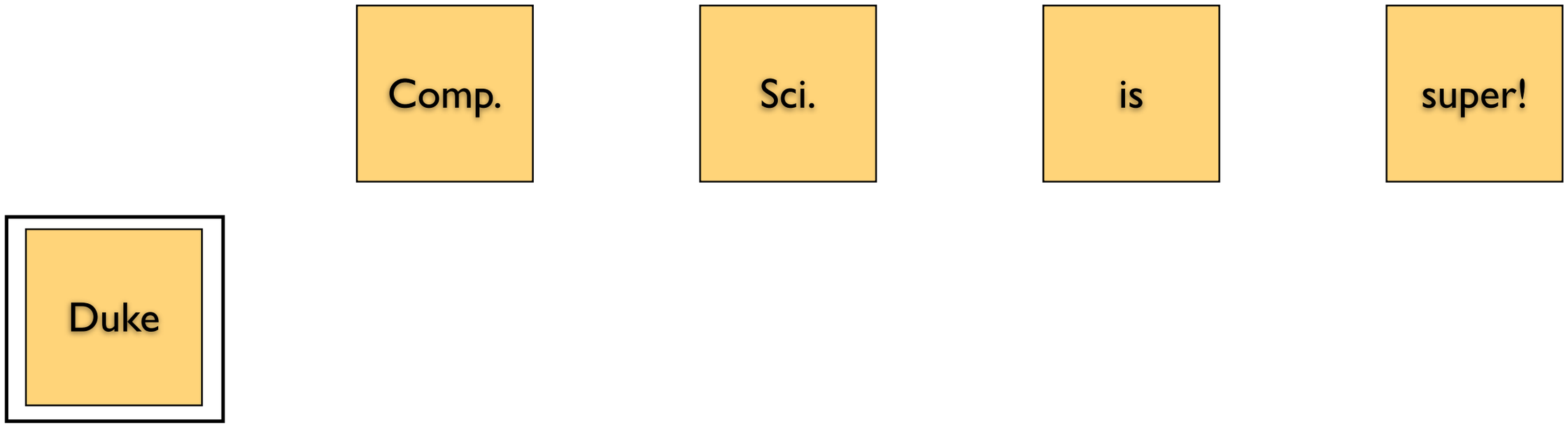
Add more than one?



Operations: 0

Adds: 0

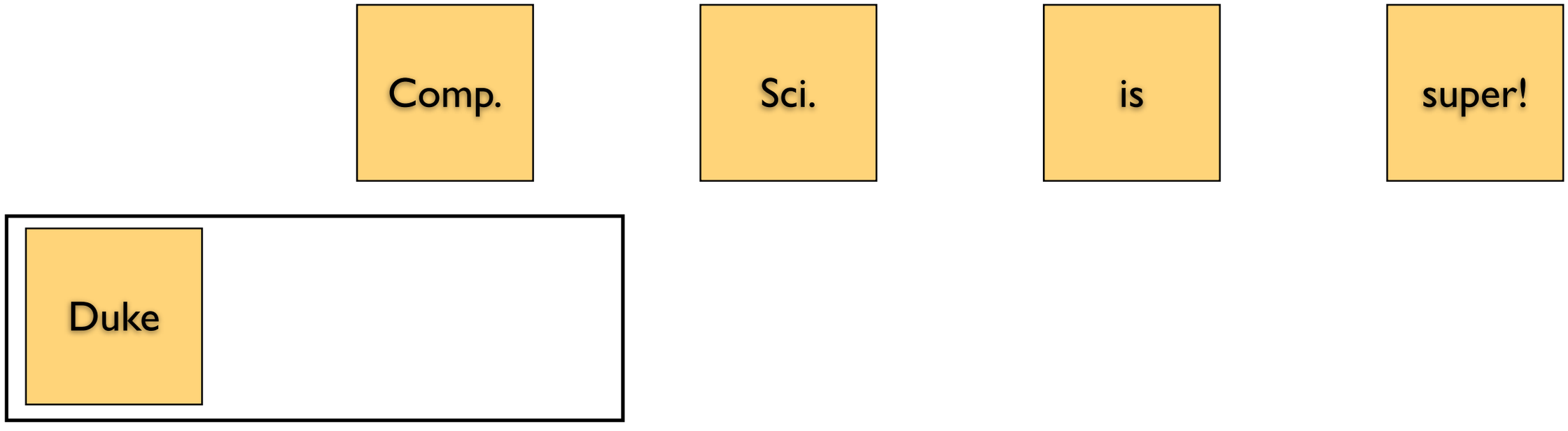
Add more than one?



Operations: 1

Adds: 1

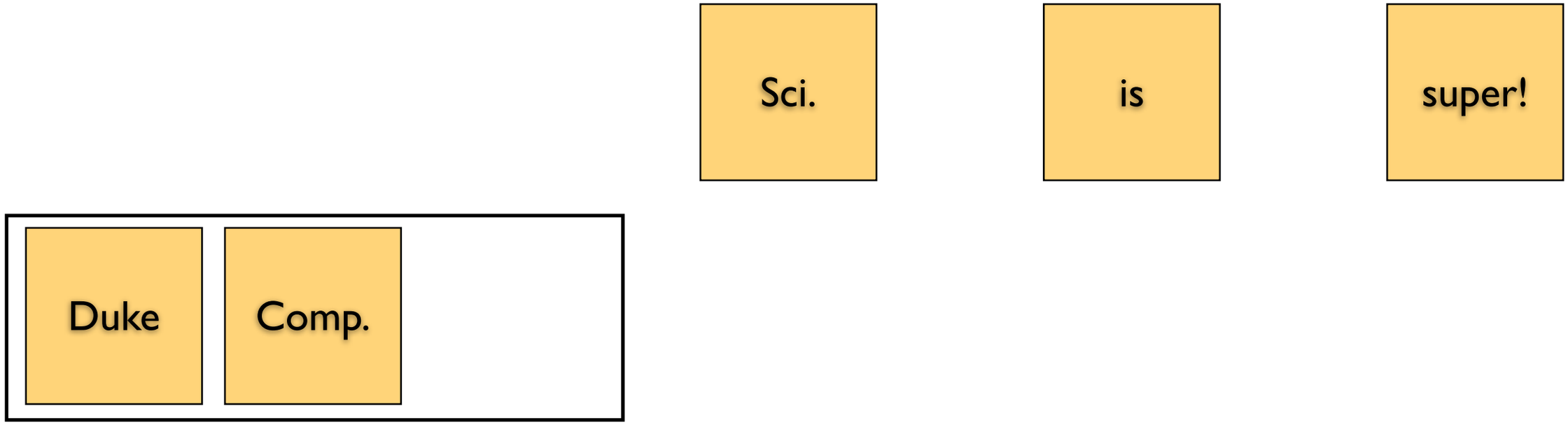
Add more than one?



Operations: 2

Adds: 1

Add more than one?



Operations: 3

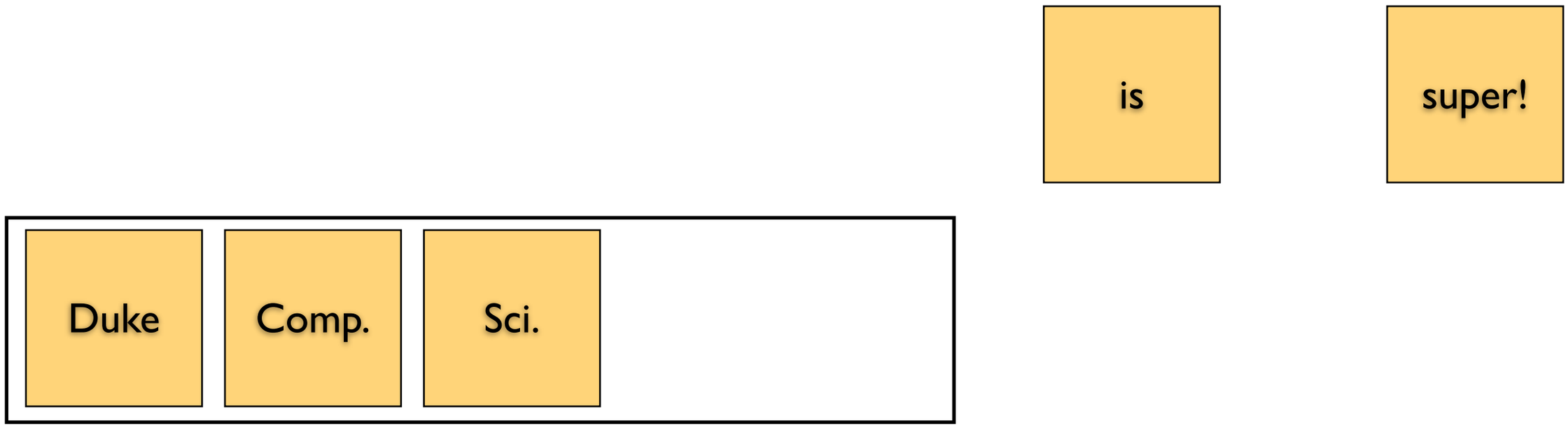
Adds: 2

Add more than one?

Operations: 4

Adds: 3

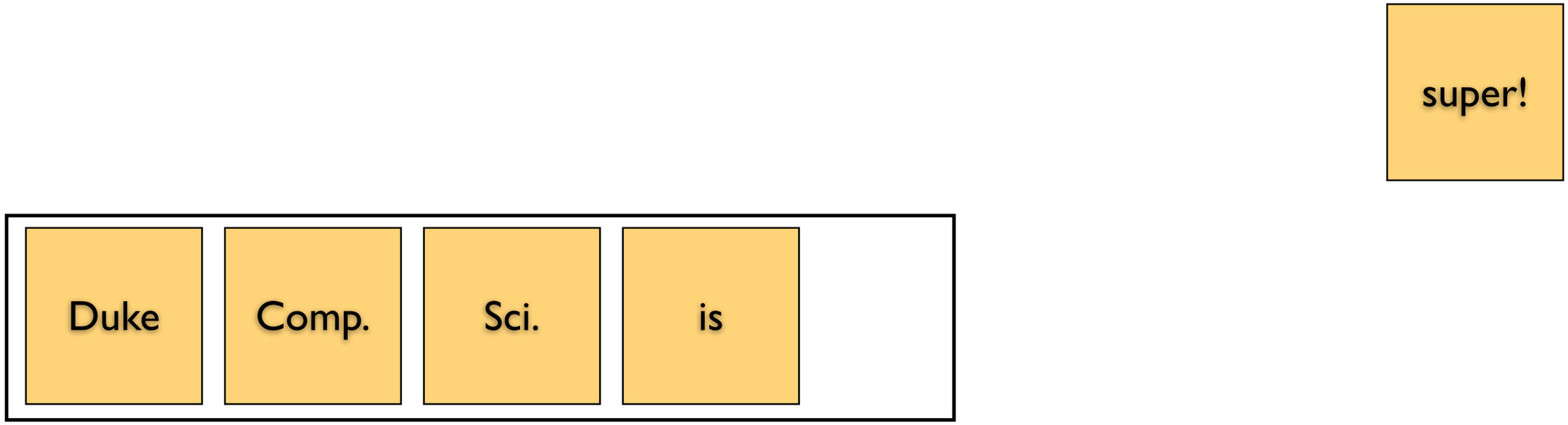
Add more than one?



Operations: 7

Adds: 3

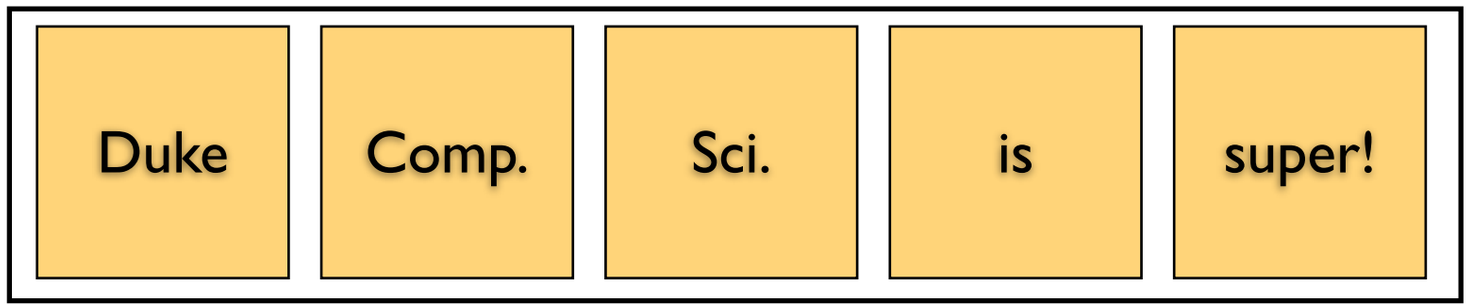
Add more than one?



Operations: 8

Adds: 4

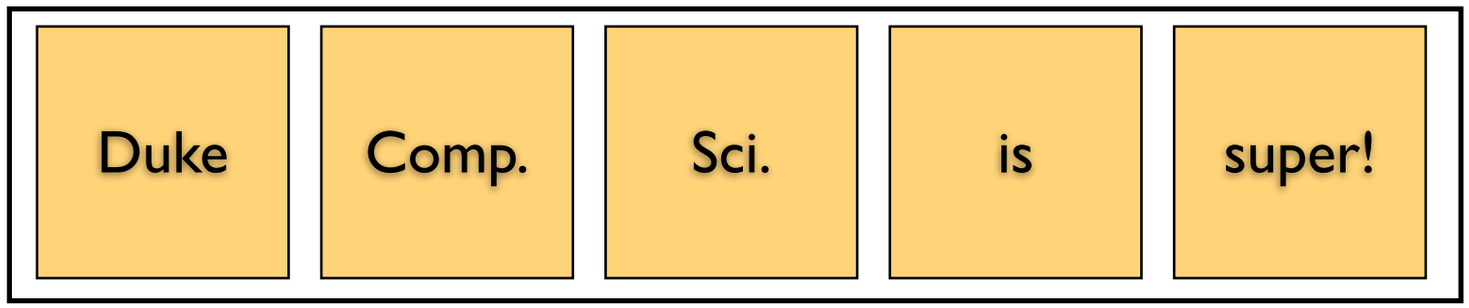
Add more than one?



Operations: 9

Adds: 5

Add more than one?



Operations: 9

Adds: 5

http://goo.gl/eLp8l

Adding *n* elements, expanding by *k*

O(1) most of the time

O(n) sometimes

$$\sum_{i=1}^{n} \left[1 + \frac{i}{k} \right]$$

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$$\sum_{i=1}^{n} \left[1 + \frac{i}{k} \right] = n + \sum_{i=1}^{n} \frac{i}{k} = n + \frac{1}{k} \sum_{i=1}^{n} i$$

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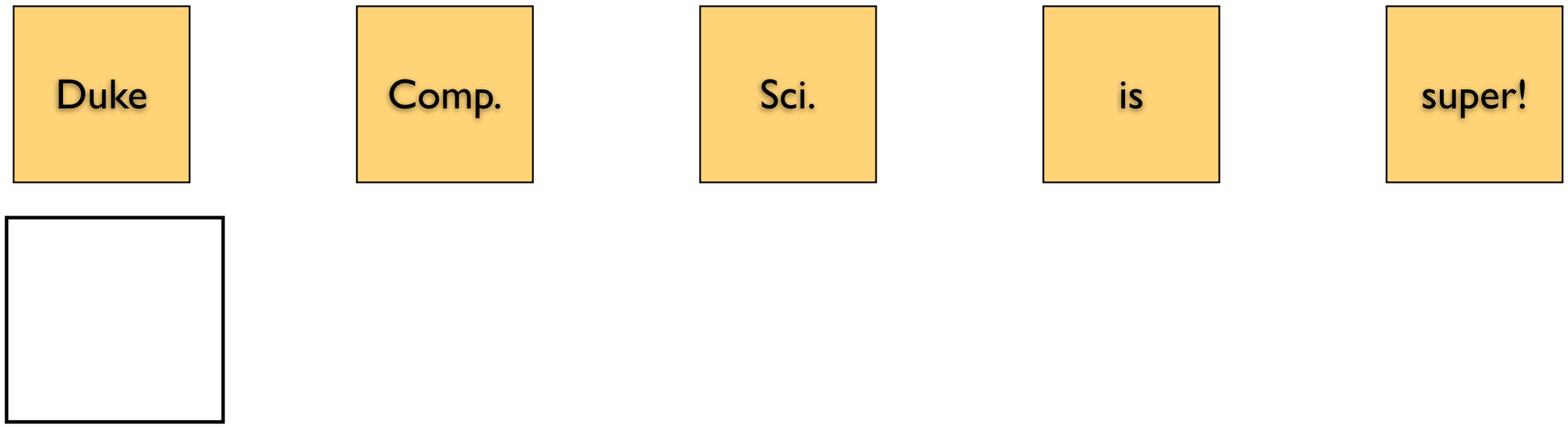
$$\sum_{i=1}^{n} \left[1 + \frac{i}{k}\right] = n + \sum_{i=1}^{n} \frac{i}{k} = n + \frac{1}{k} \sum_{i=1}^{n} i$$

Note the worst-case behavior still dominates!

$$\in O(n^2)$$

Uh-oh...

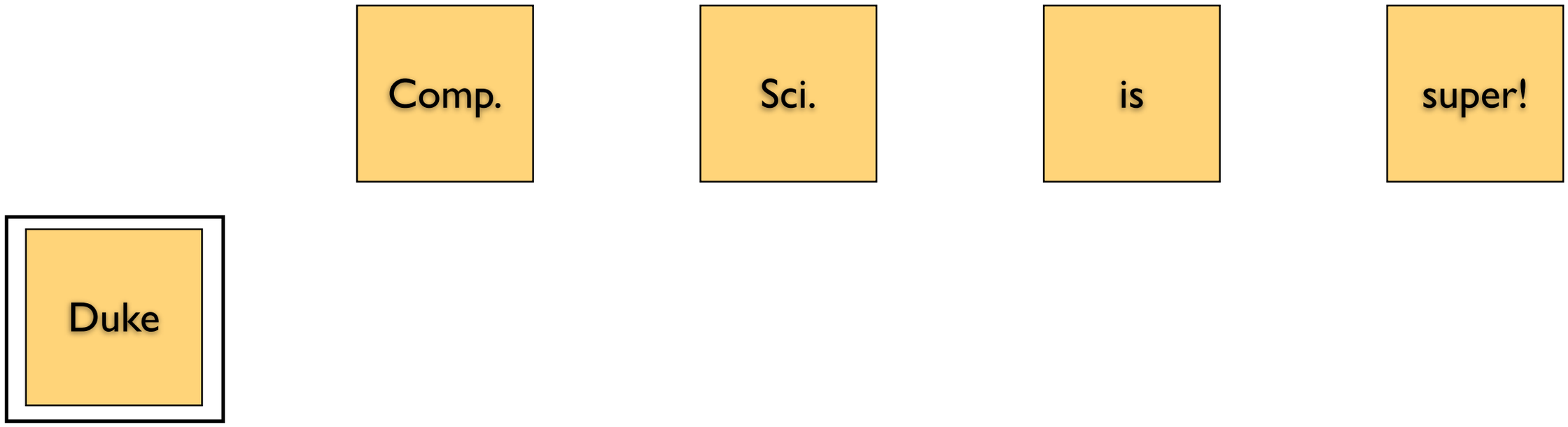
Adding didn't work...



Operations: 0

Adds: 0

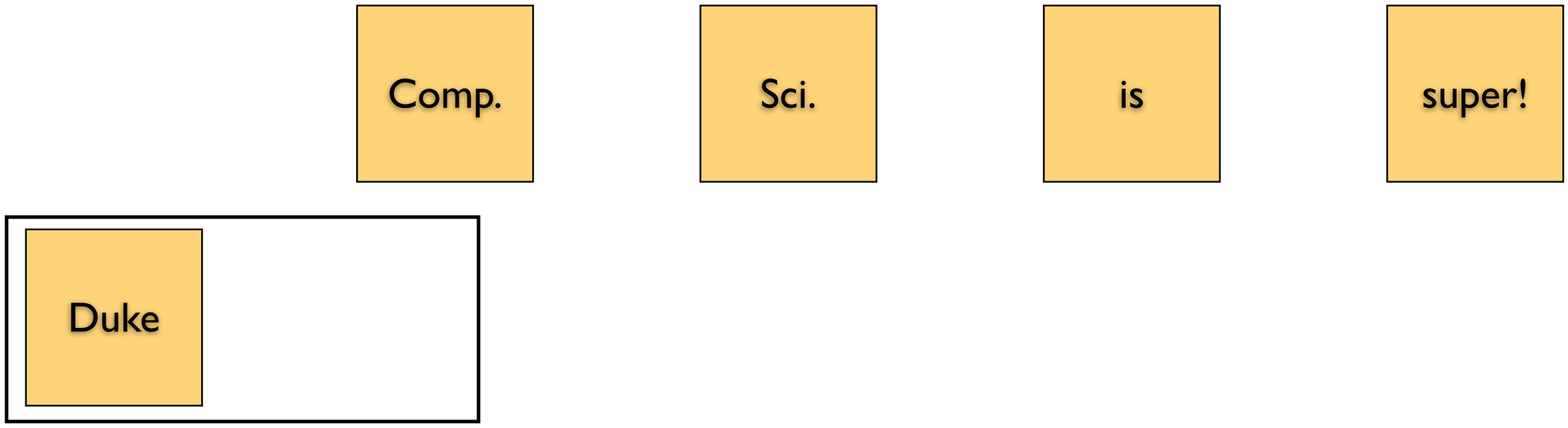
Adding didn't work...



Operations: 1

Adds: 1

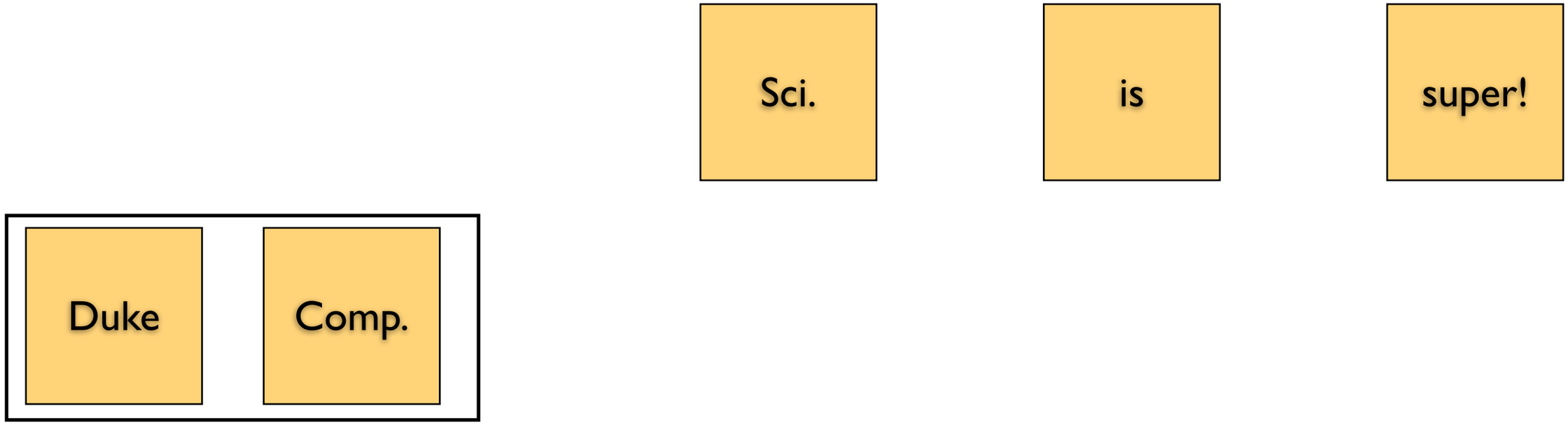
Adding didn't work...



Operations: 2

Adds: 1

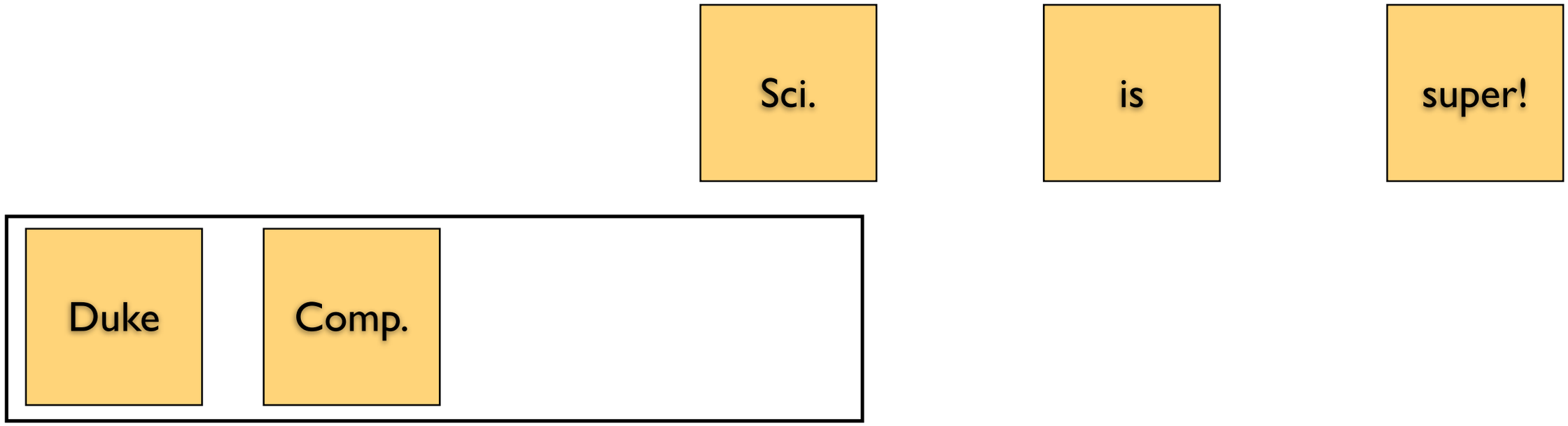
Adding didn't work...



Operations: 3

Adds: 2

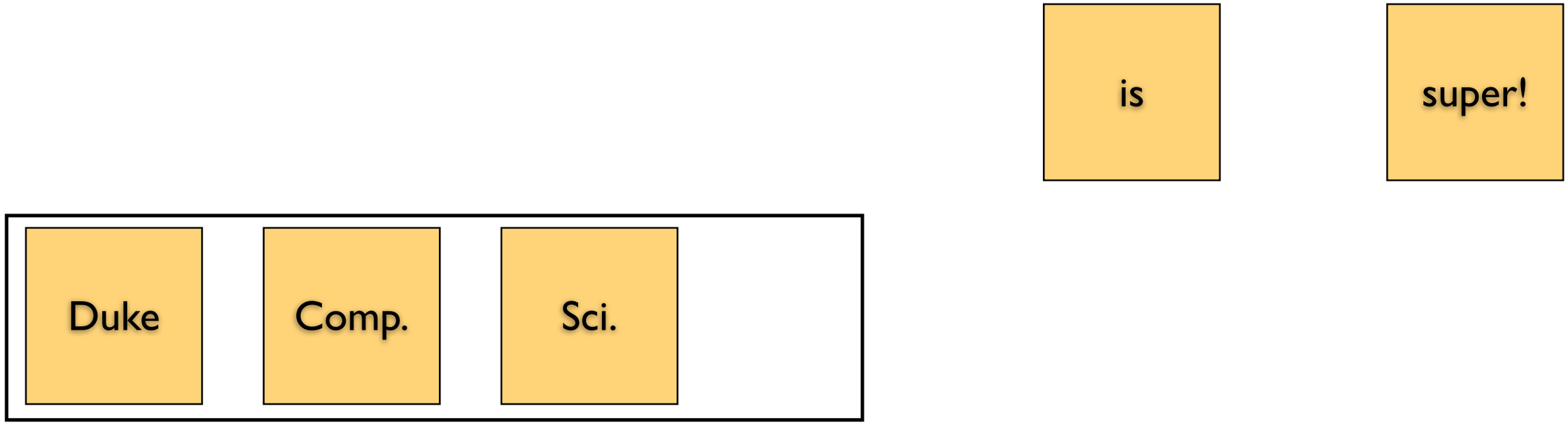
Adding didn't work...



Operations: 5

Adds: 2

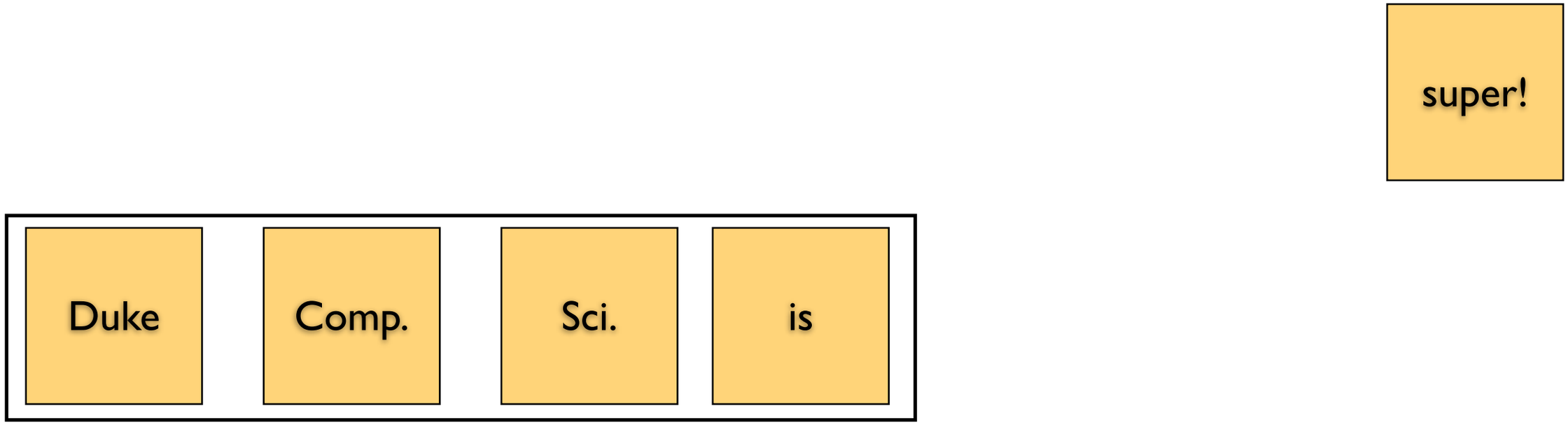
Adding didn't work...



Operations: 6

Adds: 3

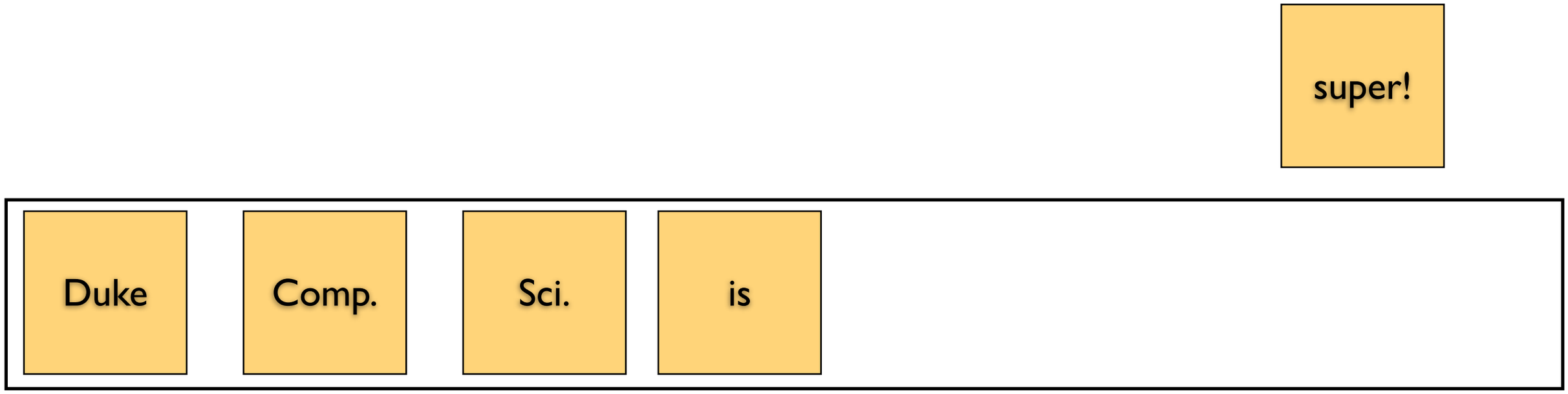
Adding didn't work...



Operations: 7

Adds: 4

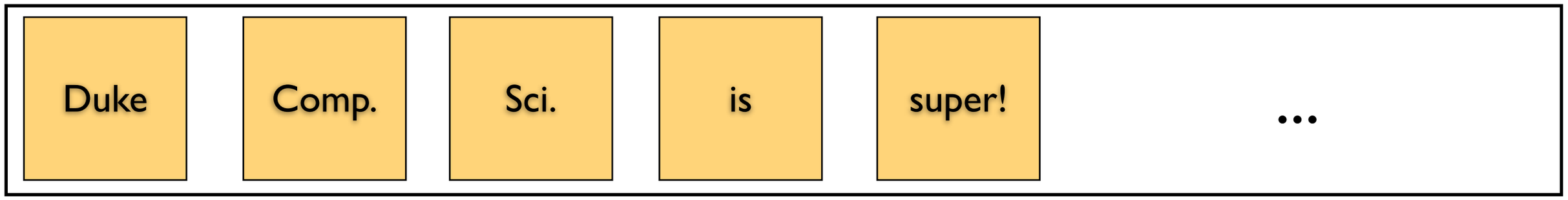
Adding didn't work...



Operations: 11

Adds: 4

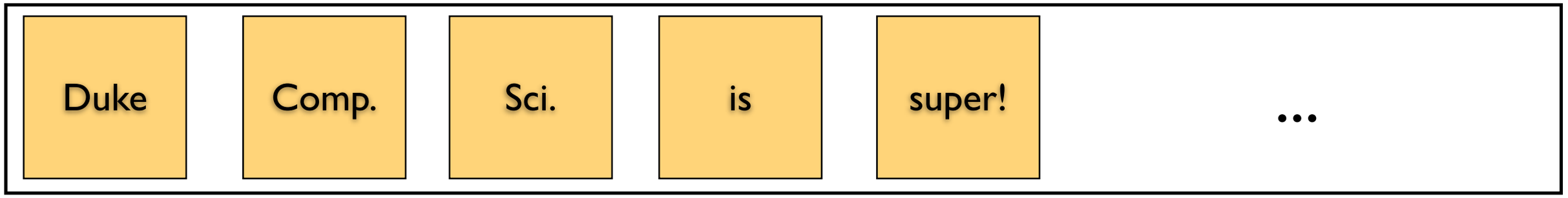
Adding didn't work...



Operations: 12

Adds: 5

Adding didn't work...



How expensive is this strategy?

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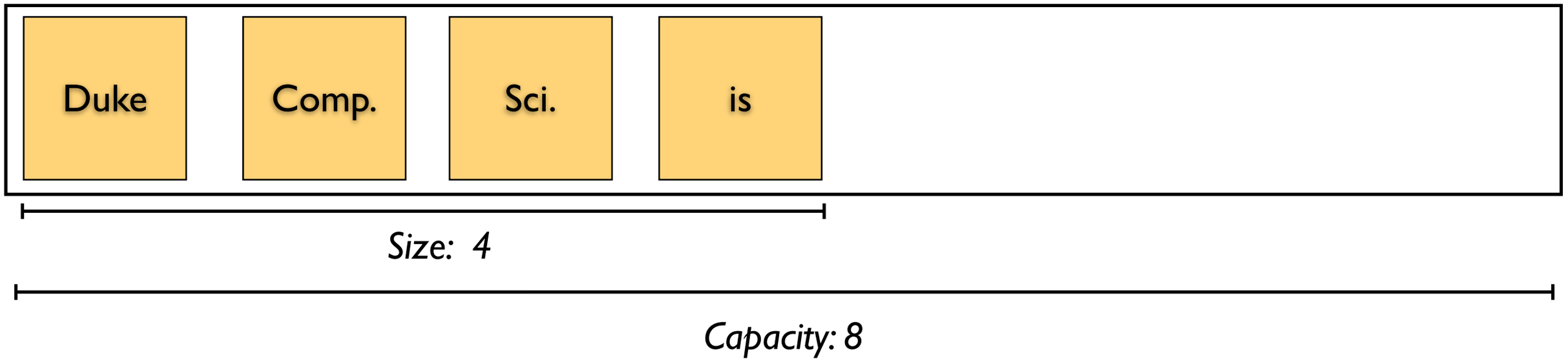
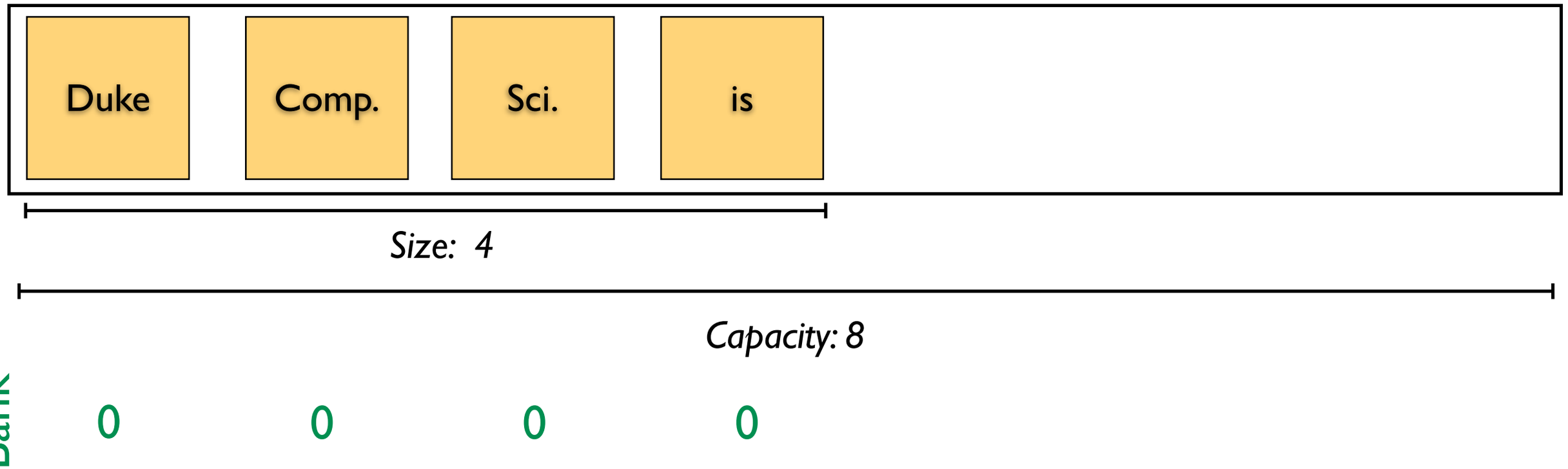
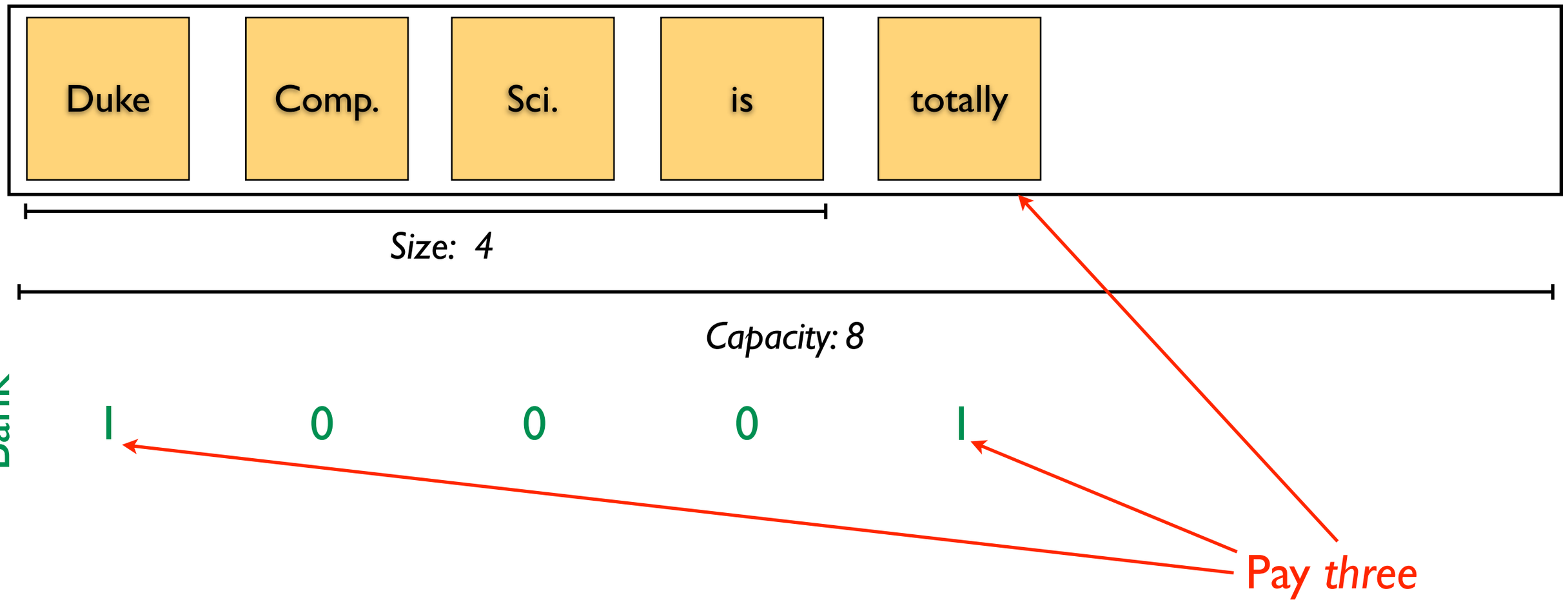


Suppose you've just doubled your array.

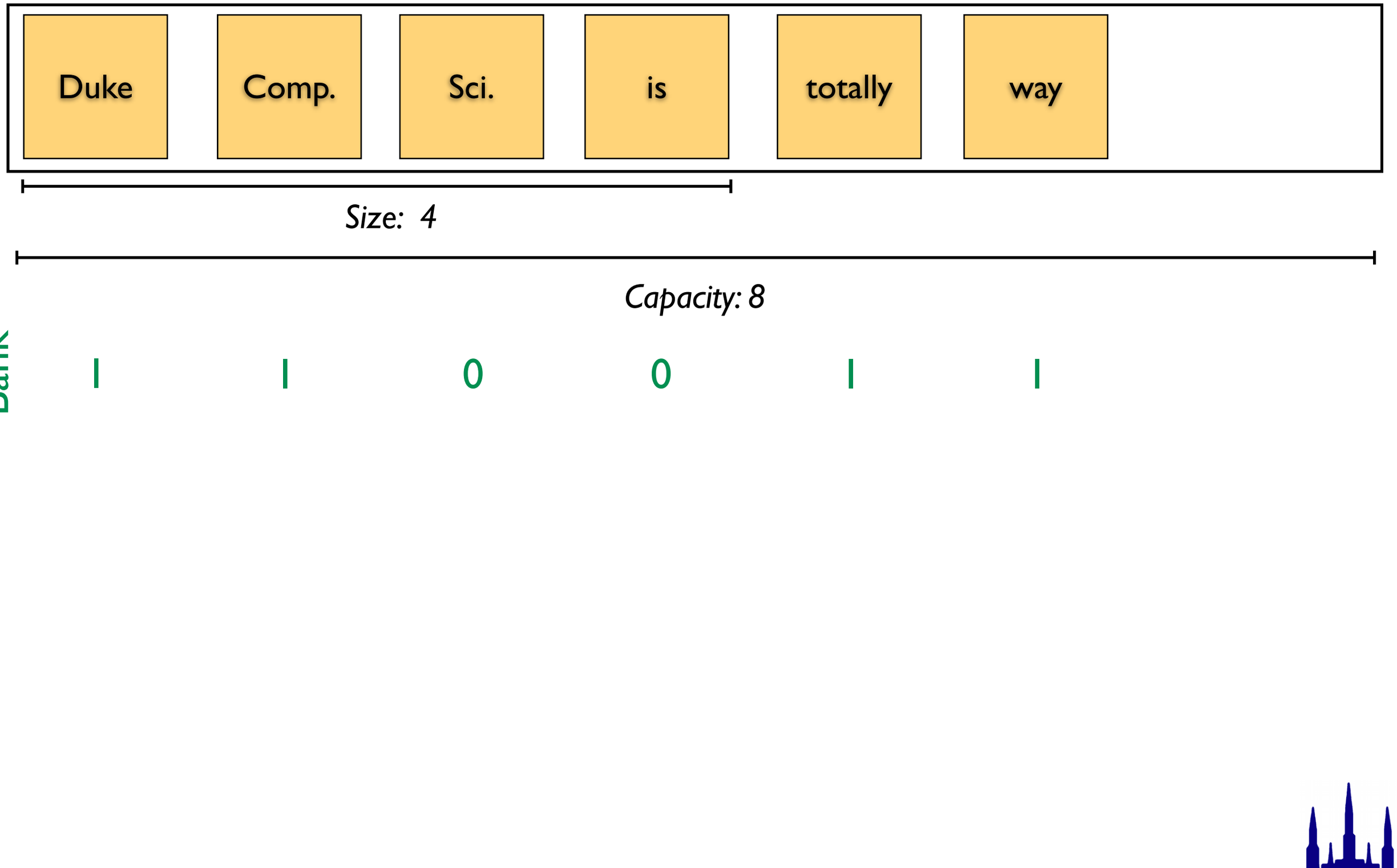
How expensive is this strategy?



How expensive is this strategy?



How expensive is this strategy?



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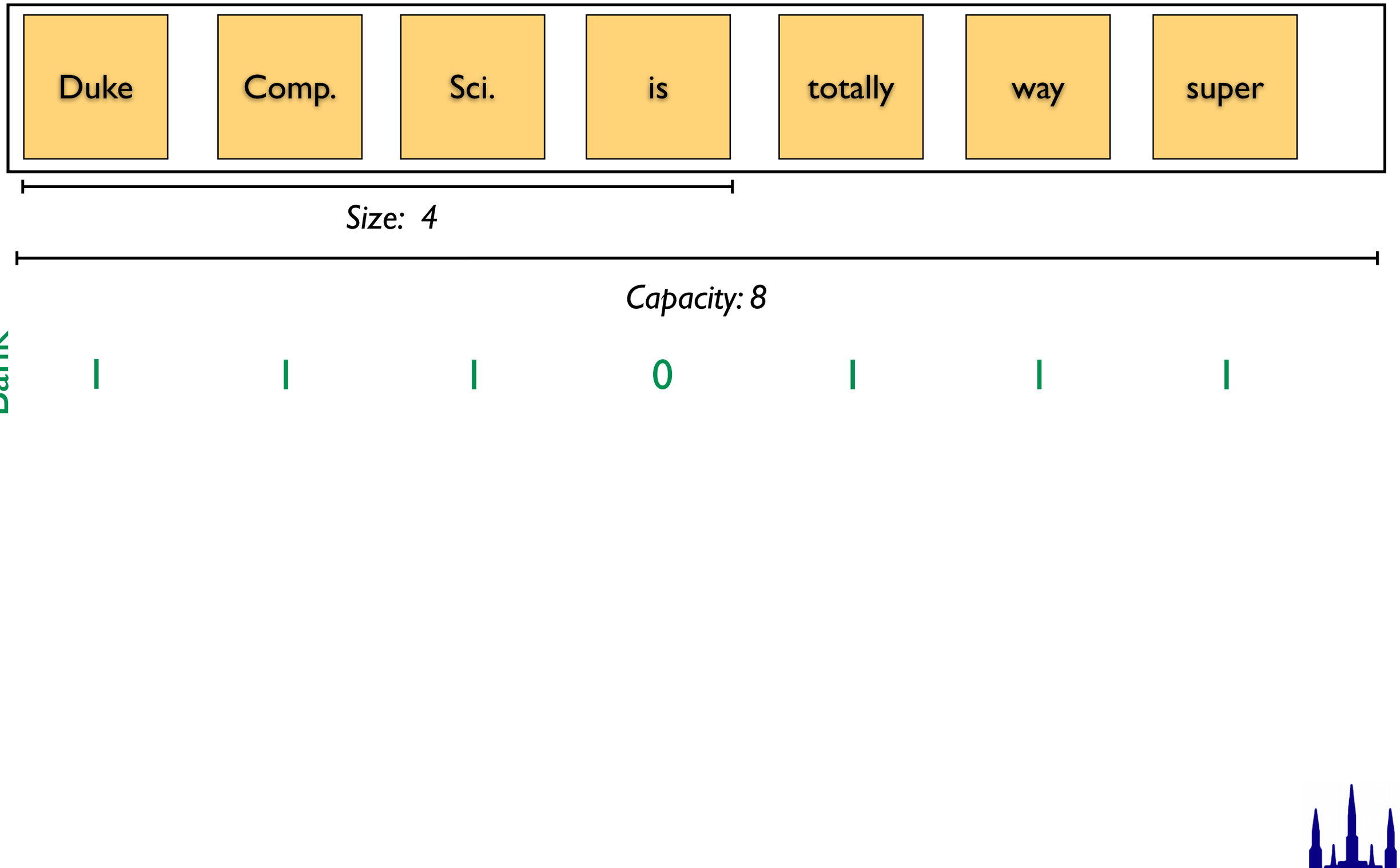

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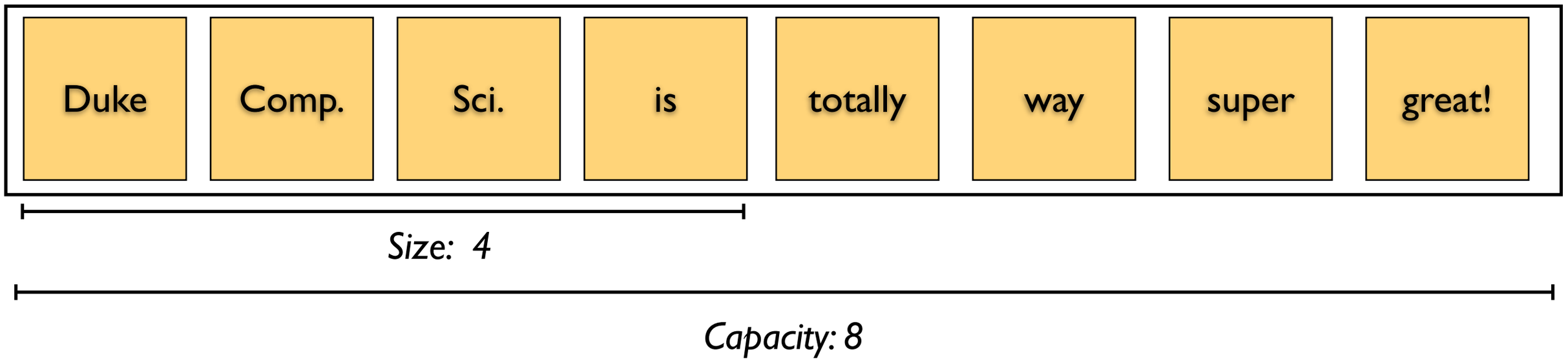
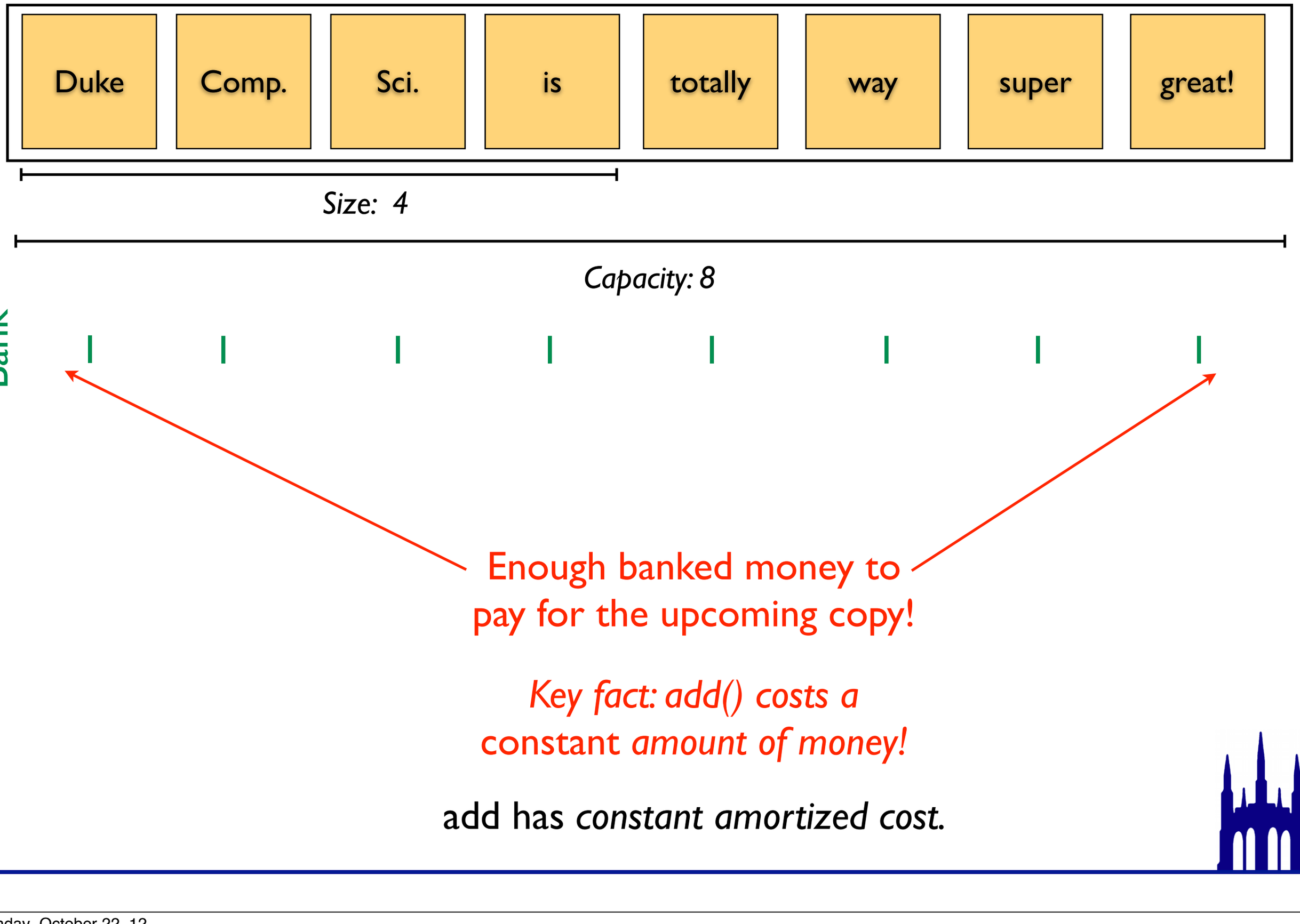


Enough banked money to pay for the upcoming copy!

How expensive is this strategy?



Key fact: add() costs a
constant *amount of money!*

add has *constant amortized cost.*

Amortization facts

Amortized analysis deals with the cost of *n* operations, not the cost of one operation.

"N calls to add cost *O(n)* total."

"One call to add might be *O(n)*, too."

Almost always good enough. So-called "realtime" applications are the exception.

The bank isn't part of the data structure (no data is stored). It's just an analytical tool.

A stack in two queues!

A stack in two queues!

Snarf Oct22InClass

http://goo.gl/4o5oN