

CPS 590.5 Computer Security

Lecture 9: Introduction to Network Security

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Previous lectures

- Worm
- Fast worm design

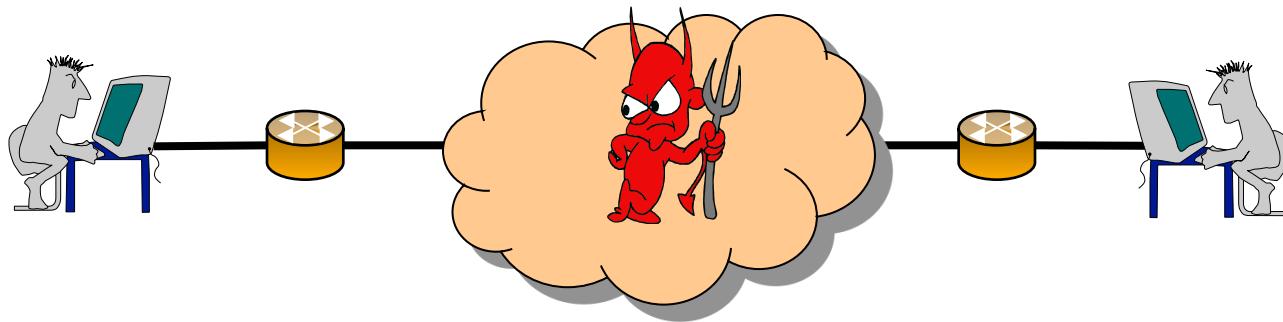
Today

- Network security
 - Cryptography building blocks
 - Existing secure protocols/applications
- Looking forward
 - What's missing

Overview

- Why studying network security?
- Basic cryptography building blocks
- Security protocols
- Non-cryptography based security: firewalls

The Internet is insecure



- Attackers may eavesdrop, modify, or drop your packets!

Network security

- Confidentiality:
 - Do you want to send your credit card #, login password over the Internet in plaintext?
- Integrity
 - Data integrity: Imagine an Amazon transaction. Do you want your payment to be modified from \$10.0 to \$100?
 - Replay attack: You do not want the same transaction confirmation to be sent multiple times!
 - Timeliness: delay a stock purchase
- Authenticity
 - Entity authentication: who are you talking to? Phishing attack
 - Message authentication: who sent this message?
- Availability
 - Denial of service attacks
- Non-repudiation
 - You've clicked the confirmation button!

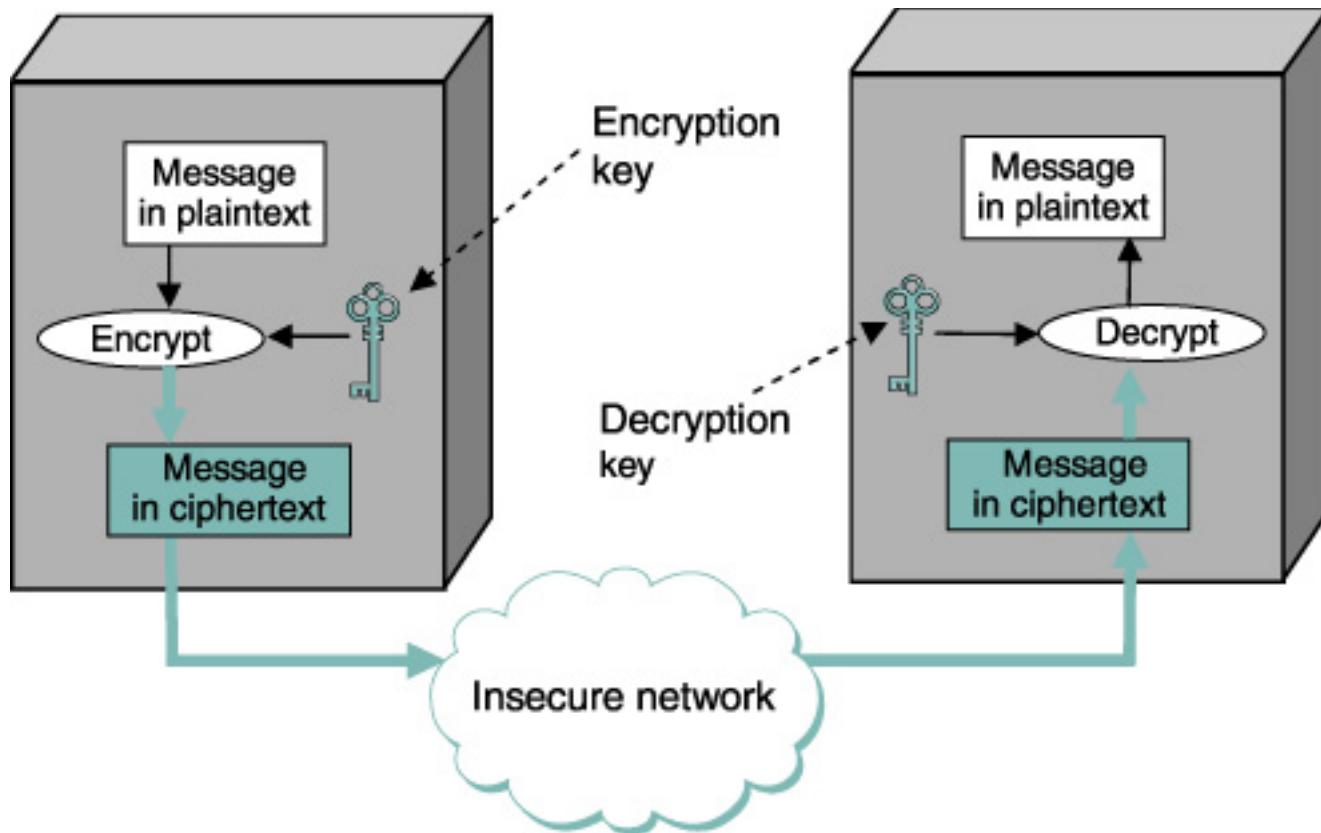
How to address those problems

- Cryptography building blocks
 - Confidentiality
 - Encryption
 - Authenticity
 - Public key signatures
 - Authentication protocols (later)
- Non-cryptographic approach
 - Firewalls

Cryptographic tools

- Cryptographic algorithms
 - Ciphers and Cryptographic hashes
 - Not a solution in themselves, but building blocks from which a solution can be built
- Key distribution
- Protocols built on cryptographic algorithms
 - System builders need to get familiar with the tools

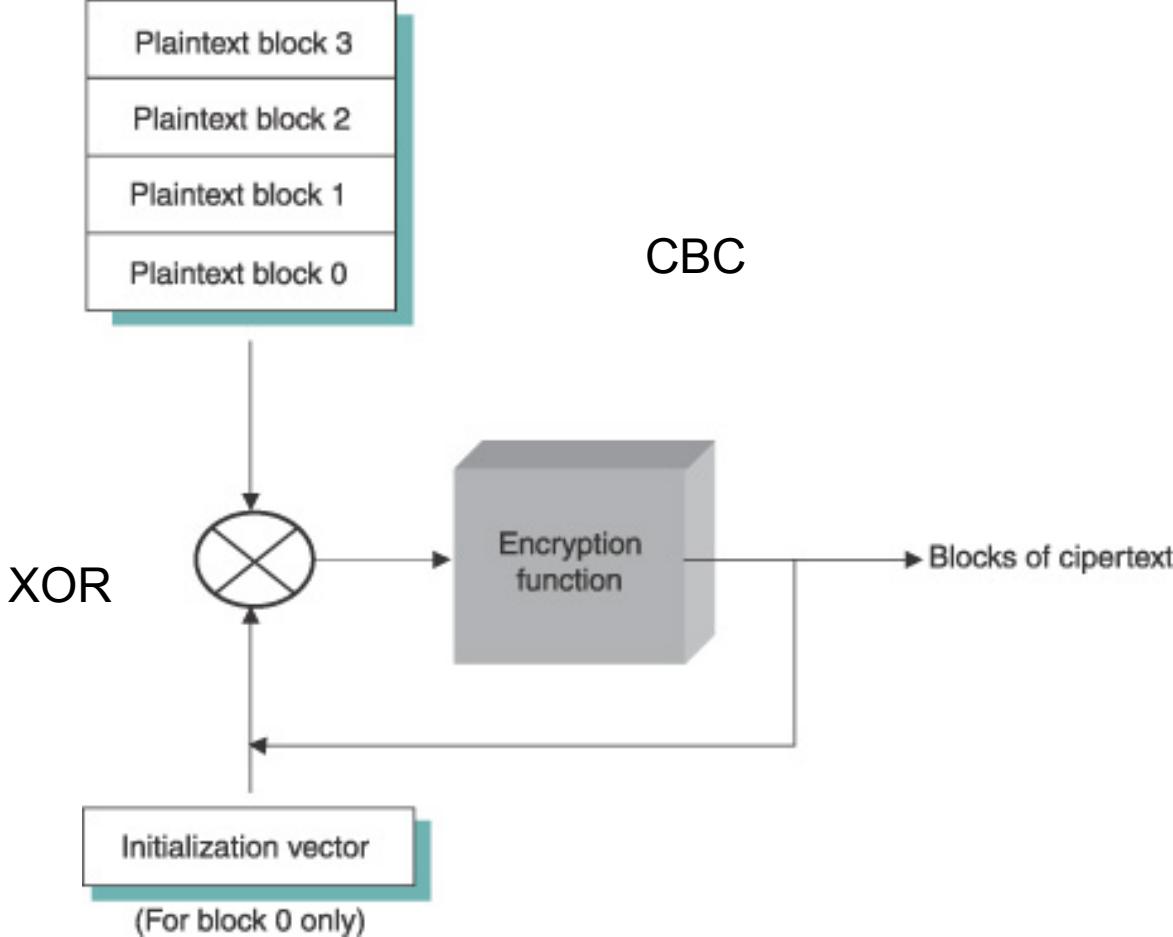
Principles of Ciphers



- $\text{Encrypt}_{\text{key}}(\text{plaintext}) \rightarrow \text{ciphertext}$
- Ciphertext is unintelligible
- $\text{Decrypt}_{\text{key}}(\text{ciphertext}) \rightarrow \text{plaintext}$
- The transformation is called a cipher

Security of a Cipher

- Encrypt() and Decrypt() are public knowledge
- Only key is secret
- Designing a cipher is like a black art
- No news is good news
- Cryptanalysis (code breaking)
 - Known plaintext
 - Know the plaintext and its encrypted version and make use of them to guess other part of secrete information such as secrete keys
 - Chosen plaintext analysis
 - An attacker can get arbitrary plaintext encrypted
 - Some plaintext has known vulnerabilities

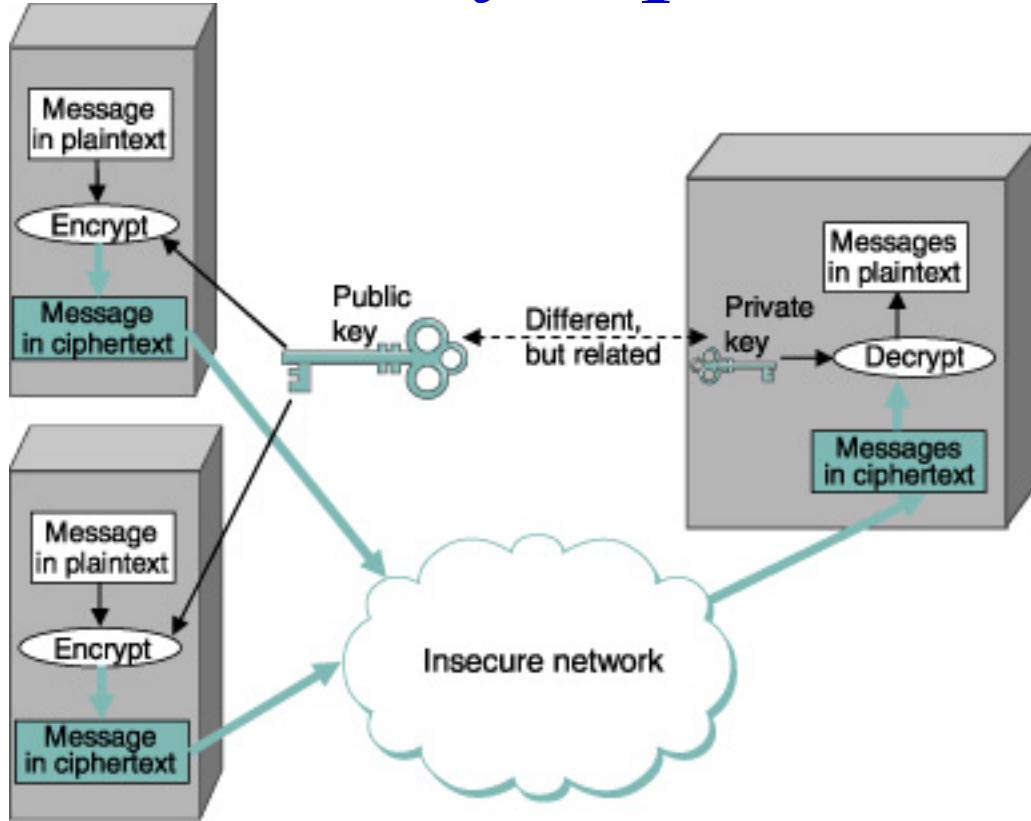


- Input is a fixed size block of text, eg, 64-128 bits
- Modes of operation
 - Electronic codebook (ECB) mode: each block is encrypted independently
 - The same block value will always result in the same cipher text block
 - Cipher block chaining
 - Each plaintext block is XORed with the previous block's ciphertext before being encrypted
 - Context dependent

Standard symmetric-key ciphers

- National Institute of Standards and Technology (NIST) issued ciphers
- Data encryption standard (DES)
 - 56-bit key
 - 64-bit block size
 - Insecure against brute-force attacks
- Triple DES (3DES)
 - First encrypt using DES-key1, decrypt using DES-key2, and encrypt using DES-key3
 - Backward compatible: can be decrypted by DES
- Advanced encryption standard (AES)
 - Originally named Rijndael
 - 128, 192, 256-bits

Public-key ciphers

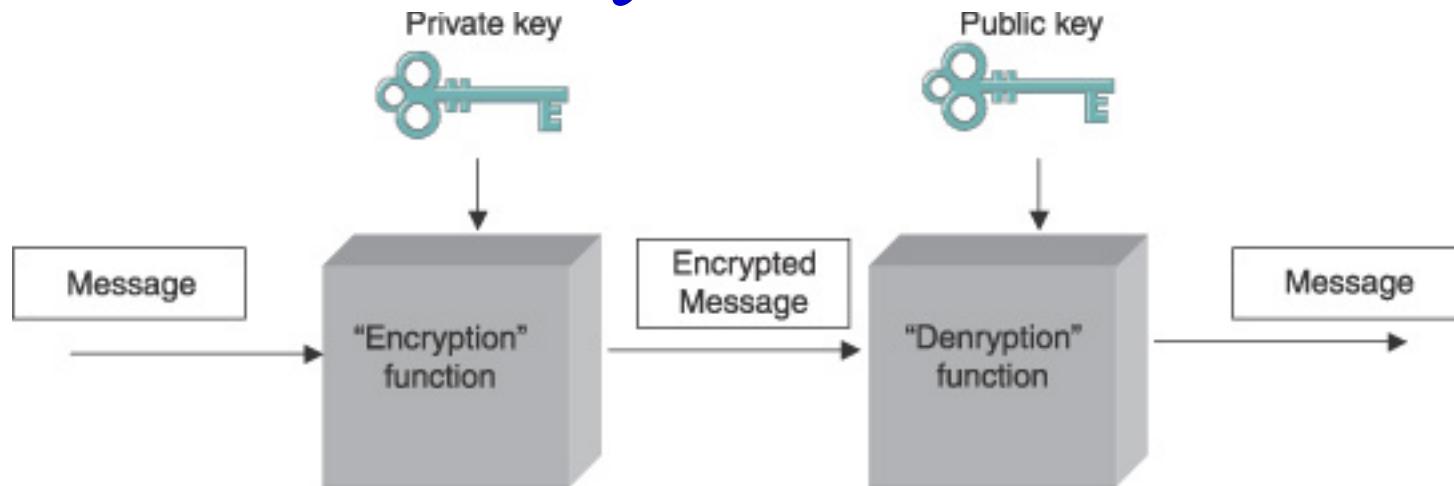


- RSA
 - Difficult to factor large numbers
 - Key length ≥ 1024 bits
- ElGamal
 - Discrete logarithm is hard
 - Key length ≥ 1024 bits
- Public-key ciphers are orders of magnitude slower than symmetric cipher

Cryptography building blocks

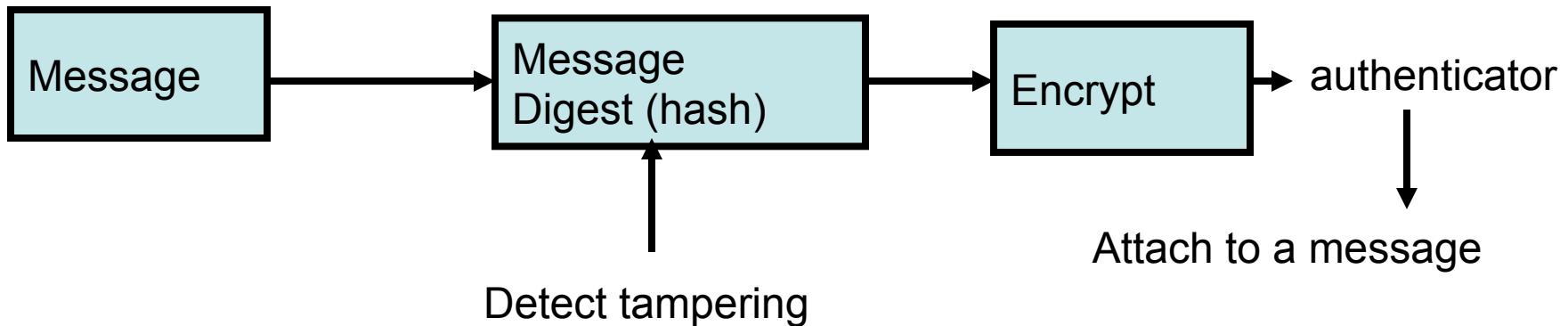
- Confidentiality
 - Encryption
- Authenticity
 - Public key signatures
 - Authentication protocols (later)

Public key authentication



- Everyone can validate who sends the message
- Not good enough
 - No data integrity: modifying a ciphertext message may produce a readable output
 - “I owe you \$10” → “I owe you \$100000”

Authenticators



- **Problem:** Encryption alone does not provide data integrity
 - Modifying a cipher may still allow decrypting to a valid plaintext
- An authenticator is a value, to be included in a transmitted message that can be used to verify simultaneously the authenticity and the data integrity of a message
 - Why are these two properties combined?
- 1. Message digest + encryption
 - Modifying the message cannot produce the correct authenticator

Authenticator methods

- Asymmetric cryptography
 - Digital signatures
- Symmetric cryptography
 - Message authentication code (MAC)
 - Another MAC!

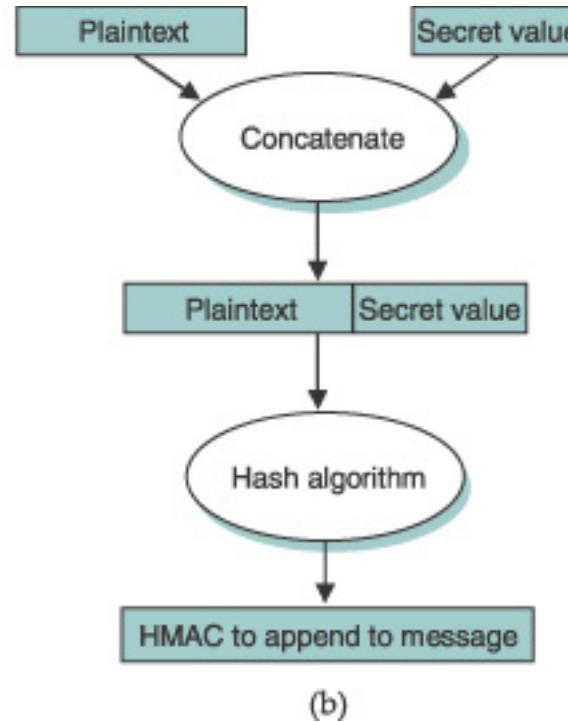
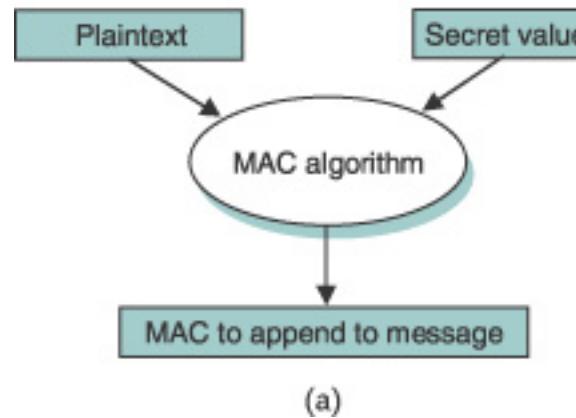
Hash functions

- A secure one-way function $f(x)$
 - Knowing $f(x)$ gives little knowledge about x
- Collision attacks
 - Attacks finding any collision
- Preimage attacks
 - A 2nd message that collides with a given first message
- Common ones: MD5, SHA-1
 - SHA-1 expires by 2010
 - SHA-2

Digital signatures

- A digest encrypted using the private key of a public-key algorithm
- Common digital signatures
 - Digital signature standard (DSS)
 - May use any one of three public-key ciphers
 - RSA, ElGamal, Elliptic Curve Digital Signature Algorithm

Authenticators – Message Authentication Code



Hashed message authentication code

- Instead of encrypting a hash, it uses a hash-like function that takes a secret value (known only to the sender and the receiver) as a parameter.
- How does two ends obtain the key?
- Security of HMAC: what if hash's not one-way?

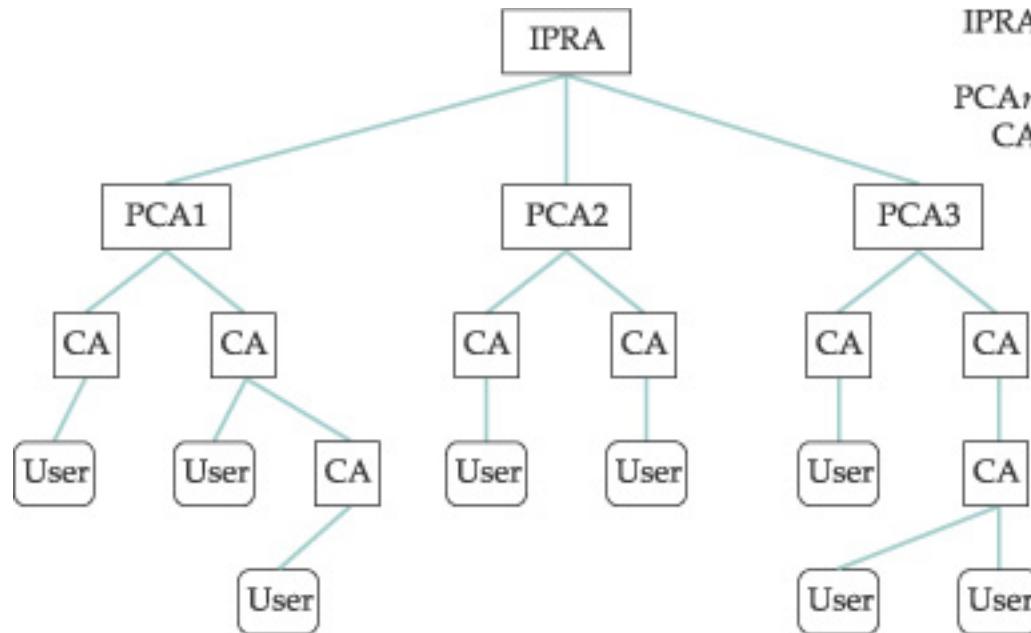
Key distribution

- Two problems:
 - How do participants know which entity has which public key?
 - A complete scheme for certifying bindings between public keys and identities – what keys belongs to who – is called a public key infrastructure (PKI)
 - Comments: not easy to scale
 - People don't use it that much
 - How do each end know the symmetric shared key?

Distributing public keys

- A public-key certificate is a digitally signed statement that binds the identity of the entity to a public key
- If A trusts B, and knows B's public key, then A can learn C's public key if B issues a public key certification of C
- X.509 certificate
 - The ID of the entity
 - The public key of the entity
 - The identity of the signer
 - The digital signature
 - A digital signature algorithm
 - Optional: expiration time

Certification authorities



IPRA = Internet Policy
Registration Authority (root)
PCAn = Policy certification authority
CA = Certification authority

- A CA is an entity claimed to be trustworthy to verify identities and issuing public key certificates
 - Verisign
- CAs can be organized into a tree
- Trust is binary: yes or no
 - Everyone trusts the root

Multiple CAs

- In the real world, there is no single rooted trust
- Multiple CAs whose public keys are trusted by different people
- Self-certifying certificates
 - Signer is self
 - Accepted by TLS

Web of Trust

- Pretty Good Privacy:
 - No single hierarchical
 - Establishing trust is a personal matter and gives users the raw material to make their own decisions
- IETF's PGP signing session:
 - Collect public keys from others whose identity one knows
 - Provide his public key to others
 - Get his public key signed by others
 - Sign the public key of others
 - Collect the certificate from other individuals whom he trusts enough to sign keys
- Trust is a matter of degree
 - A public-key certificate includes a confidence level
 - Trust dependent on the number of certificates of a key, and the confidence level of each certificate

Certificate Revocation

- Certificate revocation list
 - Periodically updated and publicly available
 - Digitally signed
 - Lists may be large
- Online certificate status protocol
 - Query the status of a certificate

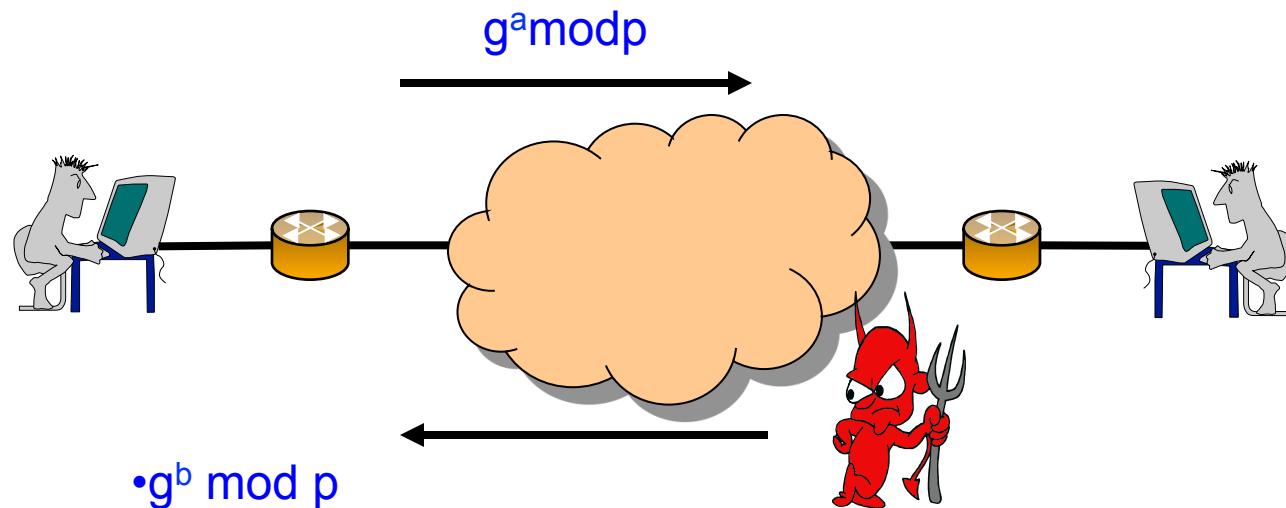
Key distribution

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Symmetric key distributions

- If there are N entities, $N(N-1)/2$ keys
- Key distribution center (KDC)
 - A trusted entity
 - Each user maintains a key with the KDC
 - KDC generates a session key when a user wants to communicate with another destination
 - Kerberos is a widely used key-distribution system

Diffie-Hellman key agreement

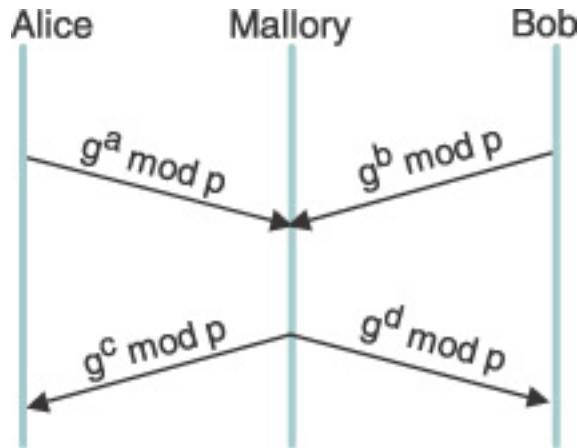


- Long considered as the invention of public key cryptography
- Establishes a session key without using any predistributed keys
- Discrete log is hard

Diffie-Hellman Key Agreement

- Two parameters: g , and p
 - p : a prime; g : a primitive root of p s.t. for every number of n from 1 through $p-1$ there must be some value k such that $n=g^k \bmod p$
 - $1=2^0 \bmod 5, 2=2^1 \bmod 5, 3=2^3 \bmod 5, 4=2^2 \bmod 5$
- Alice picks a private value a , and sends $g^a \bmod p$
- Bob picks b , and sends $g^b \bmod p$
- $g^{ab} \bmod p = g^{ba} \bmod p$
- Discrete log is hard
 - Attackers cannot guess a , or b , even when they see $g^a \bmod p$ or $g^b \bmod p$

Man in the middle attack



- Fixed DH: Alice and Bob has fixed a, and b values
- $g^a \text{ mod } p$ is certified

Overview

- Why studying network security?
 - The topic itself is worth another class
- Basic cryptography building blocks
- Security protocols
 - Case studies of using cryptography to build secure systems
- Non-cryptography based security: firewalls

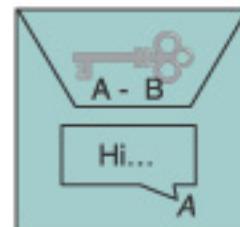
Secure systems

- PGP for email security
 - Works well for email
 - Could be replayed, but a user can detect
 - No need of prior message exchange
 - Confidentiality
 - Does not prove Alice is talking to Bob

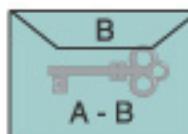
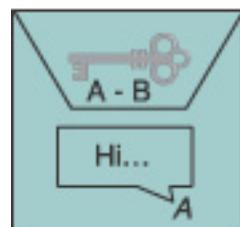
Hi... = The plaintext message



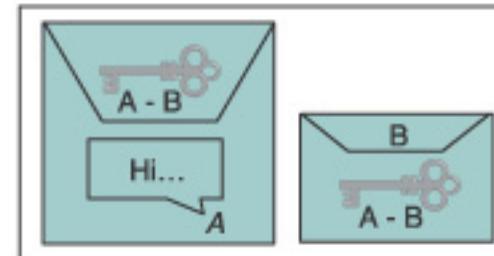
1) Digitally sign using Alice's private key



2) Encrypt using a newly generated one-time session key



3) Encrypt the session key using Bob's public key, and append that



4) Use base64 encoding to obtain an ASCII-compatible representation

base64

Secure Shell (ssh)

- Both the SSH protocol and applications that use it
- Widely used
- Replacing telnet, rsh, rcp
 - No need to send plaintext passwords to authenticate
 - Prior to it, passwords are sent in plaintext!

Components of SSH2

- SSH 2 consists of three protocols
 - **SSH-TRANS**: a transport layer protocol
 - Provides an encrypted channel between the client and server machines
 - A secure channel is established before a client authenticates
 - A client authenticates the server to establish the channel
 - **SSH-AUTH**, an authentication protocol
 - Client authenticates to the server
 - You may type your password! SSH-TRANS takes care of encryption
 - **SSH-CONN**, a connection protocol
 - Used for port forwarding

SSH's server key distribution

- A server tells the client its public key at connection time
 - Attackers are not always present
 - Vulnerability window is small
- The first time a key is sent, ssh asks the user
- If accepts, ssh remembers the key and compares the stored key with an offered key in a subsequent connection
 - Prompts the user if changed
 - Otherwise accept

Establishing a secure channel

- After a client authenticates the server
- Establish a session key
- All subsequent messages are encrypted using the session key

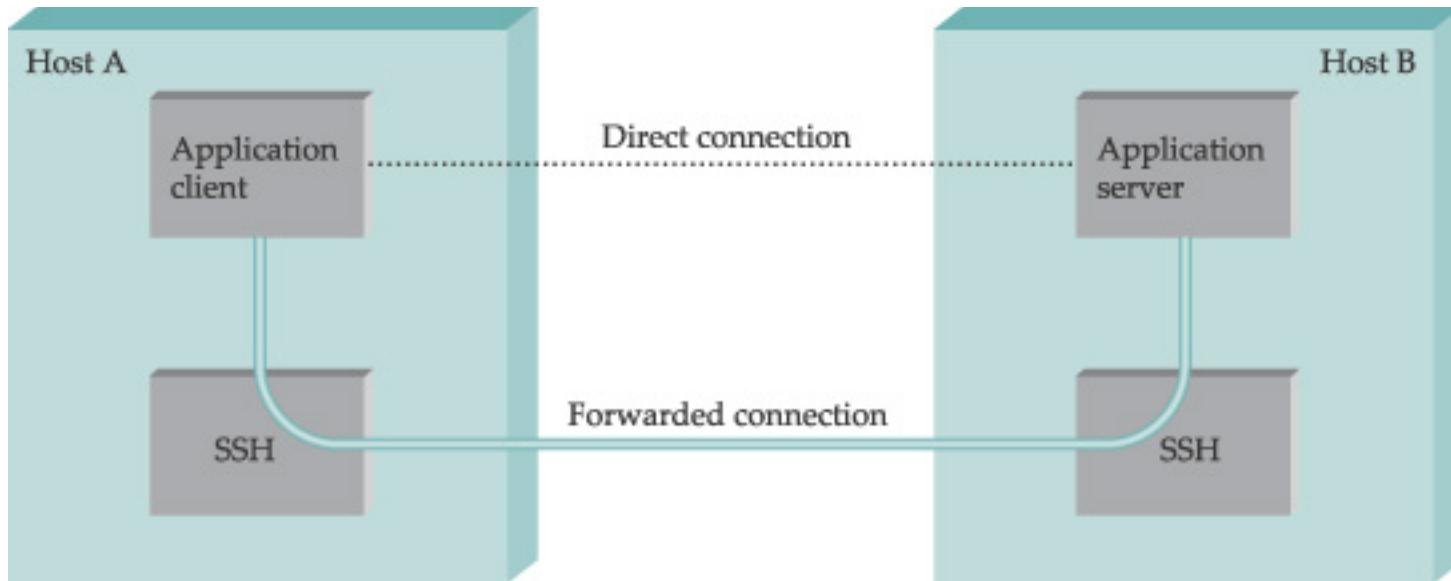
SSH's client authentication

- Password
 - A secure channel is already established!
- Public key encryption
 - Places your public key in `~/.ssh/authorized_keys`
- Host authentication
 - A user claiming to be so-and-so from a certain set of trusted hosts is automatically believed to be the same user on the server
 - The client host authenticates itself to the server
 - SSH-TRANS only authenticates server by default
 - User claims to be so-and-so from a set of trusted hosts is believed to be the same user on the server

SSH login w/o typing in your passwords

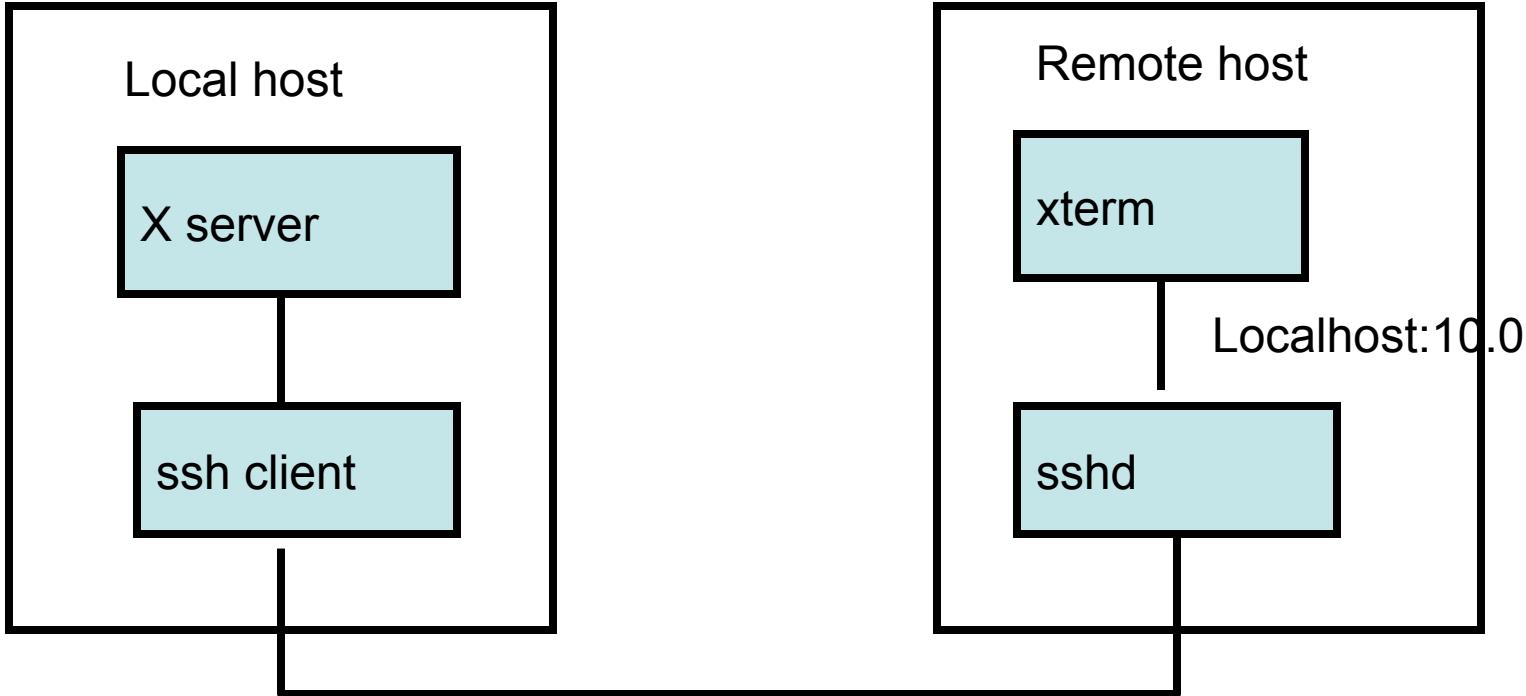
- Use ssh-keygen to generate a public/private key pair
 - ssh-keygen -t dsa
- Append .id_dsa.pub to .ssh/authorized_keys on the server
 - scp ~/.ssh/id_dsa.pub linux1.cs.duke.edu:~/
 - ssh linux1.cs.duke.edu “cat ~/id_dsa.pub >> ~/.ssh/authorized_keys”

SSH port forwarding



- SSH can be used to establish a secure channel between two hosts using the SSH-CONN protocol
- Secure legacy applications

Example: X11 forwarding

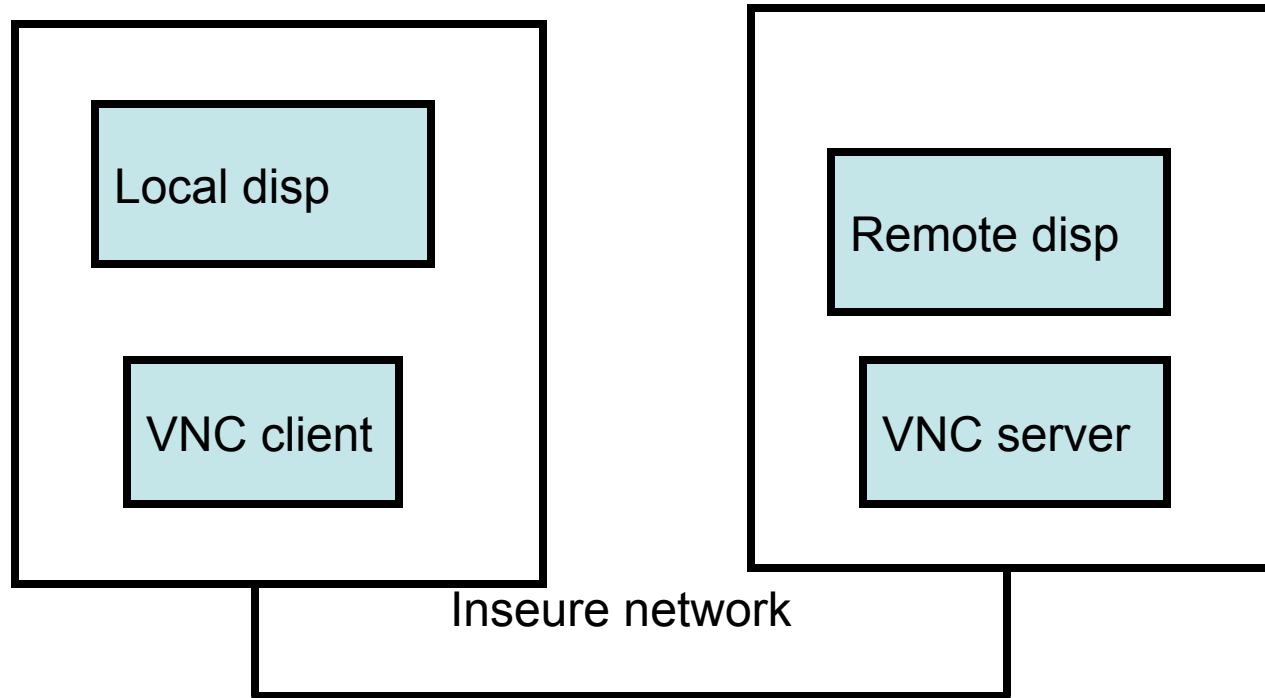


- `ssh -X remote-host`
- sshd at the remote host creates a TCP listening socket (6010), and X authentication information, and set your display variable to `localhost:10.0`
- When you type `xterm`, it speaks the X11 protocol with the sshd faked X server port (10.0), and sshd forwards it back to the ssh client at local host.
- The ssh client forwards back to the X server running on your local host
- You see the `xterm` displays on your local host, and all commands you type in the `xterm` is encrypted!

SSH port forwarding

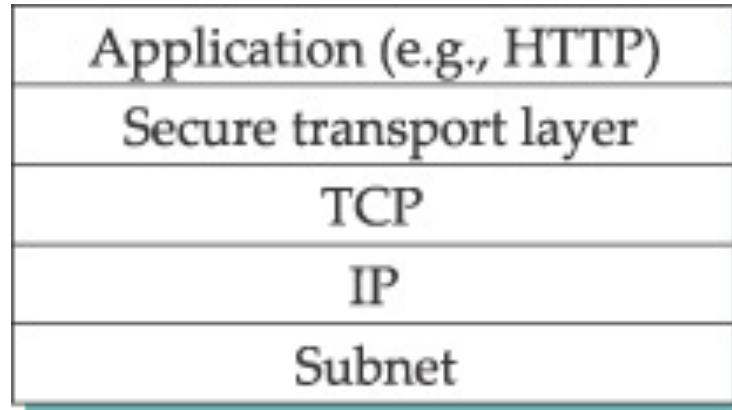
- Some legacy applications do not have security mechanism built-in: pop3
- How can you read your email without sending your password in cleartext?
 - `ssh -L 9999:localhost:110 mail.cs.duke.edu`
 - Run your pop3 mail client, and make it use localhost:9999
 - All commands will be sent via an encrypted connection
- Pop3 \longleftrightarrow localhost:9999 \longleftrightarrow ssh client \longleftrightarrow mail.cs.duke.edu:sshd \longleftrightarrow mail.cs.duke.edu:110

The VNC example



- A real world application
 - The free VNC servers do not have encryption
 - Figure out how to do it yourself!
 - Vncviewer sends your password in plaintext to a vncserver
 - Unless you purchase the non-free version
 - How can we establish a secure tunnel between the vncclient and server?

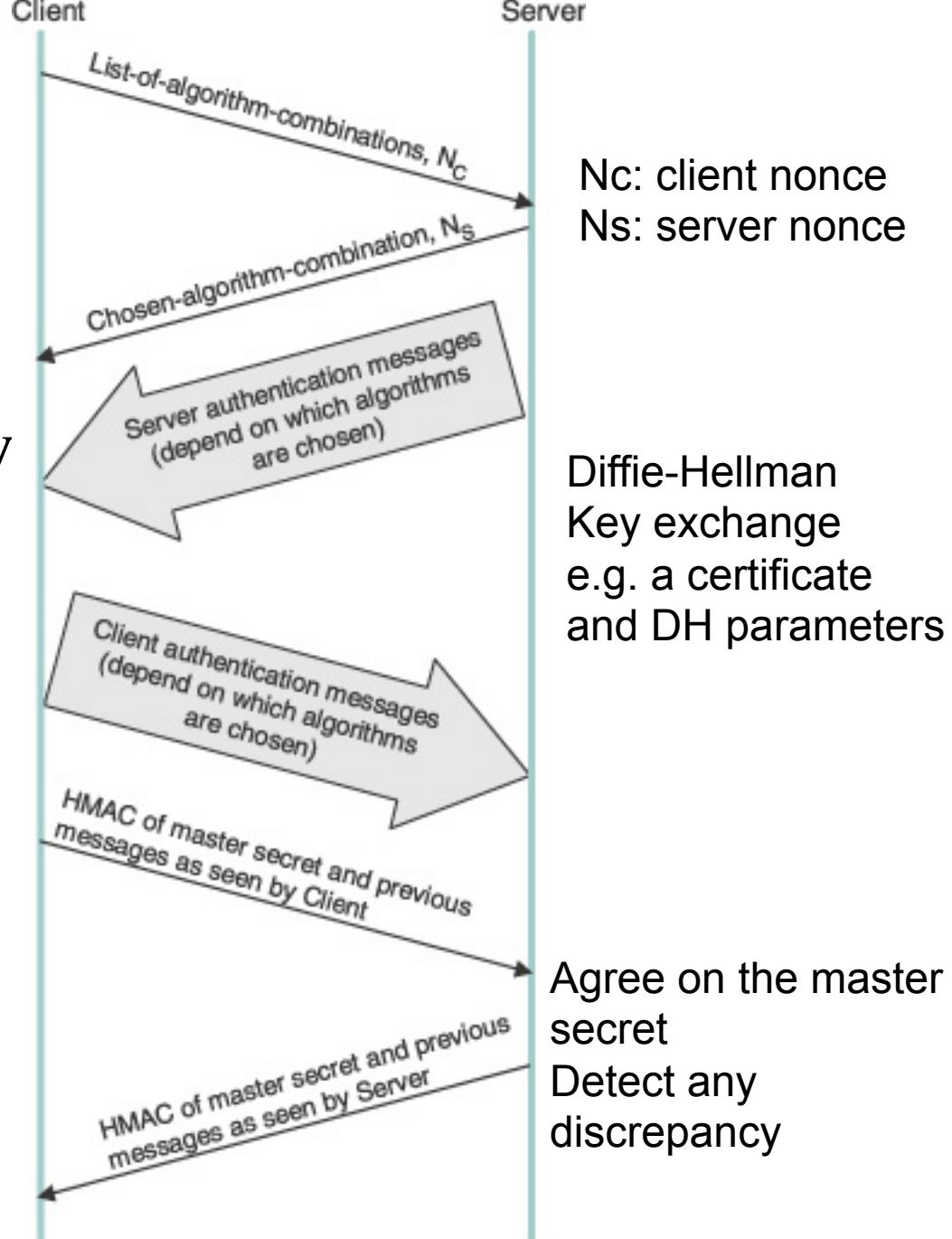
SSL/TLS



- Transport layer security (TLS) is based on Secure Socket Layer (SSL)
- https: port 443
- A handshake protocol for negotiating parameters, and a secret session key
 - Each direction has a key
- A record protocol to transmit messages

The handshake protocol

- Negotiate encryption hash, symmetric key cipher, and session key establishment protocol
 - Mutual authentication
 - Or authenticating one participant only
 - Or no authentication



The record protocol

- Fragmented or coalesced into blocks of a convenient size
- Optionally compressed
- Integrity-protected each record
- Encrypted using <https://gmail.com> was not the default for a long time
- Passed to the transport layer
- Each record message has its own sequence number to prevent replay attacks.
 - TCP's sequence number is not sufficient!
 - An active adversary can modify TCP sequence numbers in packets!

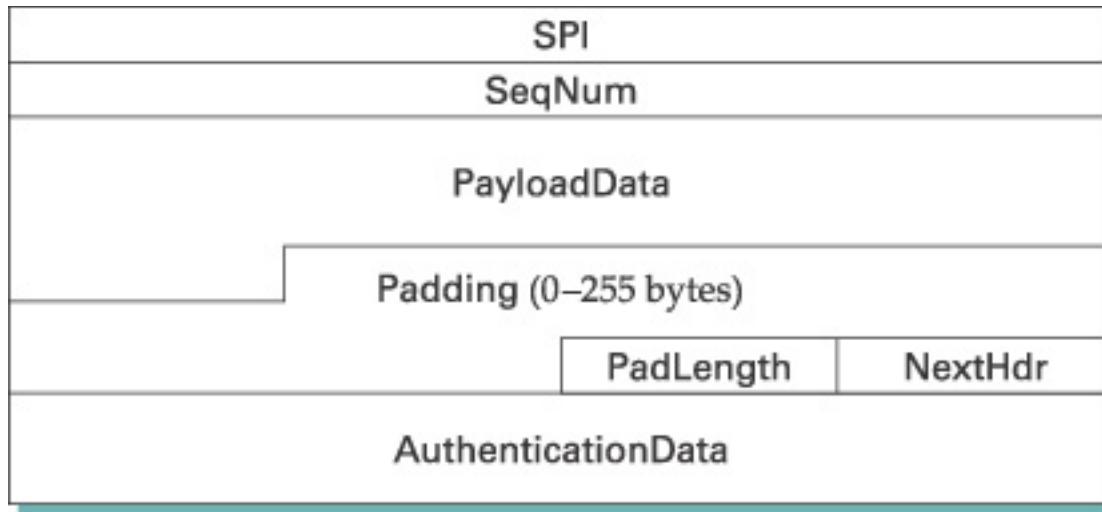
IPSec

- A framework specifies how to secure all IP traffic between two machines
- Two parts
 - Security services: Authentication header (AH): rarely used, Encapsulating Security Payload (ESP)
 - Key management: Internet security association and key management protocol (ISAKMP).
 - Defines message format, not the detailed key generation algos.

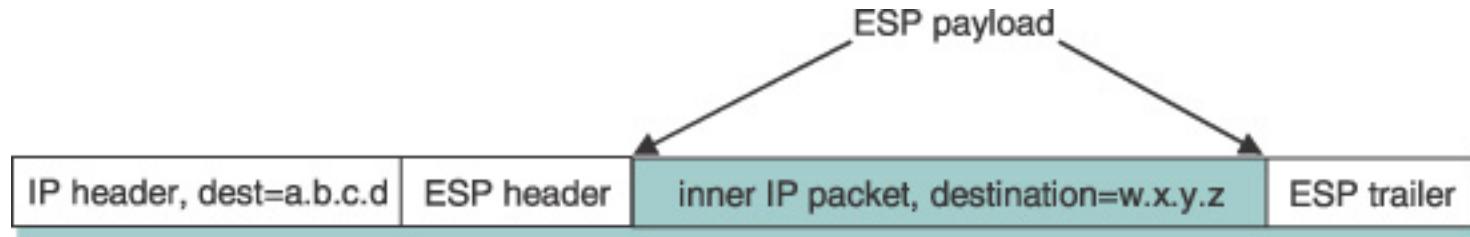
Security Association

- A security association (SA) is created for each direction
 - IP is connectionless, but IPSec is not
 - An SA includes connection state such as keys, and sequence numbers
 - An SA is identified by a security parameter SPI (a multiplexing key)
 - SPI and destination address identifies an SA
 - SAs are established, negotiated, modified and deleted using ISAKMP
 - Internet Key Exchange (IKE) is one key exchange protocol

The ESP header

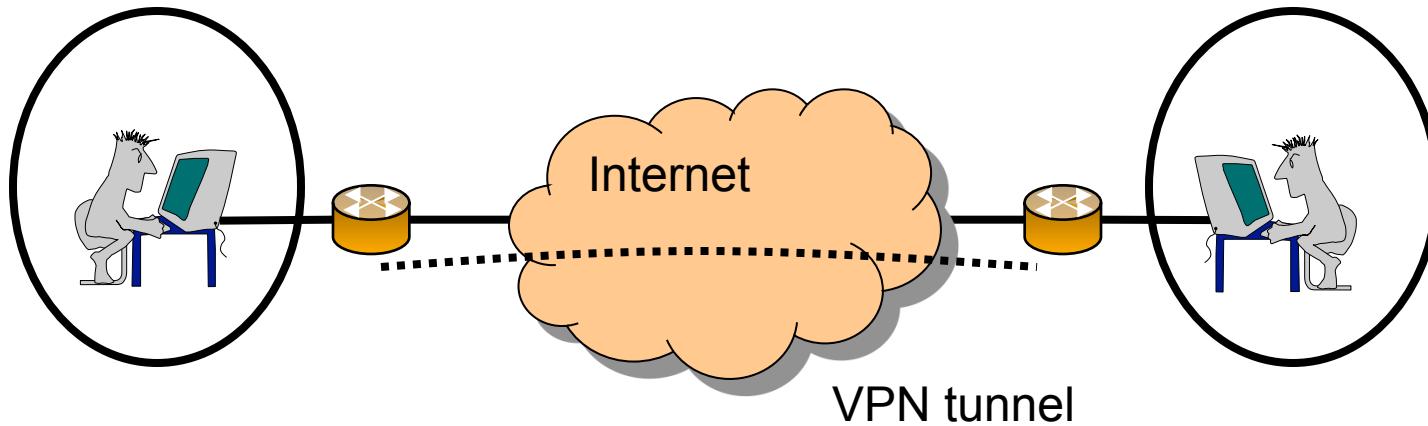


- Padding is necessary due to cipher requirement
- Payload's encrypted
- Two modes: tunnel, or transport



Company site 1

Company site 2



- Tunnel mode: bump-in-the-wire
 - Useful in creating VPNs
 - Payload is an IP packet

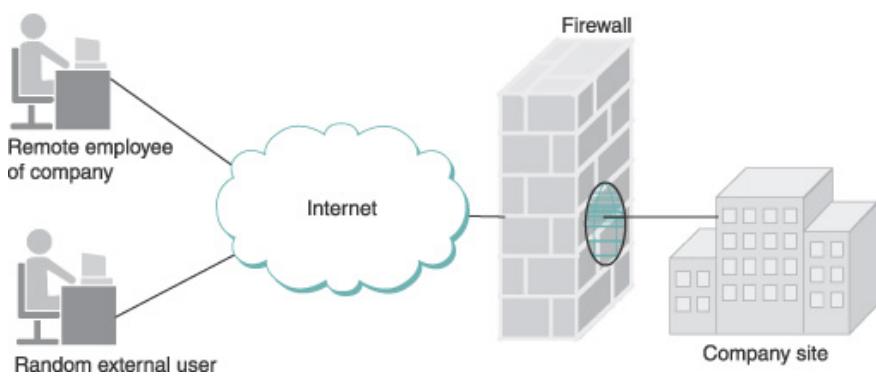
- Transport model
 - Upper layer payload is encrypted
 - UDP, TCP

Discussion

- Differences between IPsec and TLS
 - Pros and cons

Firewalls

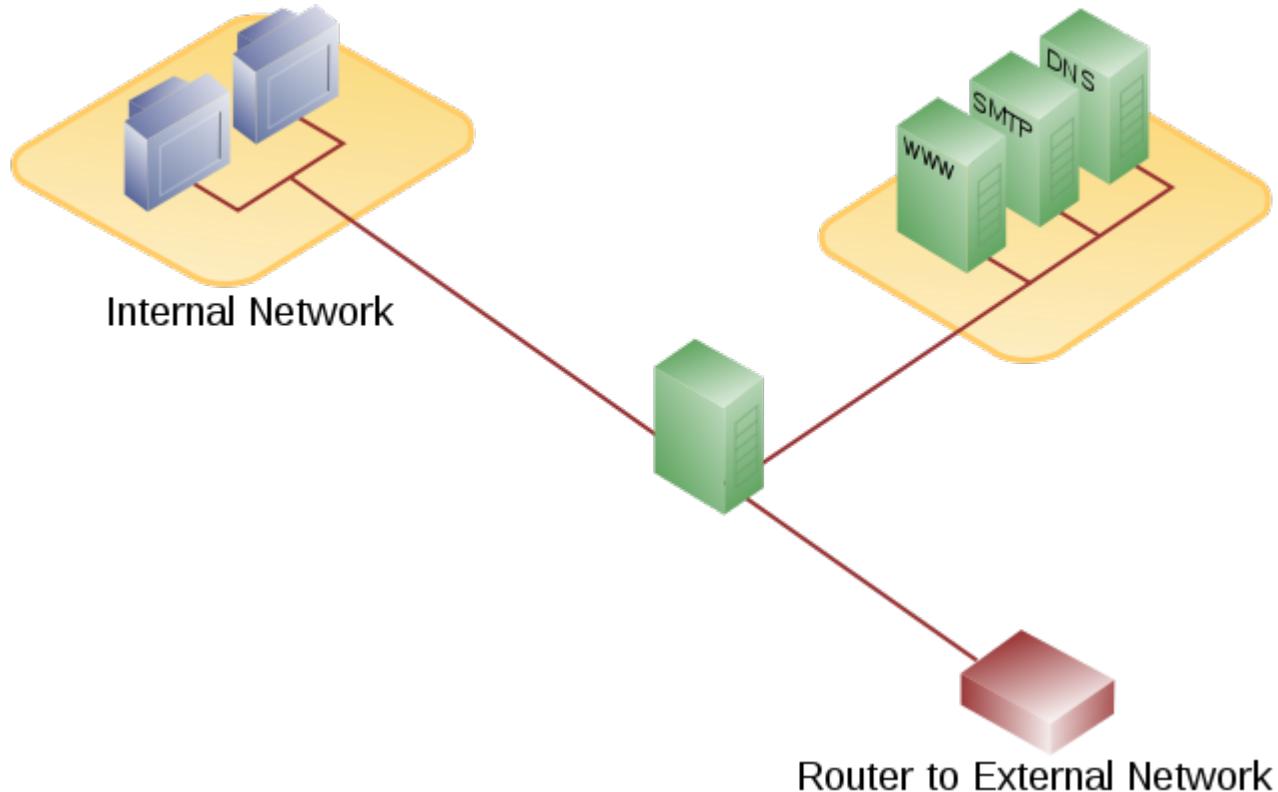
- Firewalls create zones of trust
 - The internal network
 - Demilitarized zone (DMZ)
 - DNS, email servers
 - Hosts in DMZ accessible by anyone
 - Cannot access internal hosts
 - DMZ can be periodically restored
 - The rest of the Internet
- Widely used in practice
 - Unilaterally deployed



Firewall configurations

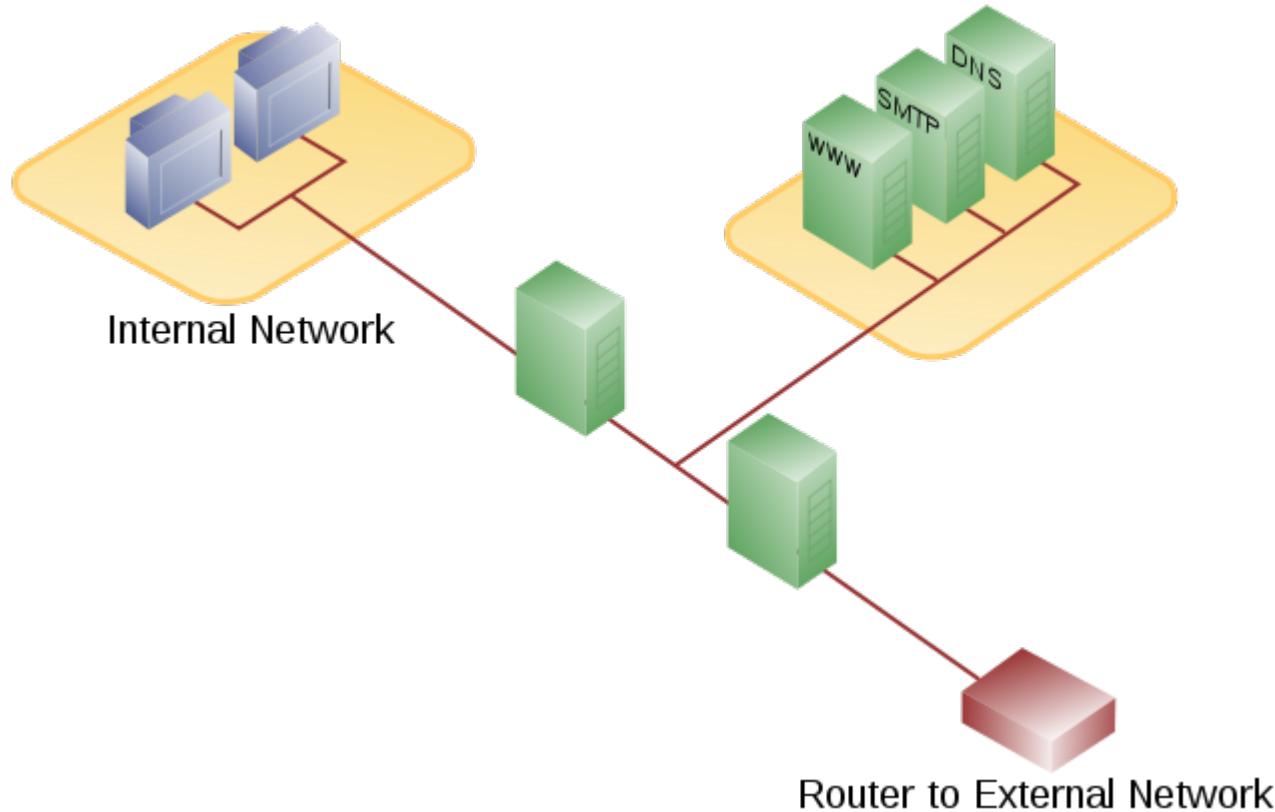
- Access lists: similar to tcpdump's filter lists
- Allows outside connection to Duke CS's main mail server one.cs.duke.edu:
 - (*, *, 152.3.140.161, 25, allow)
- Disallow to internal mail server
 - (152.3/16, *, 152.3.140.1, 25, allow)
 - (*, *, 152.3.140.1, 25, deny)

Using Firewalls to Create DMZ



- A firewall with three interfaces

Using Firewalls to Create DMZ



- Dual firewalls

Conclusion

- Network security in practice
 - Building blocks
 - Symmetric key cryptography
 - Public key cryptography
 - Protocols and applications
 - SSH
 - TLS
 - IPSec
 - Firewalls