

-  $(2^k + 1)$ -th add operation  
 running time  $\sum^{k+1}$  "heavy" / "expensive"

(allocating an array of size  $2^{k+1}$   
 copy first  $2^k$  elements  
 add the  $(2^k + 1)$ th element)

- all other add operation "light"  
 running time 1 (change length, put the element into an empty slot)

$$- \begin{cases} T_{2^k+1} = 2^{k+1} & \text{for all } k \geq 0 \\ T_i = 1 & \text{for all } i \neq 2^k+1 \end{cases}$$

- Potential argument

$$\begin{aligned} A_i &= T_i - \Phi(x_i) + \Phi(x_{i+1}) \\ &= T_i - \underbrace{(\Phi(x_i) - \Phi(x_{i+1}))}_{\text{potential difference}} \end{aligned}$$

$$\begin{cases} A_1 = T_1 - \Phi(x_1) + \Phi(x_2) \\ A_2 = T_2 - \Phi(x_2) + \Phi(x_3) \\ \dots \\ A_n = T_n - \Phi(x_n) + \Phi(x_{n+1}) \end{cases}$$

$$\underbrace{\sum_{i=1}^n A_i}_{\text{total amortized}} = \underbrace{\sum_{i=1}^n T_i}_{\text{total runtime}} - \Phi(x_1) + \Phi(x_{n+1})$$

- design the potential function

- idea: first make sure  $\Phi(x_i) - \Phi(x_{i+1})$  is large  
 if  $i = 2^k + 1$  is a heavy operation. ( $T_i = 2^{k+1}$ )

$\Phi$  is a function of # elements and capacity

$\Phi$  is a function of  $\frac{\# \text{ elements}}{l}$  and  $\frac{\text{capacity}}{c}$

$$X_i: \quad \begin{array}{cc} l & c \\ 2^k & 2^k \end{array}$$

before  $2^{k+1}$  add operation, data structure has  $2^k$  elements.

$$X_{i+1} \quad \begin{array}{cc} 2^{k+1} & 2^{k+1} \end{array}$$

difference  $\rightarrow \quad \begin{array}{cc} 1 & 2^k \end{array}$

$\Rightarrow$  if capacity  $\uparrow$  by  $2^k$ , potential  $\Phi \downarrow$  by  $2^{k+1}$

$$\Phi = \textcircled{?} - 2 \cdot c$$

- step 2:  $\Phi \geq 0$

observation:  $l \cdot 2 \geq c$  (except for the initial state)

$$1 + 2l \geq c \quad (\text{always true})$$

can use  $\Phi = 2 + 4l - 2c$

and we know  $\Phi \geq 0$

- compute amortized running time

① heavy operation  $i = 2^k + 1 \quad T_i = 2^{k+1}$

$$\Phi(X_i) = 2 + 4(2^k) - 2(2^k)$$

$$\Phi(X_{i+1}) = 2 + 4(2^k + 1) - 2(2^{k+1})$$

$$\Phi(X_i) - \Phi(X_{i+1}) = 2^{k+1} - 4$$

$$A_i = T_i - (\Phi(X_i) - \Phi(X_{i+1})) = 2^{k+1} - (2^{k+1} - 4) = 4$$

② light operation  $i \neq 2^{k+1} \quad T_i = 1$

$$X_i: \quad \begin{array}{ccc} l & c & \Phi \\ i-1 & c & 2 + 4(i-1) - 2c \end{array}$$

$$X_{i+1}: \quad \begin{array}{ccc} i & c & 2 + 4i - 2c \end{array}$$

$$\Phi(X_i) - \Phi(X_{i+1}) = -4$$

$$A_i = T_i - (\Phi(X_i) - \Phi(X_{i+1})) = 1 - (-4) = 5$$

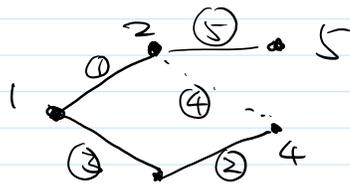
amortized cost =  $\sqrt{n} = O(1)$

## - Union-Find

- recall: Kruskal

- for each edge: want to know whether adding the edge creates a cycle.

- Sets  $\iff$  connected components



add an edge  $\iff$  union on the two connected components

edge creates a cycle  $\iff$  Find(u) = Find(v)  
(u,v) u, v in the same set.

- edge ①: Find(1) = 1    Find(2) = 2    {1} {2} {3} {4} {5}  
union(1, 2)    {1, 2} {3} {4} {5}

- edge ②: Find(3) = 3    Find(4) = 4  
union(3, 4)    {1, 2} {3, 4} {5}

- edge ③ (1, 3): Find(1) = 1    Find(3) = 3  
union(1, 3)    {1, 2, 3, 4} {5}

- edge ④ (2, 4): Find(2) = 1 = Find(4) = 1  
do not add the edge