

- Example 1 LIS (Longest Increasing Subsequence)

① find subproblems and transition function

$$\{4, 2, 3, 5, 1, 7, 10, \underline{8}\}$$

focus on last element of the array

is 8 in the LIS?

compare { ① Length of LIS if 8 is not in the sequence
② _____ if 8 is in the sequence.

- to solve ① call LIS on $\{4, 2, 3, \dots, 10\}$

- to solve ②, not as easy

- in this example $LIS\{4, 2, 3, \dots, 10\} = \{2, 3, 5, 7, 10\}$
cannot add 8 to this sequence

- $\{2, 10, 3, 5, 8\}$

$LIS\{2, 10, 3, 5\} = \{2, 3, 5\}$

can/should add 8 to this sequence.

- ideas: (a) have a subproblem $a[i, j]$

s.t. $a[i, j] =$ the LIS of first i elements whose
last element is smaller than j

(b) have a subproblem $a[i]$

$a[i] =$ the LIS of first i elements that ends at
 i -th element

A $4, 2, 3, 5, 1, \underline{7}, 10, \underline{8}$

a $1 \textcircled{1} 2 \textcircled{3} \underline{1} \textcircled{4} 5 \textcircled{5}$

how to compute $a[i]$? i -th element must be in

if $j < i$ $\underline{A[j]} < A[i]$ then we can add

$A[i]$ to the end of an IS ending at $A[j]$

$$a[i] = \max_{\substack{j < i \\ A[j] < A[i]}} \{ 1 \text{ if } A[i] < A[j] \text{ for all } j < i \\ a[j] + 1 \}$$

② figure out the base cases

③ find an appropriate order $i = 1, 2, 3, \dots, n$

for $i = 1 \text{ to } n$

$$a[i] = 1$$

{ for $j = 1 \text{ to } i-1$

if $A[j] < A[i]$ and $a[j]+1 > a[i]$ then

$$a[i] = a[j]+1$$

Example 2 Knapsack

Look at the last item

{ ① max value if last item is in Knapsack
 ② _____ if last item is not in Knapsack.

- to solve ① w_n, v_n already in my Knapsack

should try to maximize value for first $n-1$ items
 using a Knapsack of capacity $W - w_n$.

② should try to maximize value for first $n-1$ items
 using a Knapsack of capacity W .

$a[i, j] = \max$ value from first i items with capacity j .

{ ① $a[n-1, W - w_n]$
 ② $a[n-1, W]$

$a[i, j] = \max \{ a[i-1, j - w_i] + v_i \quad \text{putting } i\text{th element in } j \geq w_i \\ a[i-1, j] \quad \text{not putting } i\text{th element} \}$

$a[i-1, j - w_i] + v_i$ putting i th element in $j \geq w_i$

$a[i-1, j]$ not putting i th element

$$a[0, \text{anything}] = 0 \quad (\text{no items})$$

$$a[i, 0] = 0 \quad (\text{no capacity})$$

orderly $i = 1 \rightarrow n$

$$j = b[i, j] \leq a[i-1, j]$$

if $j \geq w_i$ and $a[i-1, j-w_i] + v_i \geq a[i-1, j]$

$$a[i, j] = a[i-1, j-w_i] + v_i$$

(other orderings can also work)