Section: Turing Machines

Review

Regular Languages

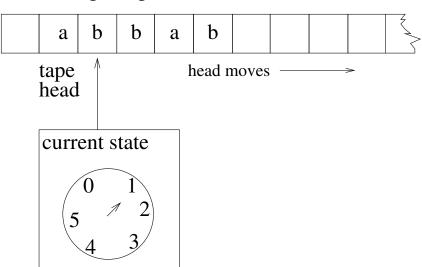
- FA, RG, RE
- recognize

Context Free Languages

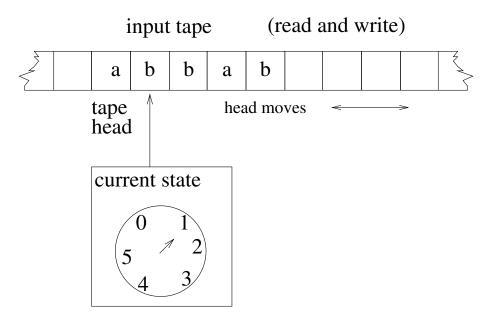
- PDA, CFG
- recognize

DFA:

input tape



Turing Machine:



Turing Machine (TM)

- invented by Alan M. Turing (1936)
- computational model to study algorithms

Definition of TM

- Storage
 - -tape
- actions
 - -write symbol
 - -read symbol
 - -move left (L) or right (R)

• computation

- initial configuration
 - * start state
 - * tape head on leftmost tape square
 - * input string followed by blanks
- processing computation
 - * move tape head left or right
 - * read from and write to tape
- -computation halts
 - * final state

Formal Definition of TM

A TM M is defined by $\mathbf{M} = (\mathbf{Q}, \Sigma, \Gamma, \delta, q_0, \mathbf{B}, \mathbf{F})$ where

- Q is finite set of states
- Σ is input alphabet
- Γ is tape alphabet
- B $\in \Gamma$ is blank
- q_0 is start state
- F is set of final states
- δ is transition function $\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{p}, \mathbf{b}, \mathbf{R})$ means

TM as Language recognizer

Definition: Configuration is denoted by \vdash .

if $\delta(q,a) = (p,b,R)$ then a move is denoted

abaqabba ⊢ ababpbba

Definition: Let M be a TM, $\mathbf{M} = (\mathbf{Q}, \Sigma, \Gamma, \delta, q_0, \mathbf{B}, \mathbf{F})$. $\mathbf{L}(\mathbf{M}) = \{w \in \Sigma^* | q_0 w \stackrel{*}{\vdash} x_1 q_f x_2 \text{ for some } q_f \in \mathbf{F}, x_1, x_2 \in \Gamma^* \}$

TM as language acceptor M is a TM, w is in Σ^* ,

- ullet if $w \in L(M)$ then M halts in final state
- if $w \notin L(M)$ then either
 - -M halts in non-final state
 - -M doesn't halt

TM as a transducer

TM can implement a function:

$$f(w)=w'$$

start with: w

 \uparrow

end with: w'

 \uparrow

Definition: A function with domain D is Turing-computable or computable if there exists TM $\mathbf{M} = (\mathbf{Q}, \Sigma, \Gamma, \delta, q_0, \mathbf{B}, \mathbf{F})$ such that

$$q_0w \overset{*}{\vdash} q_f f(w)$$

 $q_f \in \mathbf{F}$, for all $w \in \mathbf{D}$.

Example

$$\Sigma = \{a, b\}$$

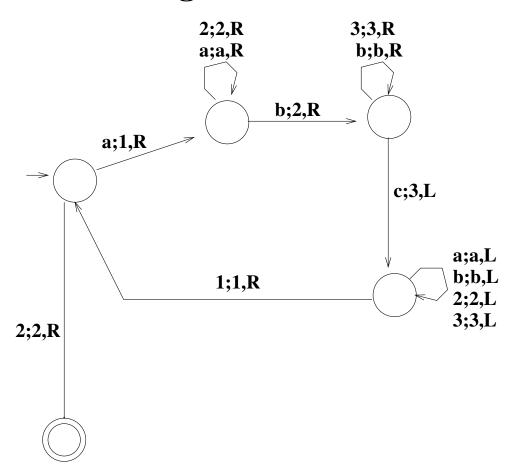
Replace every second 'a' by a 'b' if string is even length.

• Algorithm

Example:

$$\mathbf{L} = \{a^n b^n c^n | n \ge 1\}$$

Is the following TM Correct?



Example:

$$f(x) = 2x$$

x is a unary number

start with: 111

个

end with: 111111

 \uparrow

Is the following TM correct?

