

Survey of General-Purpose Computation on GPU and Introduction to CUDA

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Introduction to GPGPU

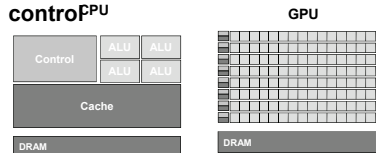
- Powerful and Inexpensive
 - Semiconductor capability, driven by advances in fabrication technology, market
- Flexible and programmable
- Limitations and difficulties

Application

- Physically Based Simulation
- Signal and Image Processing
 - Image segmentation
 - Computer Vision
 - Image Processing
- Geometric Computing

Design philosophies are different.

- The GPU is specialized for compute-intensive, massively data parallel computation (exactly what graphics rendering is about)
 - So, more transistors can be devoted to data processing rather than **data caching** and **flow control**



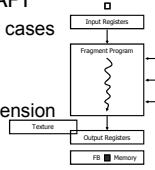
- Significant application-level speedup over uni-processor execution
 - No more “killer micros”
- Easy entrance
 - An initial, naïve code typically get at least 2-3X speedup
- Wide availability to end users
 - available on laptops, desktops, clusters, super-computers

GPGPU Movement

- General Purpose computation using GPU in applications other than 3D graphics
 - GPU accelerates critical path of application
- Data parallel algorithms leverage GPU attributes
 - Large data arrays, streaming throughput
 - Fine-grain SIMD parallelism
 - Low-latency floating point (FP) computation
- Applications – see //GPGPU.org
 - Game effects (FX) physics, image processing
 - Physical modeling, computational engineering, matrix algebra, convolution, correlation, sorting

GPGPU Constraints

- Dealing with graphics API
 - Working with the corner cases of the graphics API
- Addressing modes
 - Limited texture size/dimension
- Shader capabilities
 - Limited outputs
- Instruction sets
 - Lack of Integer & bit ops



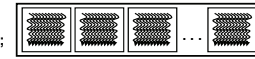
These have all changed with CUDA!

CUDA – C with no shader limitations!

- Integrated host+device app C program
 - Serial or modestly parallel parts in **host** C code
 - Highly parallel parts in **device** SPMD kernel C code

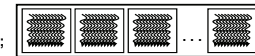
Serial Code (host)

Parallel Kernel (device)
KernelA<<< nBlk, nTid >>>(args);



Serial Code (host)

Parallel Kernel (device)
KernelB<<< nBlk, nTid >>>(args);



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CUDA Extends C

- Declspecs**
 - global, device, shared, local, constant
- Keywords**
 - threadIdx, blockIdx
- Intrinsics**
 - __syncthreads
- Runtime API**
 - Memory, symbol, execution management
- Function launch**

```

__device__ float filter[N];
__global__ void convolve (float *image) {
    __shared__ float region[M];
    ...
    region[threadIdx] = image[i];
    __syncthreads()
    ...
    image[j] = result;
}
// Allocate GPU memory
void *myimage = cudaMalloc(bytes)

// 100 blocks, 10 threads per block
convolve<<<100, 10>>>> (myimage);
    
```

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Arrays of Parallel Threads

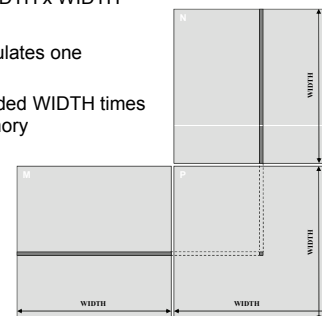
- A CUDA kernel is executed by an array of threads
 - All threads run the same code (SPMD)
 - Each thread has an ID that it uses to compute memory addresses and make control decisions

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A Simple Running Example Matrix Multiplication

- $P = M * N$ of size WIDTH x WIDTH
- Without tiling:
 - One thread calculates one element of P
 - M and N are loaded WIDTH times from global memory



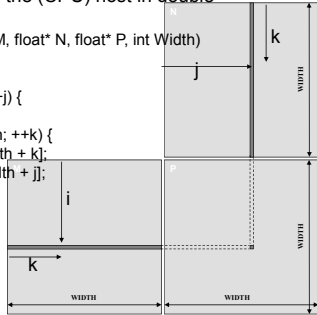
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Step 1: Matrix Multiplication A Simple Host Version in C

// Matrix multiplication on the (CPU) host in double precision

```
void MatrixMulOnHost(float* M, float* N, float* P, int Width)
{
    for (int i = 0; i < Width; ++i)
        for (int j = 0; j < Width; ++j) {
            double sum = 0;
            for (int k = 0; k < Width; ++k) {
                double a = M[i * Width + k];
                double b = N[k * Width + j];
                sum += a * b;
            }
            P[i * Width + j] = sum;
        }
}
```

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Step 2: Input Matrix Data Transfer (Host-side Code)

```
void MatrixMulOnDevice(float* M, float* N, float* P, int Width)
```

```
{
    int size = Width * Width * sizeof(float);
    float* Md, Nd, Pd;
    ...
    1. // Allocate and Load M, N to device memory
    cudaMalloc(&Md, size);
    cudaMemcpy(Md, M, size, cudaMemcpyHostToDevice);

    cudaMalloc(&Nd, size);
    cudaMemcpy(Nd, N, size, cudaMemcpyHostToDevice);

    // Allocate P on the device
    cudaMalloc(&Pd, size);
}
```

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Step 3: Output Matrix Data Transfer (Host-side Code)

```
2. // Kernel invocation code – to be shown later
...

3. // Read P from the device
cudaMemcpy(P, Pd, size, cudaMemcpyDeviceToHost);

// Free device matrices
cudaFree(Md); cudaFree(Nd); cudaFree(Pd);
}
```

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Step 4: Kernel Function

// Matrix multiplication kernel – per thread code

```
__global__ void MatrixMulKernel(float* Md, float* Nd, float* Pd, int Width)
{
    // 2D Thread ID
    int tx = threadIdx.x;
    int ty = threadIdx.y;

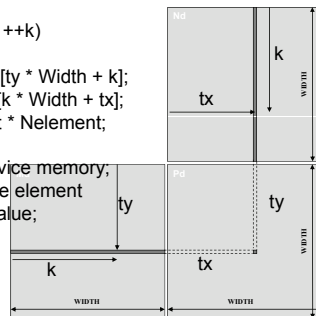
    // Pvalue is used to store the element of the matrix
    // that is computed by the thread
    float Pvalue = 0;
}
```

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Step 4: Kernel Function (cont.)

```
for (int k = 0; k < Width; ++k)
{
    float Melement = Md[ty * Width + k];
    float Nelement = Nd[k * Width + tx];
    Pvalue += Melement * Nelement;
}
// Write the matrix to device memory;
// each thread writes one element
Pd[ty * Width + tx] = Pvalue;
}
```

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Step 5: Kernel Invocation (Host-side Code)

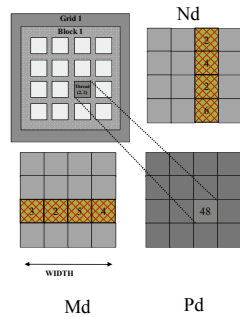
```
// Setup the execution configuration
dim3 dimBlock(Width, Width);
dim3 dimGrid(1, 1);

// Launch the device computation threads!
MatrixMulKernel<<<dimGrid, dimBlock>>>>(Md, Nd, Pd);
```

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Only One Thread Block Used

- One Block of threads compute matrix Pd
 - Each thread computes one element of Pd
- Each thread
 - Loads a row of matrix Md
 - Loads a column of matrix Nd
 - Perform one multiply and addition for each pair of Md and Nd elements
 - Compute to off-chip memory access ratio close to 1:1 (not very high)
- Size of matrix limited by the number of threads allowed in a thread block



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