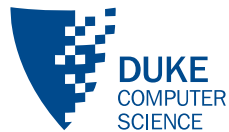


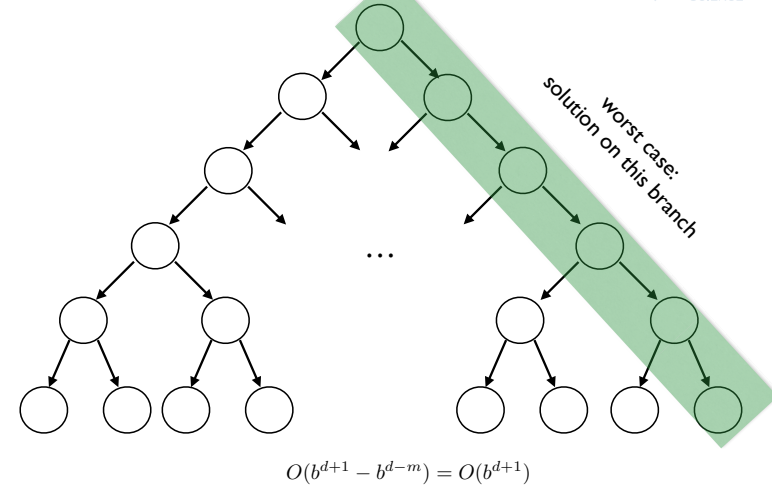
Informed Search

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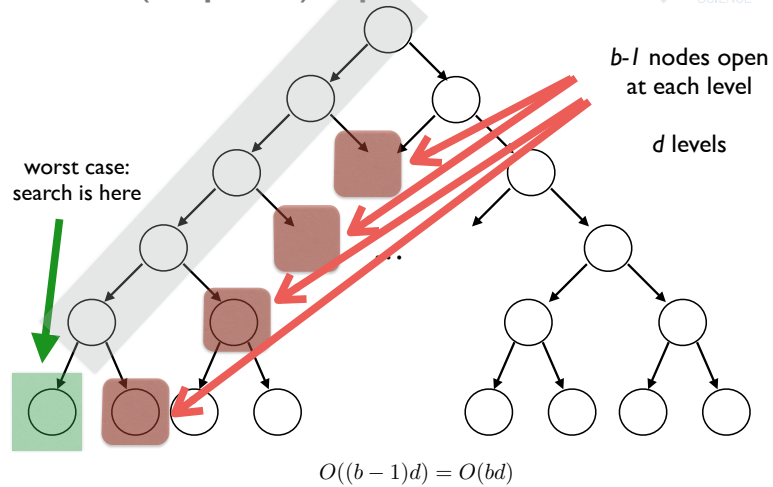


Spring 2015

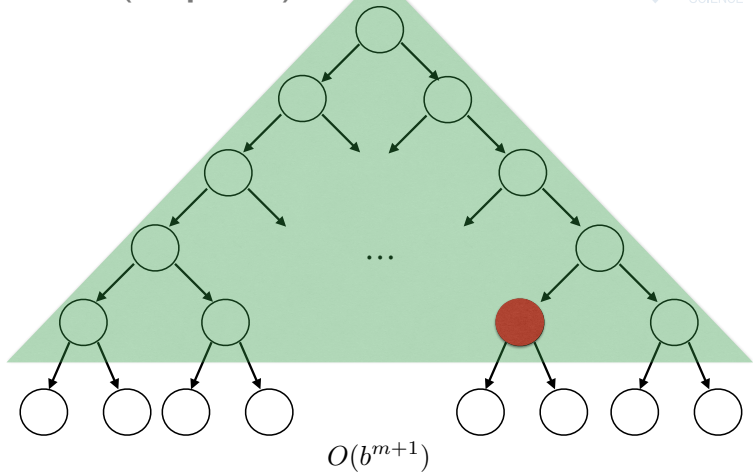
DFS: (Reprise) Time



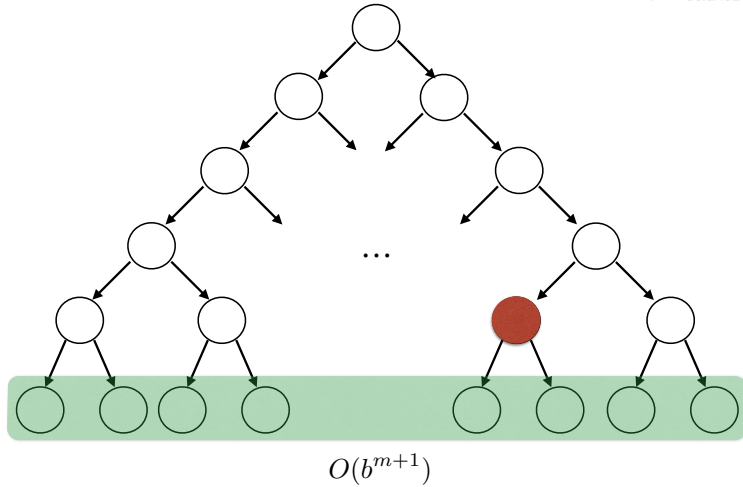
DFS (Reprise): Space



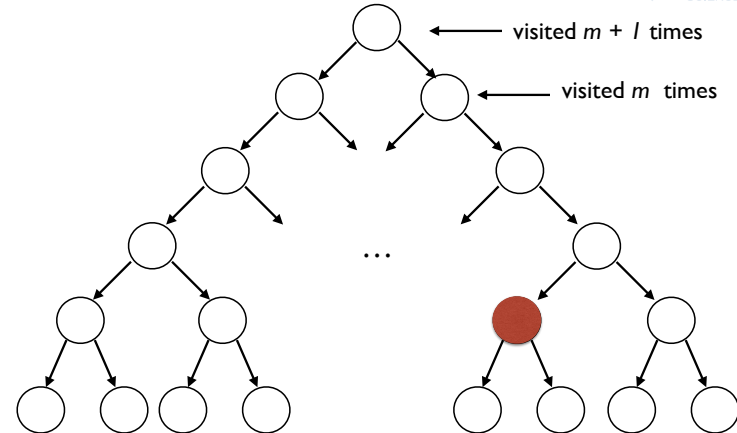
BFS (Reprise): Time



BFS: (Reprise) Space



IDS (Reprise)



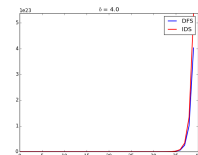
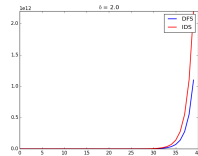
IDS (Reprise)



$$\sum_{i=0}^m b^i(m-i+1) = \frac{b(b^{m+1} - m - 2) + m + 1}{(b-1)^2}$$

nodes at level i # revisits

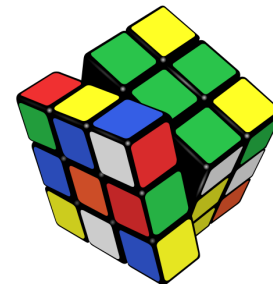
DFS worst case: $\frac{b^{m+1} - 1}{b - 1}$



Informed Search



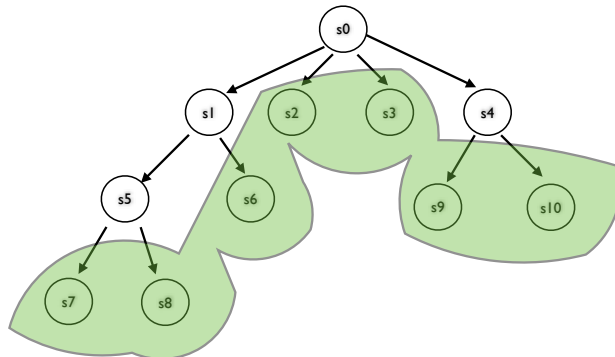
What if we *know something* about the search?



Informed Search



Key thing in search is *managing the frontier*.



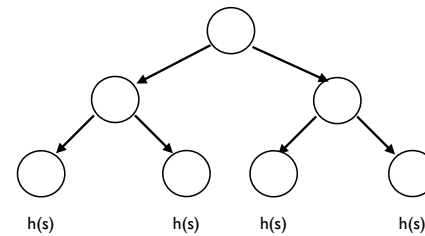
Use domain knowledge to order the nodes that you open.

Heuristics



Key idea: *heuristic function*.

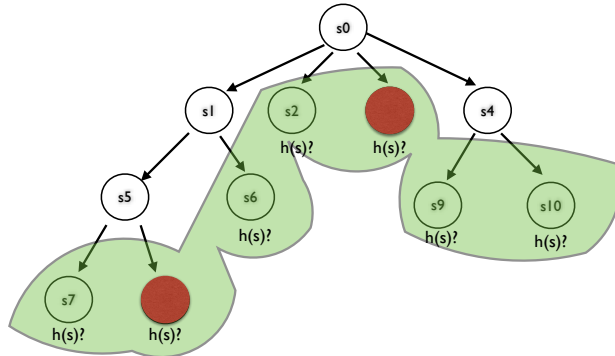
- $h(s)$ - estimates cost-to-go
- Cost to go *from* state to solution.
- Problem specific.



Greed



What if we expand the node with lowest $h(s)$?



A*



A* algorithm:

- Also $g(s)$ - cost so far (start to s).
- Expand s that minimizes $g(s) + h(s)$
- Manage frontier as priority queue.
- Admissible heuristic: *never overestimates* cost.

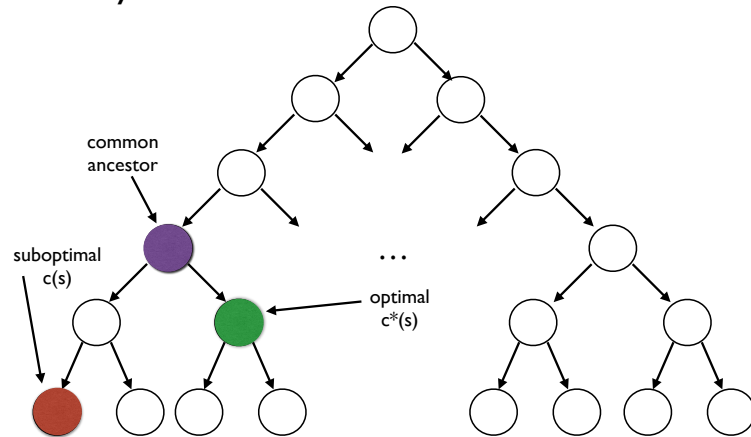
$$h(s) \leq h^*(s)$$

- $h(s) = 0$ if s is a goal state, so $g(s) + h(s) = c(s)$
- If h is admissible, A* is optimal.
- If $h(s)$ is exact, runs in $O(bd)$ time.

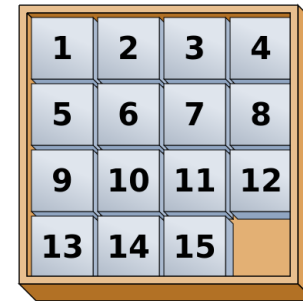
Admissible Heuristics



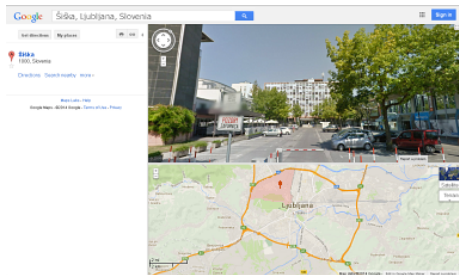
Proof by contradiction.



Example Heuristic



Example Heuristics



More on Heuristics



Heuristic $h1$ dominates $h2$ if $h1(s) \geq h2(s)$ for all s .

- Is $h1$ or $h2$ better? (If they're both admissible.)

How might you combine two heuristics?

What is $h(s) = k$ (constant) for all s ?

More on Heuristics



A* is optimally efficient: any algorithm using h **must** expand the nodes A* expands.

Why?

More on Heuristics



Ideal heuristics:

- Fast to compute.
- Close to real costs.
- Some programs *automatically generate* heuristics.