

# Game Theory

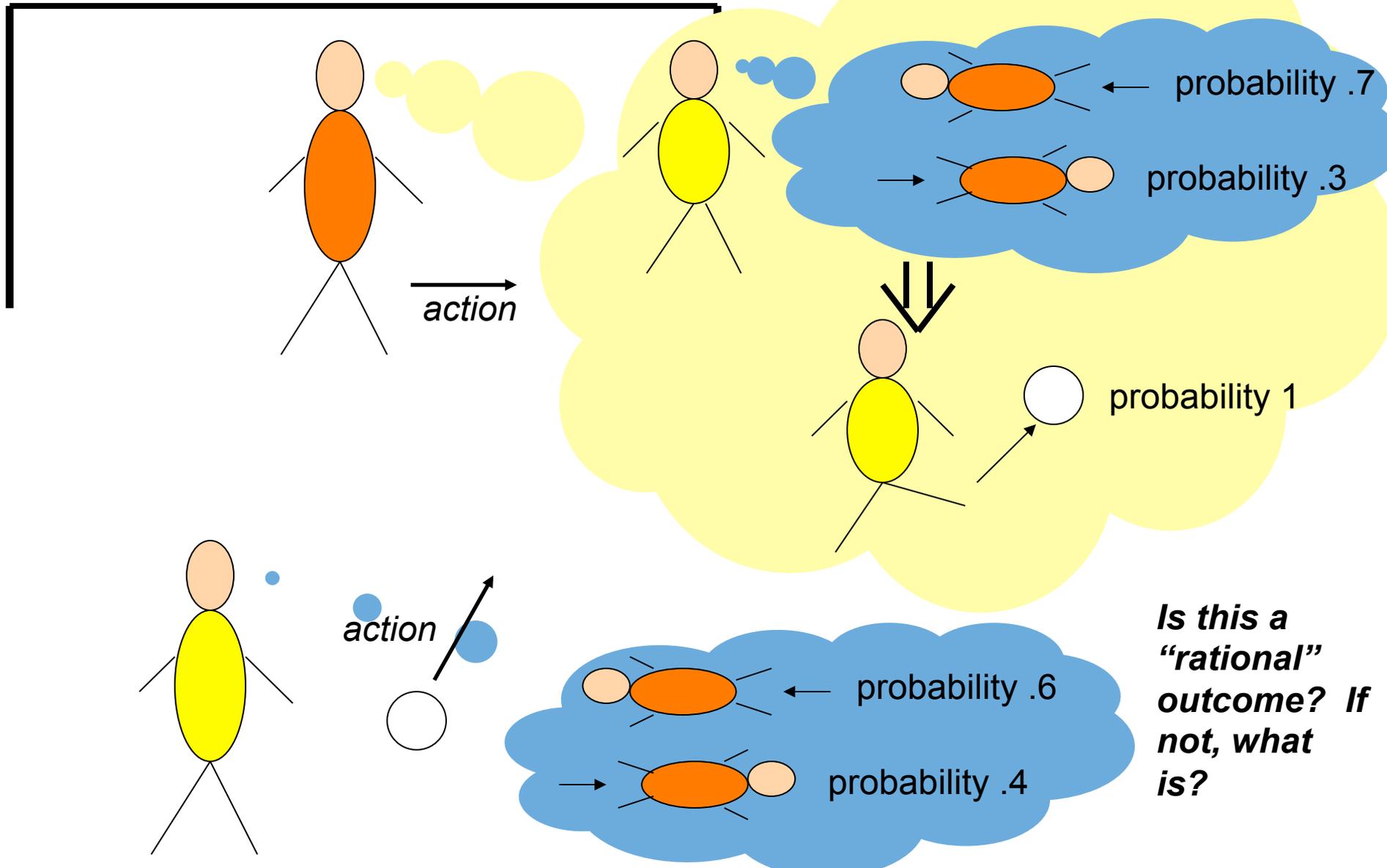
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With thanks to Ron Parr and Vince Conitzer for some contents

# What is Game Theory?

- Settings where multiple agents each have different preferences and set of actions they can take
- Each agent's utility (potentially) depends on all agents' actions
  - What is optimal for one agent depends on what other agents do!
- Game theory studies how agents can rationally form **beliefs** over what other agents will do, and (hence) how agents should **act**

# Penalty Kick Example



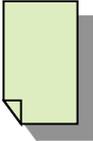
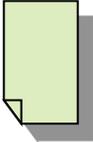
***Is this a "rational" outcome? If not, what is?***

# Overview

- Zero-sum games from Adversarial Search lecture
  - Minimax, alpha-beta pruning
- General-sum games
- Normal form vs. Extensive form games
  - Table specifying action-payoff vs. game tree with sequence of actions (and information sets)
- Solving games: dominance, iterated dominance, mixed strategy, Nash Equilibrium

# Rock-paper-scissors (zero-sum game)

Column player aka.  
player 2  
(simultaneously)  
chooses a column



|       |       |       |
|-------|-------|-------|
| 0, 0  | -1, 1 | 1, -1 |
| 1, -1 | 0, 0  | -1, 1 |
| -1, 1 | 1, -1 | 0, 0  |

An arrow points to the bottom-middle cell (1, -1) in the matrix.

Row player  
aka. player 1  
chooses a row

A row or column is  
called an **action** or  
(pure) **strategy**

Row player's utility is always listed first, column player's second

**Zero-sum** game: the utilities in each entry sum to 0 (or a constant)  
Three-player game would be a 3D table with 3 utilities per entry, etc.

# General-sum games

- You could still play a minimax strategy in general-sum games
  - pretend that the opponent is only trying to hurt you!
- But this is not rational:

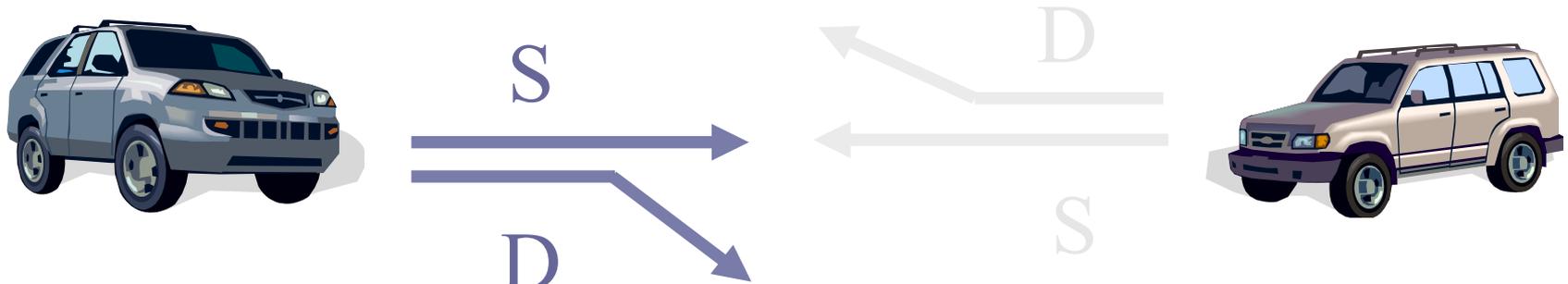
|      |      |
|------|------|
| 0, 0 | 3, 1 |
| 1, 0 | 2, 1 |

not zero-sum

- If Column was trying to hurt Row, Column would play Left, so Row should play Down
- In reality, Column will play Right (strictly dominant), so Row should play up

# Chicken

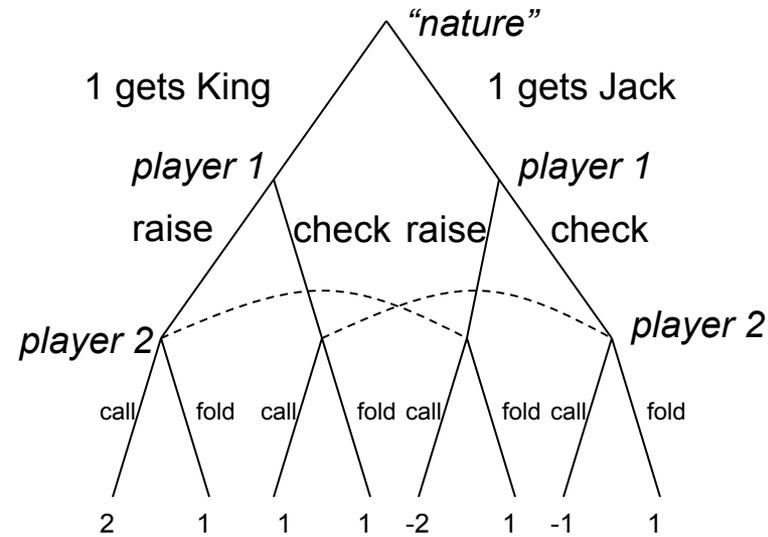
- Two players drive cars towards each other
- If one player goes straight, that player wins
- If both go straight, they both die



|   | D     | S      |
|---|-------|--------|
| D | 0, 0  | -1, 1  |
| S | 1, -1 | -5, -5 |

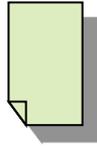
not zero-sum

# A “poker-like” game

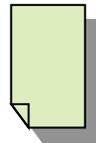
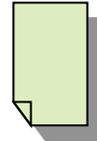


|    | cc      | cf        | fc    | ff    |
|----|---------|-----------|-------|-------|
| rr | 0, 0    | 0, 0      | 1, -1 | 1, -1 |
| rc | .5, -.5 | 1.5, -1.5 | 0, 0  | 1, -1 |
| cr | -.5, .5 | -.5, .5   | 1, -1 | 1, -1 |
| cc | 0, 0    | 1, -1     | 0, 0  | 1, -1 |

# Rock-paper-scissors – Seinfeld variant



MICKEY: All right, rock beats paper!  
(Mickey smacks Kramer's hand for losing)  
KRAMER: I thought paper covered rock.  
MICKEY: Nah, rock flies right through paper.  
KRAMER: What beats rock?  
MICKEY: (looks at hand) Nothing beats rock.

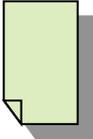
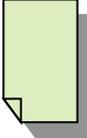


|          |       |       |          |
|----------|-------|-------|----------|
|          | Rock  | Paper | Scissors |
| Rock     | 0, 0  | 1, -1 | 1, -1    |
| Paper    | -1, 1 | 0, 0  | -1, 1    |
| Scissors | -1, 1 | 1, -1 | 0, 0     |

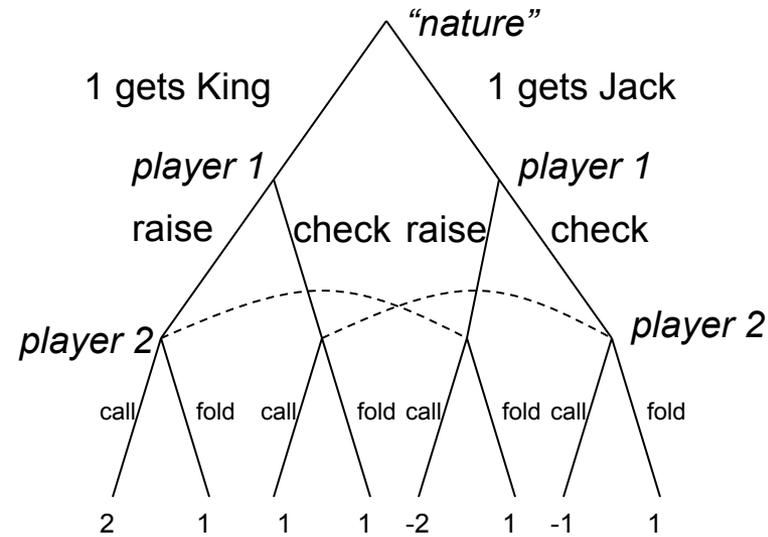
# Dominance

- Player  $i$ 's strategy  $s_i$  **strictly dominates**  $s_i'$  if
  - for any  $s_{-i}$ ,  $u_i(s_i, s_{-i}) > u_i(s_i', s_{-i})$
- $s_i$  **weakly dominates**  $s_i'$  if
  - for any  $s_{-i}$ ,  $u_i(s_i, s_{-i}) \geq u_i(s_i', s_{-i})$ ; and
  - for some  $s_{-i}$ ,  $u_i(s_i, s_{-i}) > u_i(s_i', s_{-i})$

*-i = "the player(s) other than i"*

|   |  |  |  |
|---|---|---|---|
|    | 0, 0  | 1, -1   | 1, -1   |
|  | -1, 1   | 0, 0  | -1, 1   |
|  | -1, 1   | 1, -1   | 0, 0  |

# Back to the poker like game



|    | cc                 | cf                   | fc    | ff               |
|----|--------------------|----------------------|-------|------------------|
| rr | 0, 0               | <del>0, 0</del>      | 1, -1 | <del>1, -1</del> |
| rc | .5, -.5            | <del>1.5, -1.5</del> | 0, 0  | <del>1, -1</del> |
| cr | <del>-.5, .5</del> | <del>-.5, .5</del>   | 1, -1 | <del>1, -1</del> |
| cc | 0, 0               | <del>1, -1</del>     | 0, 0  | <del>1, -1</del> |

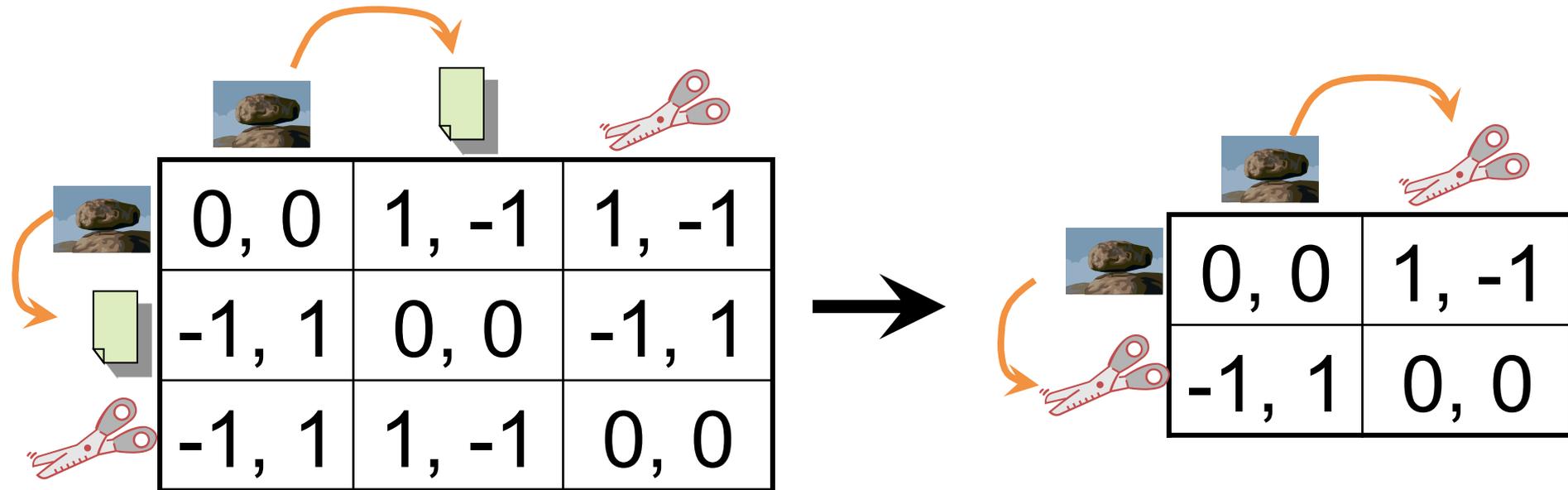
# Prisoner's Dilemma

- Pair of criminals has been caught
- District attorney has evidence to convict them of a minor crime (1 year in jail); knows that they committed a major crime together (3 years in jail) but cannot prove it
- Offers them a deal:
  - If both confess to the major crime, they each get a 1 year reduction
  - If only one confesses, that one gets 3 years reduction

|               | confess | don't confess |
|---------------|---------|---------------|
| confess       | -2, -2  | 0, -3         |
| don't confess | -3, 0   | -1, -1        |

# Iterated Dominance

- Iterated dominance: remove (strictly/weakly) dominated strategy, repeat
- Iterated strict dominance on Seinfeld's RPS:

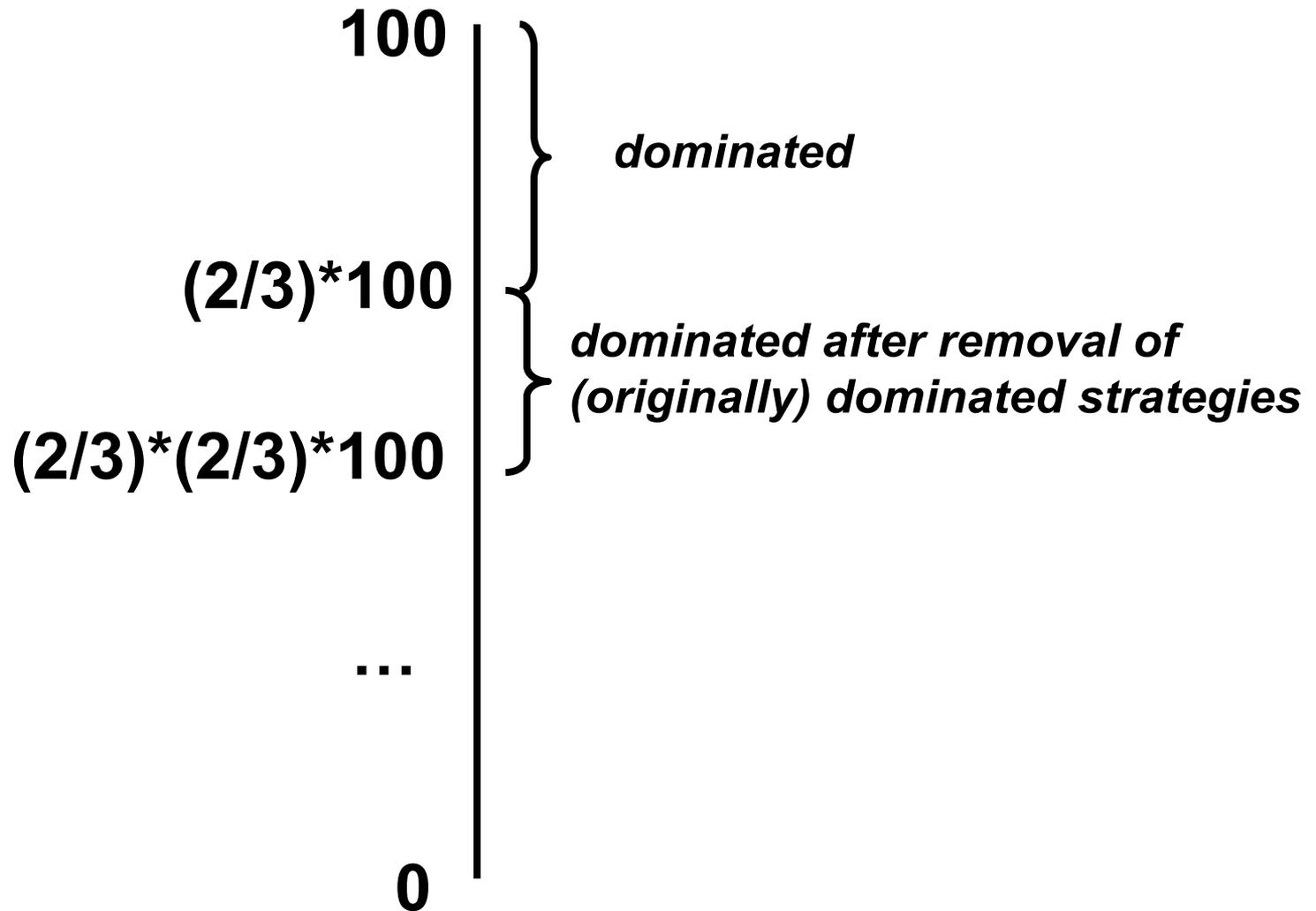


# “2/3 of the average” game

- Everyone writes down a number between 0 and 100
- Person closest to 2/3 of the average wins
- Example:
  - A says 50
  - B says 10
  - C says 90
  - Average(50, 10, 90) = 50
  - 2/3 of average = 33.33
  - A is closest ( $|50 - 33.33| = 16.67$ ), so A wins

Try?

# “2/3 of the average” via dominance



# Mixed strategy

- **Mixed strategy** for player  $i$  = probability distribution over player  $i$ 's (pure) strategies
- E.g.  $1/3$  ,  $1/3$  ,  $1/3$  
- Example of dominance by a mixed strategy:

|       |      |      |
|-------|------|------|
| $1/2$ | 3, 0 | 0, 0 |
| $1/2$ | 0, 0 | 3, 0 |
|       | 1, 0 | 1, 0 |

An orange bracket on the left side of the table groups the first two rows, with a label  $1/2$  next to each row. An orange arrow points from the bottom of this bracket to the third row of the table.

# Best-Response

- Let  $A$  be a matrix of player 1's payoffs
- Let  $s_2$  be a mixed strategy for player 2
- $As_2$  = vector of expected payoffs for each strategy for player 1
- Highest entry indicates **best response** for player 1
- Any mixture of ties is also BR
- Generalizes to  $>2$  players

|       |        |            |
|-------|--------|------------|
| 0, 0  | -1, 1  | $\sigma_2$ |
| 1, -1 | -5, -5 |            |

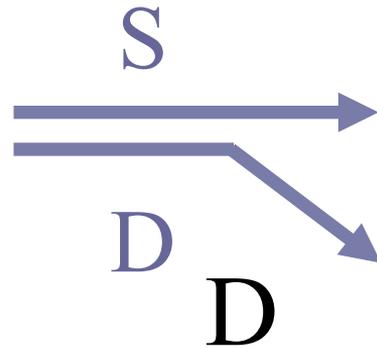
# Nash Equilibrium

## [Nash 50]



- A vector of strategies (one for each player) = a **strategy profile**
- Strategy profile  $(\sigma_1, \sigma_2, \dots, \sigma_n)$  is a **Nash equilibrium** if each  $\sigma_i$  is a **best response** to  $\sigma_{-i}$
- Does not say anything about multiple agents changing their strategies at the same time
- In any (finite) game, at least one Nash equilibrium (possibly using mixed strategies) exists **[Nash 50]**

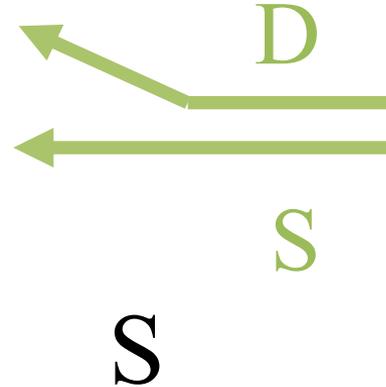
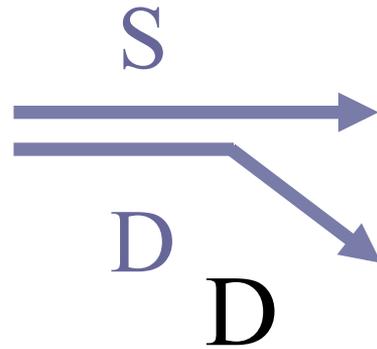
# NE of “Chicken”



|   |       |        |
|---|-------|--------|
|   | D     | S      |
| D | 0, 0  | -1, 1  |
| S | 1, -1 | -5, -5 |

- (D, S) and (S, D) are Nash equilibria
  - They are **pure-strategy Nash equilibria**: nobody randomizes
  - They are also **strict Nash equilibria**: changing your strategy will make you strictly worse off
- No other pure-strategy Nash equilibria

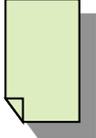
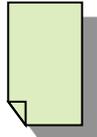
# Equilibrium Selection



|   |       |        |   |
|---|-------|--------|---|
|   |       | D      | S |
| D | 0, 0  | -1, 1  |   |
| S | 1, -1 | -5, -5 |   |

- (D, S) and (S, D) are Nash equilibria
- Which do you play?
- What if player 1 assumes (S, D), player 2 assumes (D, S)
- Play is (S, S) = (-5, -5)!!!
- This is the **equilibrium selection** problem

# Rock-paper-scissors revisited

|   |   |  |   |
|---|---|--|---|
|   |  |  |  |
|  | 0, 0  | -1, 1  | 1, -1   |
|  | 1, -1   | 0, 0   | -1, 1   |
|  | -1, 1   | 1, -1  | 0, 0  |

- Any pure-strategy Nash equilibria?
- But it has a **mixed-strategy Nash equilibrium**:  
Both players put probability  $1/3$  on each action
- If the other player does this, every action will give you expected utility 0
  - Might as well randomize

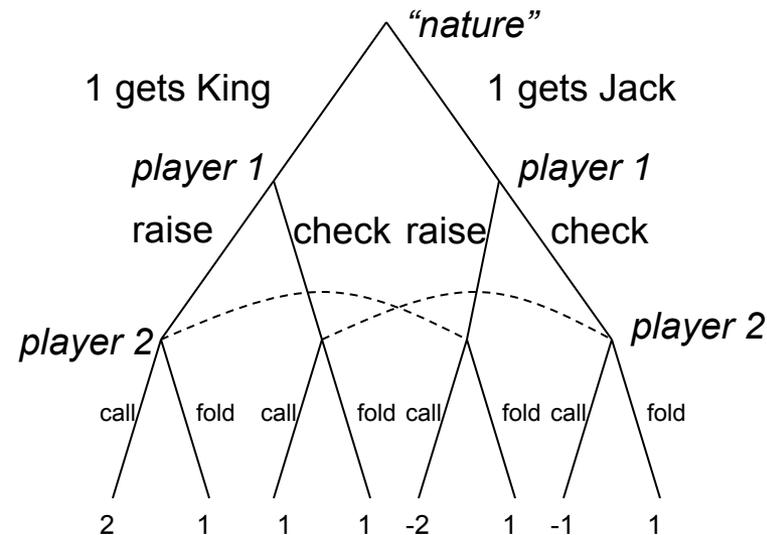
# NE of “Chicken”

|   | D     | S      |
|---|-------|--------|
| D | 0, 0  | -1, 1  |
| S | 1, -1 | -5, -5 |

- Is there a Nash equilibrium that uses mixed strategies -- say, where player 1 uses a mixed strategy?
- *If a mixed strategy is a best response, then all of the pure strategies that it randomizes over must also be best responses*
- So we need to make player 1 **indifferent** between D and S
- Player 1's utility for playing D =  $-p^c_S$
- Player 1's utility for playing S =  $p^c_D - 5p^c_S = 1 - 6p^c_S$
- So we need  $-p^c_S = 1 - 6p^c_S$  which means  $p^c_S = 1/5$
- Then, player 2 needs to be indifferent as well
- Mixed-strategy Nash equilibrium: ((4/5 D, 1/5 S), (4/5 D, 1/5 S))
  - People may die! Expected utility -1/5 for each player

$-p^c_S$  = probability that column player plays s

# The “poker-like game” again



|               |    | $\frac{2}{3}$<br>cc | cf                 | $\frac{1}{3}$<br>fc | ff               |
|---------------|----|---------------------|--------------------|---------------------|------------------|
| $\frac{1}{3}$ | rr | 0, 0                | <del>0, 0</del>    | 1, -1               | <del>1, -1</del> |
| $\frac{2}{3}$ | rc | .5, -.5             | 1.5, -1.5          | 0, 0                | <del>1, -1</del> |
|               | cr | <del>-.5, .5</del>  | <del>-.5, .5</del> | 1, -1               | <del>1, -1</del> |
|               | cc | 0, 0                | <del>1, -1</del>   | 0, 0                | <del>1, -1</del> |

- To make player 1 indifferent between rr and rc, we need:  

$$\text{utility for rr} = 0 \cdot P(\text{cc}) + 1 \cdot (1 - P(\text{cc})) = .5 \cdot P(\text{cc}) + 0 \cdot (1 - P(\text{cc})) = \text{utility for rc}$$
 That is,  $P(\text{cc}) = \frac{2}{3}$
- To make player 2 indifferent between cc and fc, we need:  

$$\text{utility for cc} = 0 \cdot P(\text{rr}) + (-.5) \cdot (1 - P(\text{rr})) = -1 \cdot P(\text{rr}) + 0 \cdot (1 - P(\text{rr})) = \text{utility for fc}$$
 That is,  $P(\text{rr}) = \frac{1}{3}$

# Computational considerations

- Zero-sum games - solved efficiently as LP
- General sum games may require exponential time (in # of actions) to find a single equilibrium (no known efficient algorithm and good reasons to suspect that none exists)
- Some better news: Despite bad worst-case complexity, many games can be solved quickly

# Extensions

- Partial information
- Uncertainty about the game parameters, e.g., payoffs (Bayesian games)
- Repeated games: Simple learning algorithms can converge to equilibria in some repeated games
- Multistep games with distributions over next states (game theory + MDPs = stochastic games)
- Multistep + partial information (Partially observable stochastic games)
  
- Game theory is so general, that it can encompass essentially all aspects of strategic, multiagent behavior, e.g., negotiating, threats, bluffs, coalitions, bribes, etc.