CS 356: Computer Network Architectures

Lecture 9: The Internet Protocol (IP) Ch 3.2

Xiaowei Yang

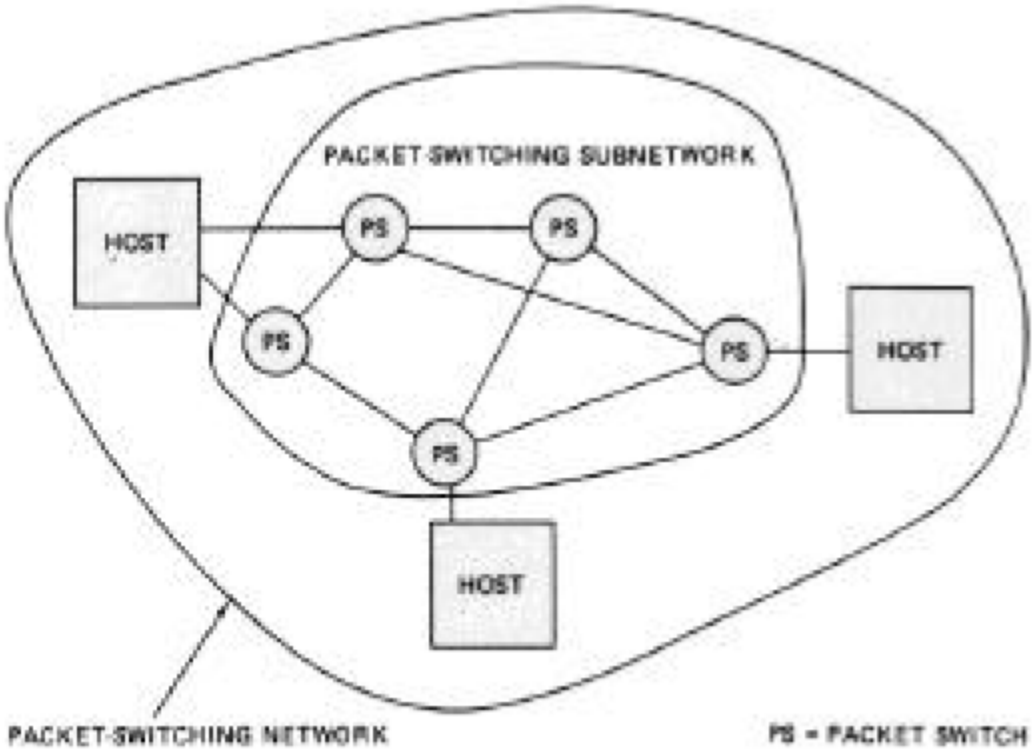
xwy@cs.duke.edu

Overview

- History of IP
- IP header format
- IP addressing
- IP forwarding
  - Forwarding algorithm
  - Fragmentation

History of the Internet

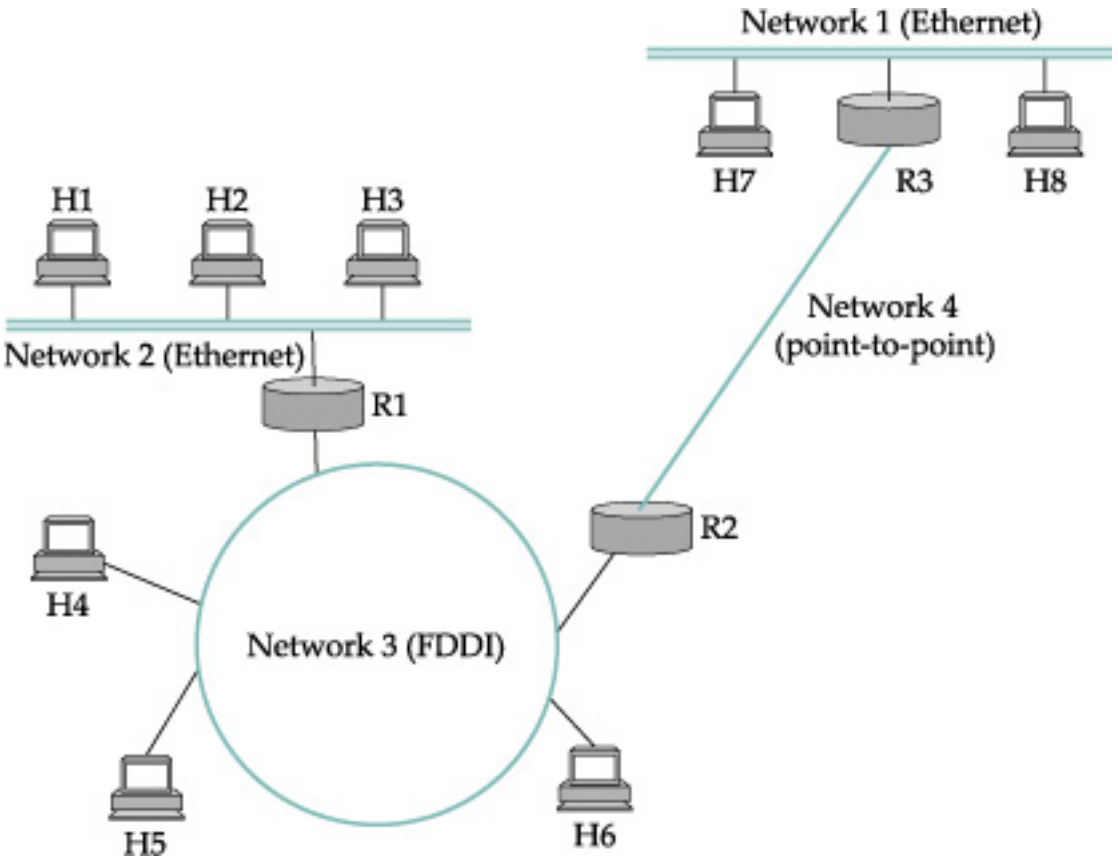
- Original design goal: Interconnecting different networks
- Many different types of packet switch networks
  - ARPANET, packet satellite networks, ground-based packet radio networks, and other networks.
- Each has
  - Hosts, packet switches, processes
  - A protocol for communication
- Q: what would you do differently given such a design task?



Challenges

1. Different addressing schemes and host communication protocols
   - Ethernet, FDDI, ATM
2. Different Maximum Transmission Units (MTUs)
3. Different success or failure indicators
4. End-to-end reliability: failures may occur at each network
5. Different control protocols
   - Status information, routing, fault detection/isolation

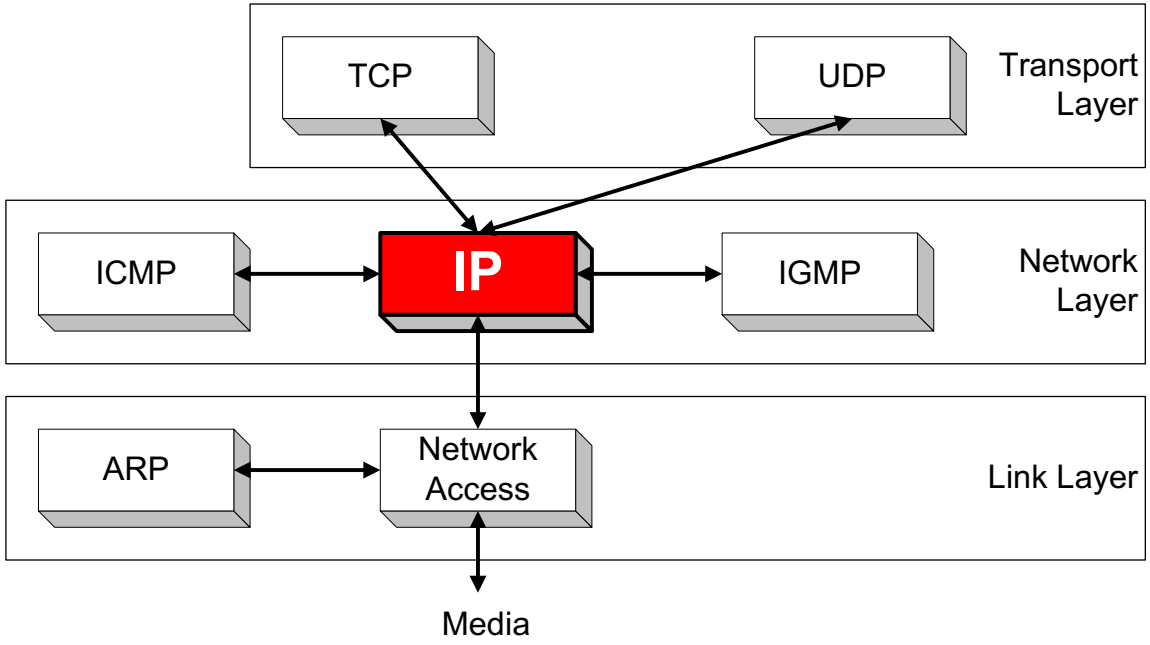
Inter-networking



- One level of indirection
  - Routers interface different networks
- Uniform addressing (IP)
- Routers send packets to their destination IP addresses

Internet Protocol

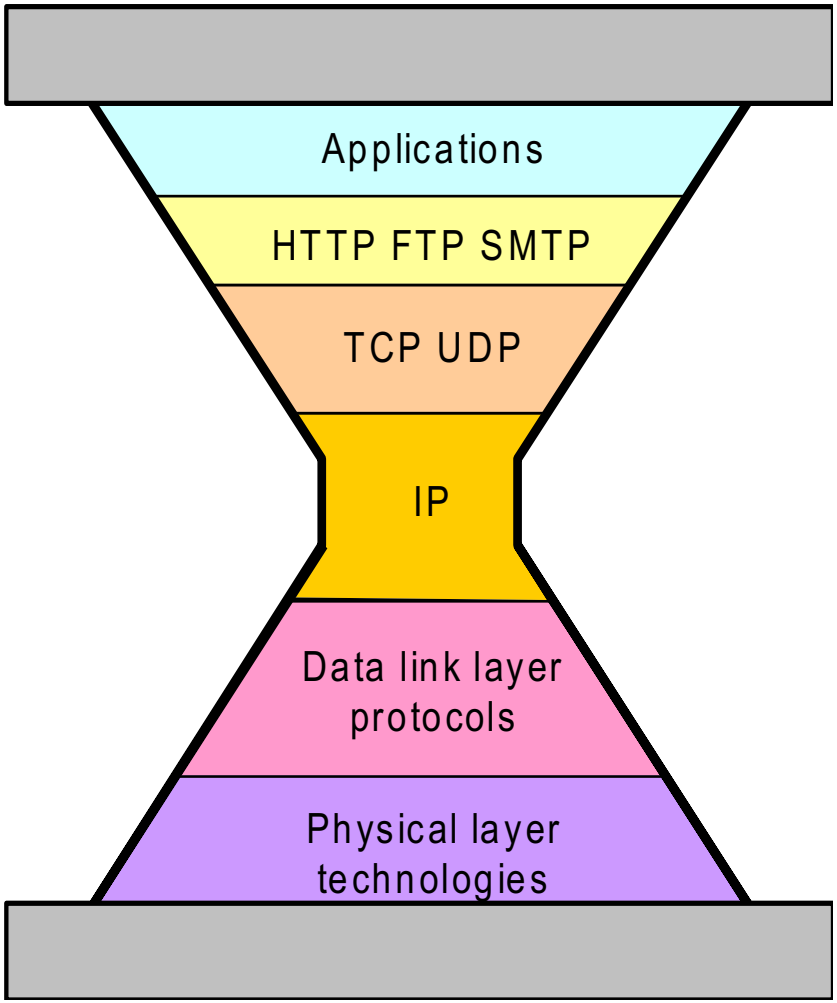
- IP (Internet Protocol) is a Network Layer Protocol
- IP's current version is Version 4 (IPv4). It is specified in RFC 791.
- IPv6 is also deployed



IP: the thin waist of the hourglass

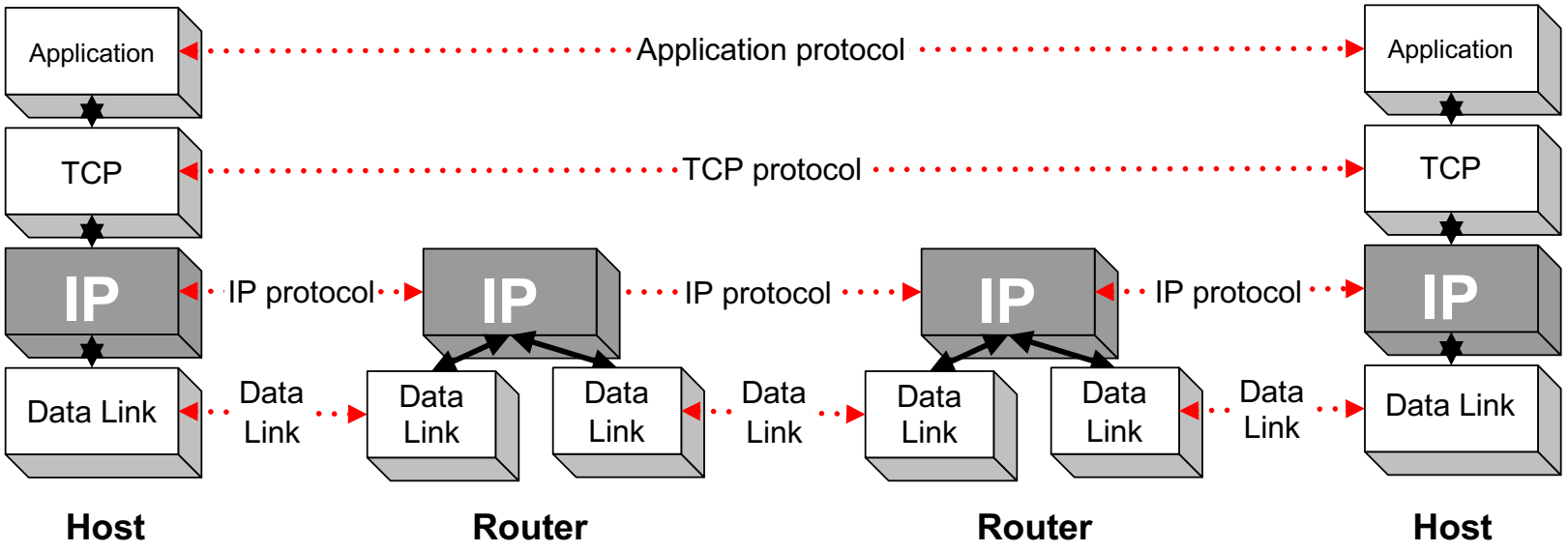
- **IP is the waist of the hourglass of the Internet protocol architecture**

- Multiple higher-layer protocols
- Multiple lower-layer protocols
- Only one protocol at the network layer.
- What is the advantage of this architecture?
  - To avoid the N * M problem

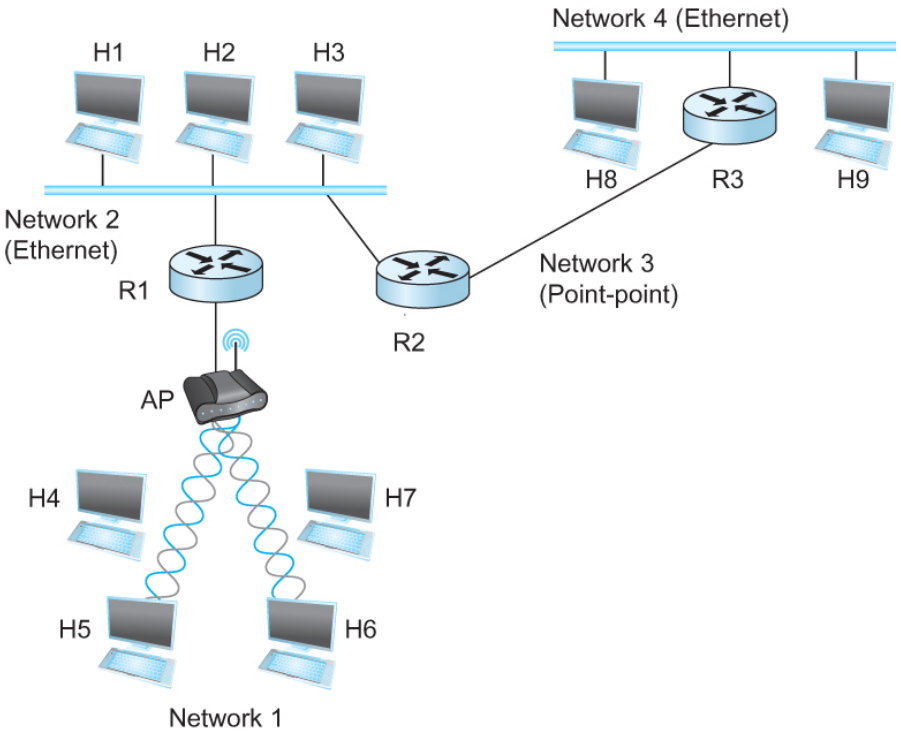


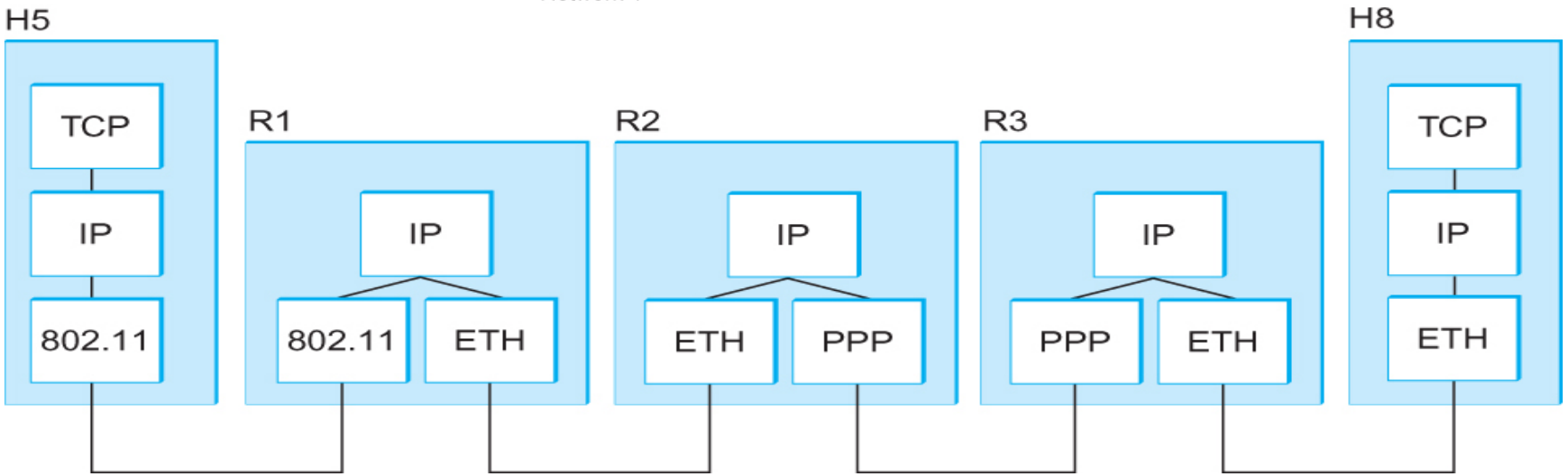
Application protocol

- Routers look at a packet's IP header and link layer header



A simple network





IP Service Model

- Delivery service of IP is minimal
- IP provides an unreliable connectionless best effort service (also called: “datagram service”).
  - **Unreliable**
  - **Connectionless**
  - **Best effort**
- Consequences
  - Loss, out of order, and duplicate must be handled at the upper layer

Basic IP router functions

- Things you need to understand to do lab2
  - Internet protocol
    - IP header
    - IP addressing
    - IP forwarding
  - Address resolution protocol
  - Error reporting and control
    - Internet Control Message Protocol

IP header format

IPv4 Header Format

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Version				IHL				DSCP						ECN		Total Length															
4	32	Identification																Flags			Fragment Offset											
8	64	Time To Live								Protocol								Header Checksum															
12	96	Source IP Address																															
16	128	Destination IP Address																															
20	160	Options (if IHL > 5)																															
24	192																																
28	224																																
32	256																																

- 20 bytes fixed length header + variable length options
- **Internet Header Length (IHL 4 bits):** the length of header in 32-bit words
  - Maximum header length?

IP header format

IPv4 Header Format

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
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- **DSCP (Differentiated Services Code Point 6 bits):** old Type of Service
  - Real-time, VoIP
- **Explicit Congestion Notification (ECN)**
  - Early Congestion notice

IP header format

IPv4 Header Format

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- **Total length (16 bits):** packet length in bytes, including the header
  - 65535 bytes
  - Fragmentation and reassembly

IP header format

IPv4 Header Format

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
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- **Identification (16 bits):** Unique datagram identifier from a host
  - Incremented whenever a datagram is transmitted (in some OS)
  - Used by many researchers for various purposes

IP header format

IPv4 Header Format

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
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- **Flags (3 bits):**
  - First bit always set to 0
  - DF bit (Do not fragment)
  - MF bit (More fragments)
- **Fragment offset (13 bits)**
- **Identification, Flags, Fragment offset**
  - fragmentation and assembly

IP header format

IPv4 Header Format

Offsets	Octet	0								1								2								3							
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- **Time To Live (TTL) (1byte):**
  - Specifies the longest path before a datagram is dropped
  - Role of TTL field: Ensure that a packet is eventually dropped when a routing loop occurs

  Used as follows:
  - Sender sets the value (e.g., 64)
  - Each router decrements the value by 1
  - When the value reaches 0, the datagram is dropped

IP header format

IPv4 Header Format

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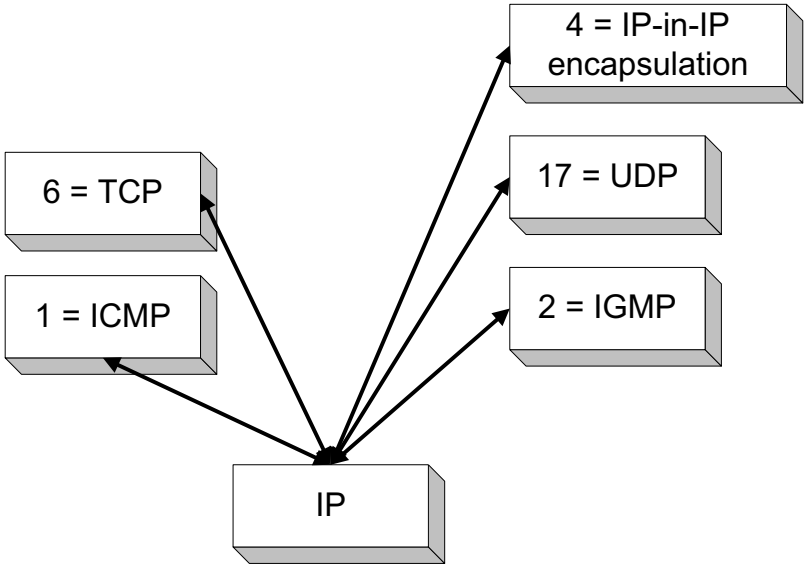
- **Protocol (1 byte):**
  - Specifies the higher-layer protocol.
  - De-multiplexing to higher layers.



IP header format

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- **Header checksum (16 bits):** header checksum
  - Header only
  - Must be computed at every hop!

Fields of the IP Header

- **Options:**
  - Record Route: each router that processes the packet adds its IP address to the header.
  - Timestamp: each router that processes the packet adds its IP address and time to the header.
  - (loose) Source Routing: specifies a list of routers that must be traversed.
  - (strict) Source Routing: specifies a list of the only routers that can be traversed.
  - IP options increase routers processing overhead
- **Padding:** Padding bytes are added to ensure that header ends on a 4-byte boundary

Global IP addresses

What is an IP Address?

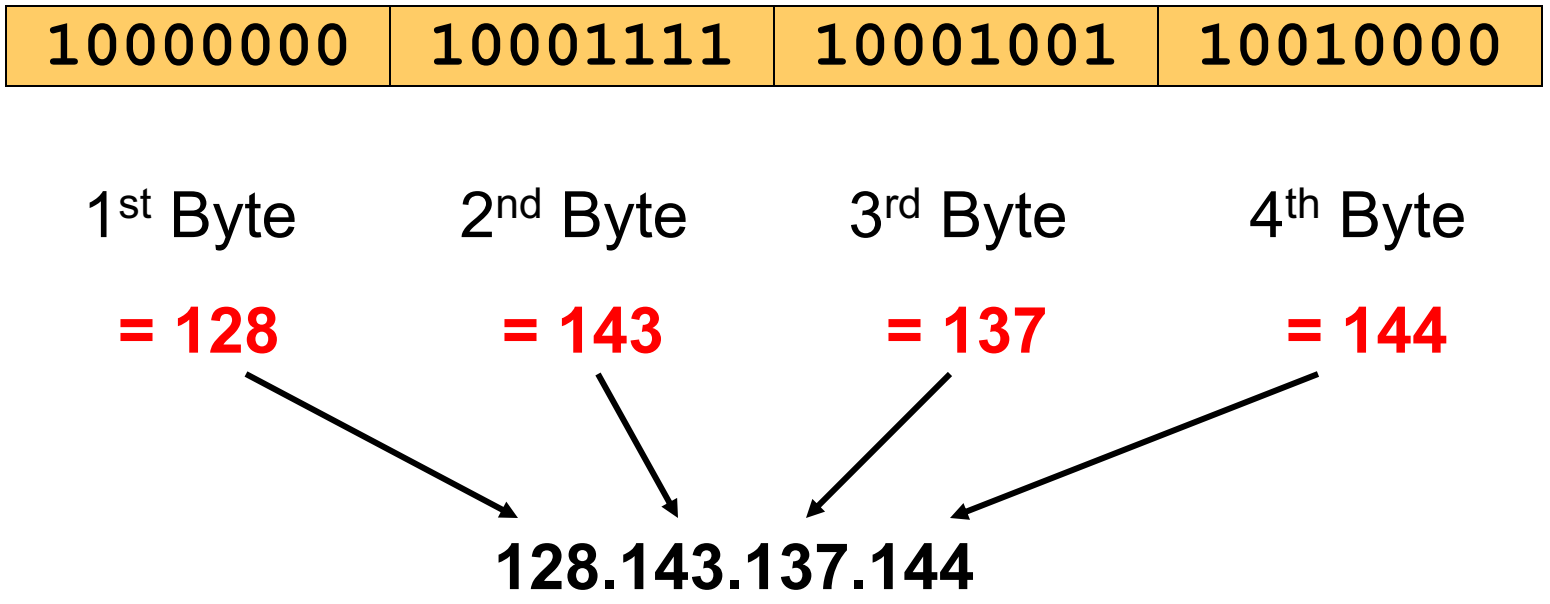
- An IP address is a unique global identifier for a network interface
  - An IP address uniquely identifies a network location
- Routers forwards a packet based on the destination address of the packet
- Uniqueness ensures global reachability

IP versions

- IPv4 (32-bit)
  - Classful IP addresses (obsolete)
  - Classless inter-domain routing (CIDR) (RFC 854, current standard)
- IP Version 6 addresses (128-bit)

Dotted Decimal Notation

- Each byte is identified by a decimal number in the range [0…255]:



Structure of an IP address

- An IP address has a structure
  - Network prefix identifies a network
  - Host number identifies a specific host interface
- Improves the scalability of routing
  - Scales better than flat addresses

How long is a network prefix?

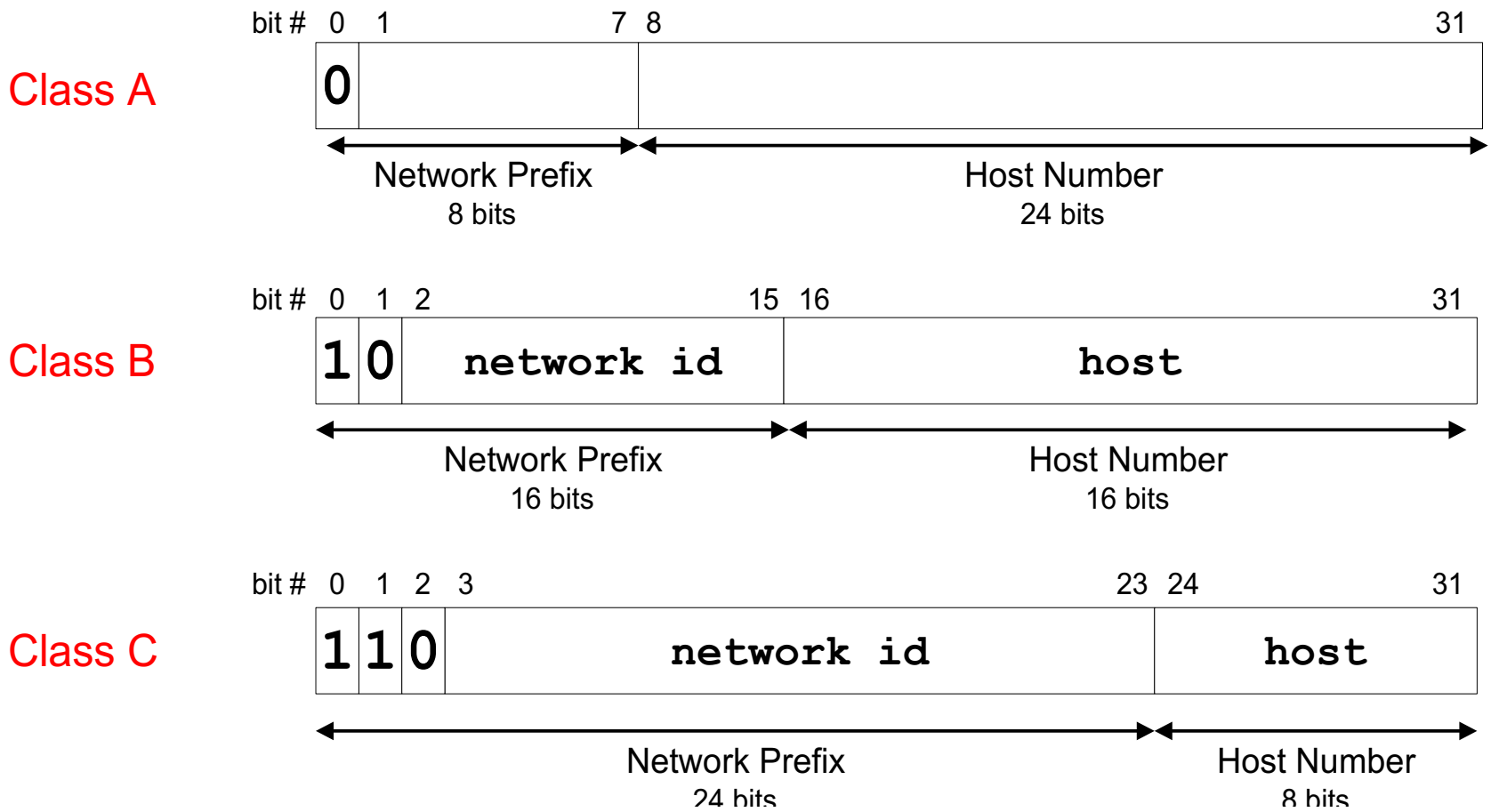
- **Before 1993:** The network prefix is implicitly defined (**class-based addressing**)

- **After 1993:** The network prefix is indicated by a **netmask**

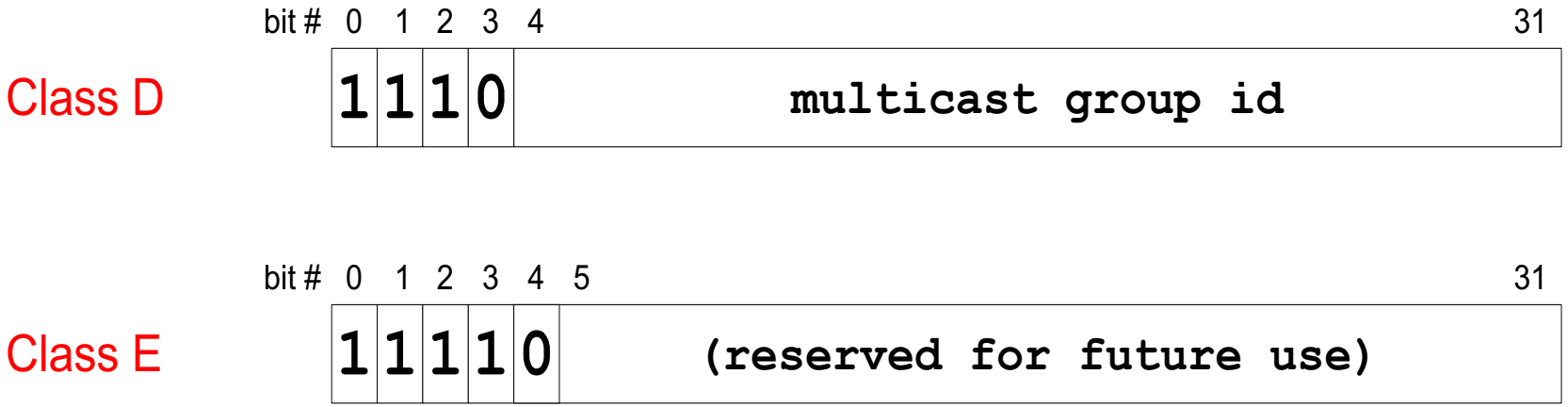
Before 1993: Class-based addressing

- The Internet address space was divided up into classes:
  - **Class A:** Network prefix is 8 bits long
  - **Class B:** Network prefix is 16 bits long
  - **Class C:** Network prefix is 24 bits long
  - Class D is multicast address
  - Class E is reserved

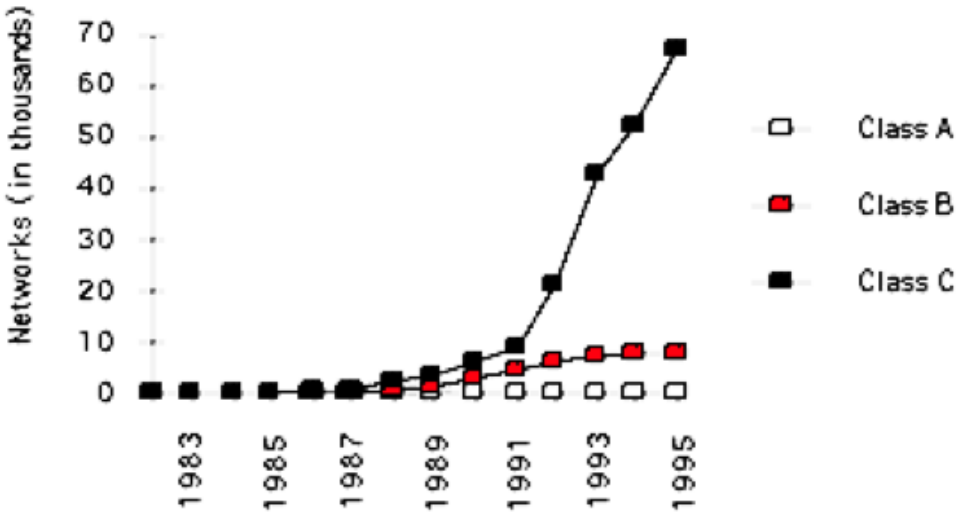
Classful IP Addresses (before 1993)



Classful IP Addresses (before 1993)



Problems with Classful IP Addresses

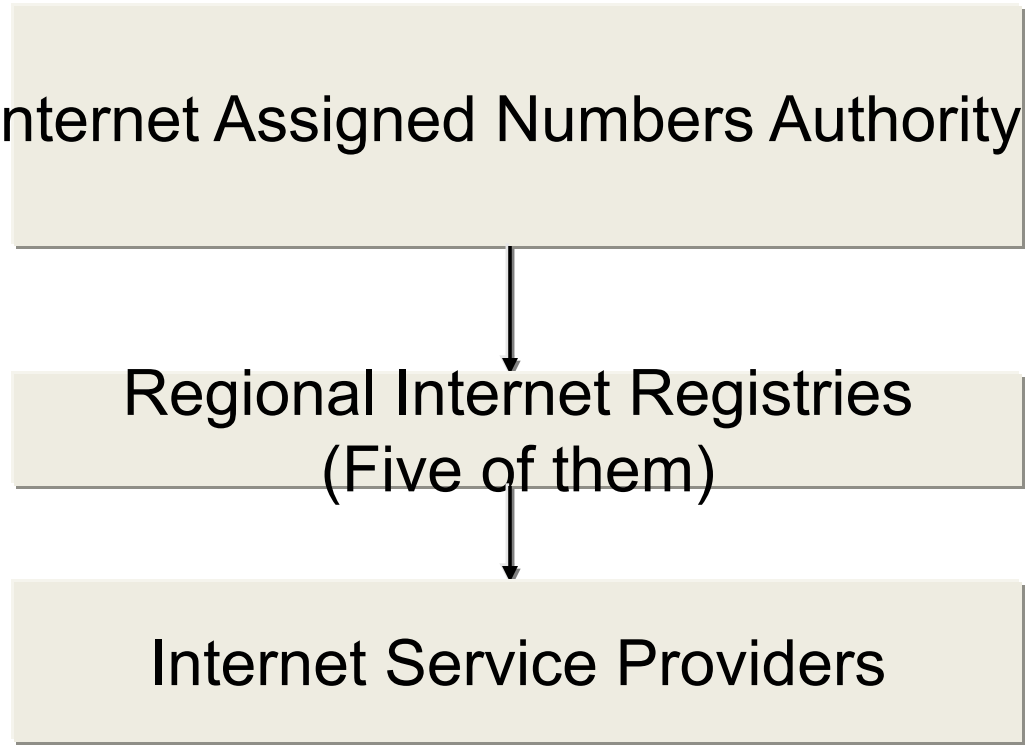


- Fast growing routing table size
  - Each router must have an entry for every network prefix
  - ~ 2^{21} = 2,097,152 class C networks
  - In 1993, the size of routing tables started to outgrow the capacity of routers
- Local admins must request another network number before installing a new network at their site

Solution: Classless Inter-domain routing (CIDR)

- Network prefix is of variable length
  - No rigid class boundary
- Addresses are allocated hierarchically
- Routers can aggregate multiple address prefixes into one routing entry
- Hierarchy is the key

Hierarchical IP Address Allocation



- American Registry for Internet Numbers (ARIN)
- RIPE, APNIC, LACNIC, AfriNIC

CIDR network prefix has variable length

	128	143	137	144
Addr	10000000	10001111	10001001	10010000
	255	255	255	0
Mask	11111111	11111111	1111111	00000000

- A network mask specifies the number of bits used to identify a network in an IP address.

CIDR notation

- CIDR notation of an IP address:
  - 128.143.137.144/24
  - /24 is the prefix length. It states that the first 24 bits are the network prefix of the address (and the remaining 8 bits are available for specific host addresses)

- CIDR notation can nicely express blocks of addresses
  - An address block

    [128.195.0.0, 128.195.255.255]

    can be represented by an address prefix

    128.195.0.0/16
  - How many IP addresses are there in a /x address block?
    - $2^{(32-x)}$

Output of ifconfig

```
en0: flags=8863<UP,BROADCAST,SMART,RUNNING,SIMPLEX,MULTICAST> mtu 1500
        ether a8:66:7f:16:02:08
        inet6 fe80::10cf:731b:1d54:e775%en0 prefixlen 64 secured scopeid 0x5
        inet 10.194.131.251 netmask 0xffffe000 broadcast 10.194.159.255
        nd6 options=201<PERFORMNUD,DAD>
        media: autoselect
        status: active
```

IP Forwarding

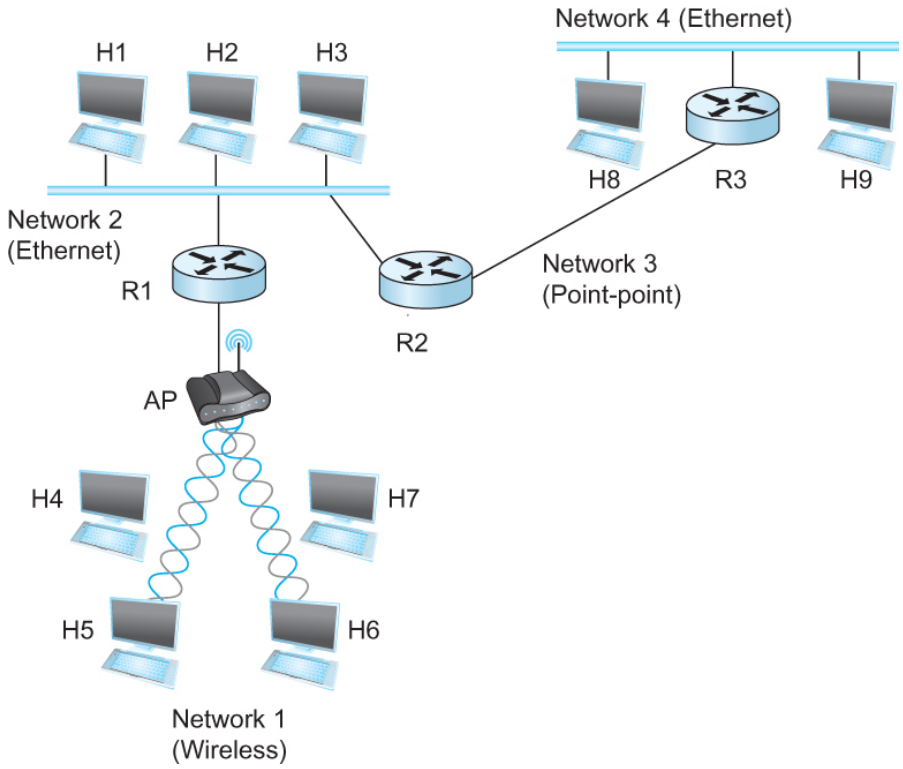
Forwarding of IP datagrams

- There are two distinct processes to delivering IP datagrams:

  1. **Forwarding (data plane):** How to pass a packet from an input interface to the output interface?

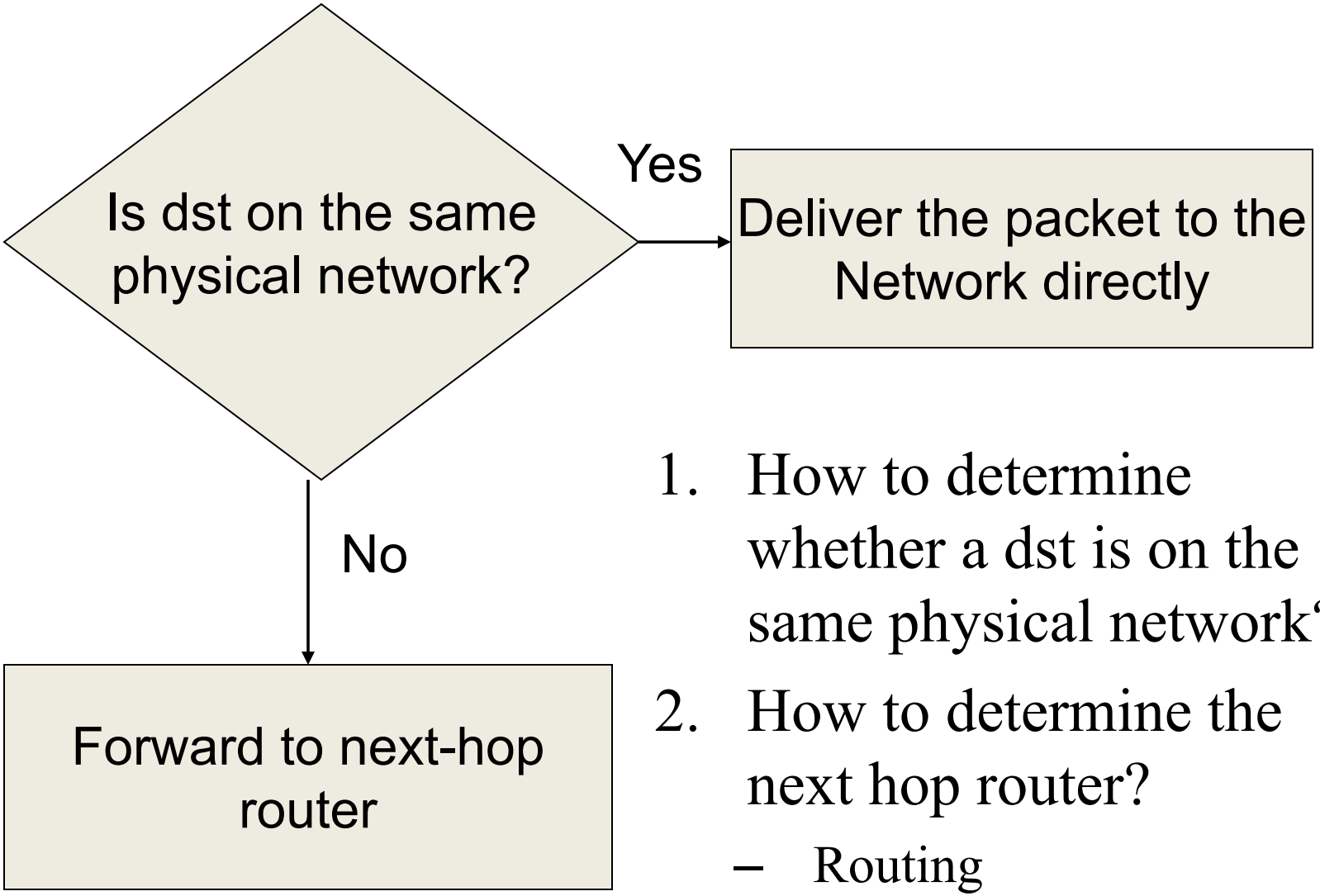
  2. **Routing (control plane):** How to find and setup the forwarding tables?

Forwarding basics



- Routers forward according to network prefixes
- All interfaces on the same network have the same network prefixes

Forwarding algorithm



1. How to determine whether a dst is on the same physical network?
2. How to determine the next hop router?
   - Routing

Detailed forwarding algorithm

- If (networkNum == networkNum of one of my interfaces) then
  - Deliver packet over the interface

- Else
  - if (NetworkNum is in my forwarding table) then
    - Deliver to the NextHop router
  - Else
    - Deliver packet to the default router

How does a host/router determine the network number of a destination address?

- Destination address & network mask = NetworkNumOfDestination

- If (NetworkNumOfDestination == my network Number) then
  - Send through my direct interfaces

Forwarding table lookup

- **Forwarding table lookup:** Use the IP destination address as a key to search the routing table
- Result of the lookup is the IP address of a next hop router, and/or the name of a network interface

Destination address	**Next hop/ interface**
network prefix *or* host IP address *or* loopback address *or* default route	IP address of next hop router *or* Name of a network interface

Type of forwarding table entries

- **Network route**
  - Destination addresses is a network address (e.g., 10.0.2.0/24)
  - Most entries are network routes
- **Host route**
  - Destination address is an interface address (e.g., 10.0.1.2/32)
  - Used to specify a separate route for certain hosts
- **Default route**
  - Used when no network or host route matches
- **Loopback address**
  - Routing table for the loopback address (127.0.0.1)
  - The next hop lists the loopback (lo0) interface as outgoing interface

Unified forwarding algorithm

- Observation:
  - A directly physical network can be an entry in the forwarding table
  - A default route can be an entry

1. Look up destination address in the forwarding table using longest prefix match
2. Forward the packet to the next hop indicated by the matched entry

The longest prefix matching algorithm

1. Search for a match on all 32 bits
2. Search for a match for 31 bits

   …..

32. Search for a match on 0 bits

Host route, loopback entry

→ 32-bit prefix match

Default route is represented as 0.0.0.0/0

→ 0-bit prefix match

Why longest prefix match?

- Longest → smallest network
- Network prefixes may be aggregated

Example

128.143.71.21

Destination address	Next hop
10.0.0.0/8	eth0
128.143.0.0/16	R2
128.143.64.0/20	R3
128.143.192.0/20	R3
128.143.71.0/24	R4
128.143.71.55/32	R3
0.0.0.0/0 (default)	R5

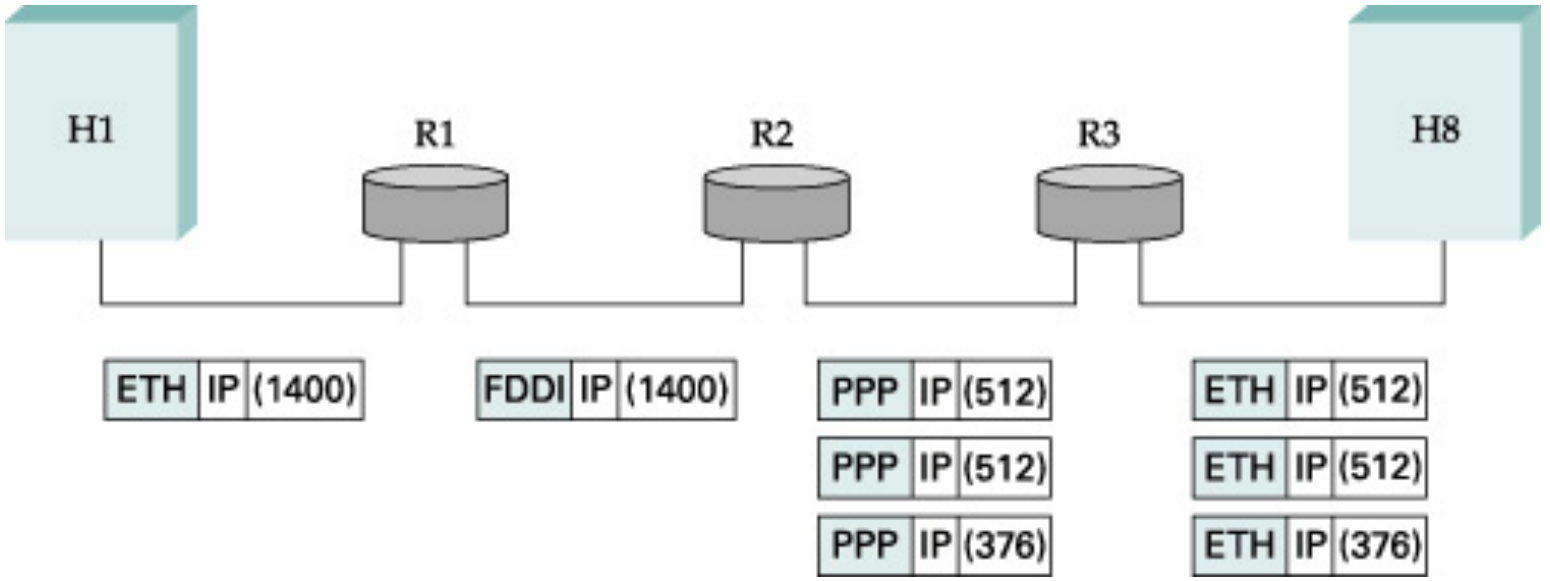
The longest prefix match for 128.143.71.21 is for 24 bits with entry 128.143.71.0/24

Datagram will be sent to R4

Fragmentation and Reassembly

(not required for Lab 2)

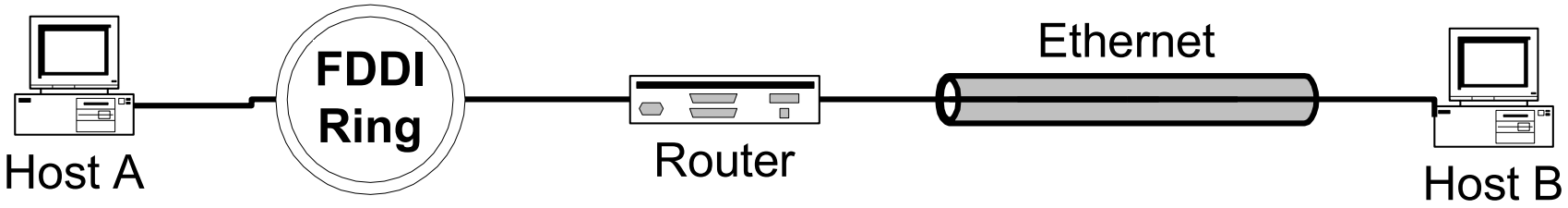
Different networks have different Maximum Transmission Units (MTUs)



IP Fragmentation and Reassembly

- What if the size of an IP datagram exceeds the MTU?
  IP datagram is fragmented into smaller units.

- What if the route contains networks with different MTUs?

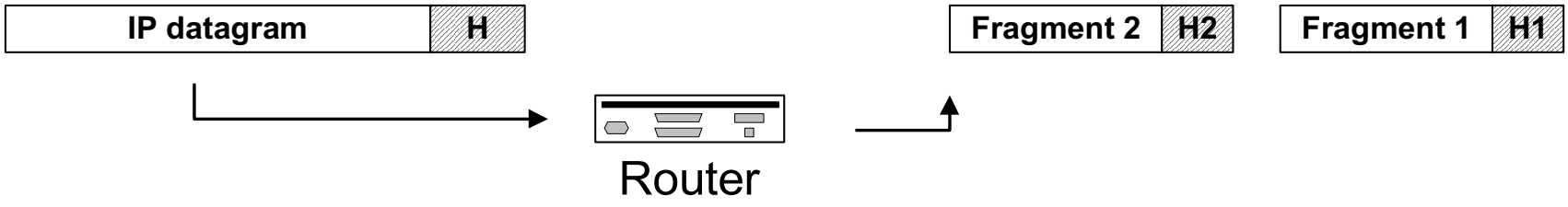


MTUs: FDDI: 4352 Ethernet: 1500

- **Fragmentation**:
  - IP router splits the datagram into several datagrams

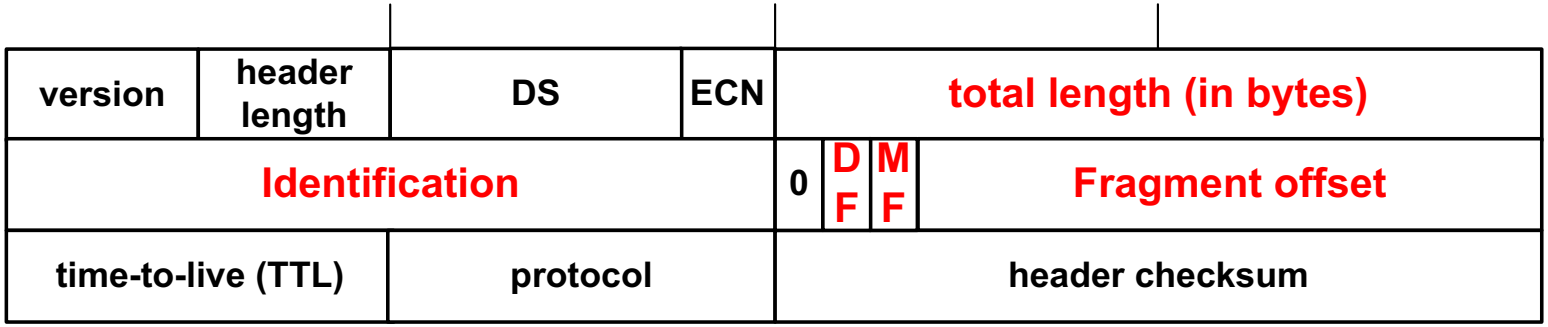
Design question: Where is Fragmentation/reassembly done?

- Fragmentation can be done at the sender or at intermediate routers
- The same datagram can be fragmented several times.
- Reassembly of original datagram is only done at destination hosts !! (why?)



What's involved in Fragmentation?

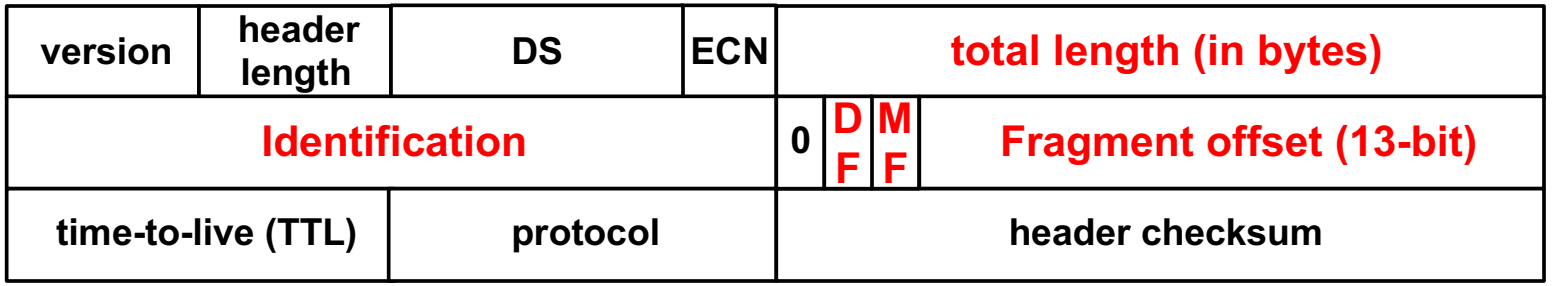
- The following fields in the IP header are involved:



- Identification
  - When a datagram is fragmented, the identification is the same in all fragments
  - Used to reassemble the original packet
- Flags
  - DF bit is set: datagram cannot be fragmented and must be discarded if MTU is too small
    - ICMP sent
  - MF bit:
    - 1: this is not the last fragment
    - 0: last fragment

What's involved in Fragmentation?

- The following fields in the IP header are involved:

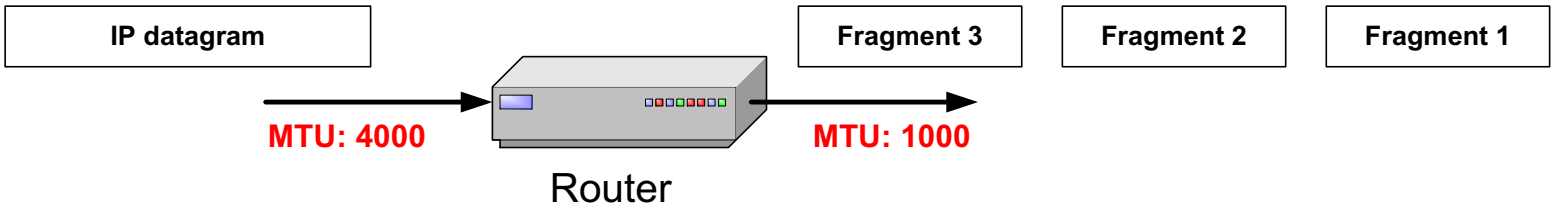


- *Fragment offset*
  - Offset of the payload of the current fragment in the original datagram in units of 8 bytes
    - Why?
    - Because the field is only 13 bits long, while the total length is 16 bits.
- Total length
  - Total length of the current fragment

Example of Fragmentation

- A datagram with size 2400 bytes must be fragmented according to an MTU limit of 1000 bytes

Header length: 20	Header length: 20	Header length: 20	Header length: 20
Total length: 2400	Total length: 448	Total length: 996	Total length: 996
Identification: 0xa428	Identification: 0xa428	Identification: 0xa428	Identification: 0xa428
DF flag: 0	DF flag: 0	DF flag: 0	DF flag: 0
MF flag: 0	MF flag: 0	MF flag: 1	MF flag: 1
Fragment offset: 0	Fragment offset: 244	Fragment offset: 122	fragment offset: 0



Determining the length of fragments

- Maximum payload length = 1000 – 20 = 980 bytes
- Offset specifies the bytes in multiple of 8 bytes. So the payload must be a multiple of 8 bytes.
- 980 - 980 % 8 = 976 (the largest number that is less than 980 and divisible by 8)
- The payload for the first fragment is 976 and has bytes 0 ~ 975 of the original IP datagram. The offset is 0.
- The payload for the second fragment is 976 and has bytes 976 ~ 1951 of the original IP datagram. The offset is 976 / 8 = 122.
- The pay load of the last fragment is 2400 – 976 * 2 = 428 bytes and has bytes 1952 ~ 2400 of the original IP datagram. The offset is 244.
- Total length of three fragments: 996 + 996 + 448 = 2440 > 2400
  - Why?
  - Two additional IP headers.

Path MTU discovery

- Fragmentation slows down the router
- → should be done by end hosts

- How does a sender know the MTU of a path?
  - A host only knows the MTU of its links

- Solution
  - send large packets with DF set
  - If receive ICMP Fragmentation needed messages, reduce maximum segment size

Summary

- History of IP
- IP header format
- IP addressing
- IP forwarding
  - Forwarding algorithm
  - Fragmentation