Compsci 201 Graphs, APTS, and More Part 1 of 5





Susan Rodger April 17, 2020

W is for ...

- World Wide Web
 - Invented in 1989 Sir Tim Berners Lee



• Wifi

• We need this everyday



Announcements

- APT-7 due Thursday, April 16
- APT-8 due Tuesday, April 21
- Assignment P6 Huffman due April 22
 - All late work turned in by April 22 (APTs and Asgns)
 - Except Huffman grace through April 23
- Assignment P7 Optional out Extra Credit!
- Exam 2 will focus on grading that the next few days
- APT Quiz 2 is April 12-18 Your own work!
- Final Exam will be on April 30 any time on this day
- Discussion 13, Monday, April 20, Last one!

Plan for the Day

- Logistics to Wrap up the Course
 - Extensions must end, Final Exam
- APTs and Algorithmic Concepts
 - Reminder about greedy
 - Memoizing
- Graphs and Graph Algorithms
 - Concepts, terminology, APTs
 - Toward Dijkstra's Algorithm

Last Chance Last Chance

- Extensions on APTs and Assignments
 - You must fill out the extension form
 - We can't take anything late after April 22!
 - Except P6 Huffman
 - only one grace day April 23!
 - NO LATE SUBMISSIONS!!!!!!
 - Except P7 Create
 - til Sunday April 26 with no penalty

Calculating your grade

- Discussion sections (6% of your grade)
 - Add up your total points, max points are 13 disc*4pt =52
 - We drop 8 points
 - Divide your total points by 44.
 - Examples:
 - 42pts/44 = 95%
 - 50pts/44 = 100% (can't get > 100)

Calculating your grade (part 2)

- Programming and Analysis (23% of your grade)
 - Max points are: 178pts
 - Divide your total by 178 for your score
- WOTOS (3.75% of your grade)
 - If you have 40% of the points you get 100%
- Reading Zybooks or 6 extra APTs (1.25% of grade)
- APTs (6%)
 - 38 APTs, 10 pts each 380+ points is 100%

Calculating your grade (part 3)

• 90 A-, 94 A

Discussion Sections	6%
Programming and analysis assignments	23%
WOTOs(75%)/Reading(25%)	5%
APTs	6%
APT Quizzes (2)	10%
Exam grade: Exam1, Exam2 and Final Exam	50%

Final Exam

- Final exam must take on Thursday, April 30
- Must take in a 24 hour window.
- Once you start have 3 hours plus 1 hour
 - That is 4 hours total
- About a 2 hour exam.
- Similar format to Exam 2
- If you don't take it get a 0.
- Exam grade is MAX of (exam1, exam2, final exam)

P7 Create assignments coming in ...



P7 Create assignments coming in ...



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imgflip.com

JAKE-CLARK.TUMBLR

Be A UTA https://www.cs.duke.edu/undergrad/uta

- SIGN UP NOW by April 29 or before...
- CompSci 201
- CompSci 101
 - Python
- CompSci 94
 - Programming with Alice, so easy to learn

CS Concepts Coming Alive

• What data structure is this?



Compsci 201, Spring 2020

2D-range tree

- Search in x-y plane
- Main tree organized by x-values
- Subtree organized by y values



Binary Search tree of points in the plane – sorted by X-value



Compsci 201 Graphs, APTS, and More Part 2 of 5





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On to Graphs: From Bacon to Erdös

• <u>https://mathscinet.ams.org/mathscinet/freeTools.html?version=2</u>





Bacon Number and Erdös Number

- Some actors are prolific: lots of movies
 - Chris in movie with Sam in movie with K. Bacon
 - Chris has a Bacon number of two
- Some authors are prolific: lots of papers/articles
 - Tina wrote paper with Tom wrote with P. Erdös
 - Tina has an Erdös number of two
- Graph terminology: connecting nodes with edges
 - In-movie-with or wrote-paper-with is edge

Erdös Numbers

Authors connected by authorship/paper-writing

https://mathscinet.ams.org/mathscinet/freeTools.html?version=2



Erdös Numbers

Authors connected by authorship/paper-writing

Search MSC	Collaboration Di	stance Current Jo	ournals Current	Publications
MR Erdos Number = 2				
Susanne	E. Hambrusch	coauthored with	Alfred J. Boals	MR1140478
Alfred J.	Boals	coauthored with	Paul Erdős ¹	MR0944683
Change First Author Change Second Author New Search				

Bacon Number

Actors connected by acting/movie-roles

https://oracleofbacon.org/





Obama in Movie



Widely Used Mobile Apps

- Google Maps, Uber, Lyft
 - Why is UI important for drivers
 - Why shortest path algorithms important
 - Which use graphs?



Why Graphs are Important

- Google/Pagerank models webpages as a graph
 - Nodes are webpages
 - Links are hyperlinks between pages
 - Weights based on "importance" of link/page
 - <u>https://en.wikipedia.org/wiki/PageRank</u>



Is the Internet a Graph? It depends

- Internet as graph
 - Nodes are anything with an IP address (IP)
 - Nodes are Autonomous Systems (AS)
 - Edges connect Thermostat to Website or ...



The Coronavirus graph you don't see

- Who infected who
 - Nodes people
 - Edges person A infected person B
 - Need a lot of testing to make this graph



You can see the growth is exponential, HUGE!

Graphs

- Graphs are collections of vertices and edges
 - Vertices or nodes, edges or links
 - Undirected graph Tom-Kevin and Meg-Kevin
 - Sometimes edges have weights



Directed (weighted) Graph

- Edges can have direction: directed graph
 - Not Facebook. Yes Tinder?



Data Structures for Graphs

- Use number for vertex, index into array
 - Can use string and map as well
- Adjacency List Representation
 - Good for sparse graphs





http://lagodiuk.github.io/computer science/2016/12/19/efficient adjacency lists in c.html

Adjacency Matrix

- Good for dense graphs, vertices still numbers
 - Symmetric matrix if undirected
 - Can have weights instead of 0,1



https://www.oreilly.com/library/view/php-7-data/9781786463890/32fd15e8-423f-49aa-84c2-db1518023299.xhtml

Theory and Practice

- Code is often simpler with Adjacency "list"
 - Map<String, Set<String>> for "list"
 - Vertex identified by String
 - Connected-by-edge? set of vertices
 - Need something more for weighted graphs
- For APTs, this is a good approach as we'll see
 - Simple to make, simple to use, scaling? meh

Compsci 201 Graphs, APTS, and More Part 3 of 5





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Simple Graph Algorithms

- What vertices are reachable from starting vertex?
 - Can use DFS or BFS to find connected vertices
 - Must avoid visiting same vertex more than once
- Find connected components
 - Many applications

https://en.wikipedia.org/wiki/Connected component (graph theory)

```
Breadth First Search
public Set<String> bfs(String start) {
    Set<String> visited = new TreeSet<>();
    Queue<String> qu = new LinkedList<>();
    visited.add(start);
    qu.add(start);
    while (qu.size() > 0) {
      String v = qu.remove();
      for(String adj : myGraph.getAdjacent(v)) {
        if (! visited.contains(adj)) {
            visited.add(adj)'
            qu.add(adj);
    return visited;
```

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BFS Example

- BFS at A: B, C, D
 - from B: E, from C: ..., from D: F
 - from E: ..., from F, ...



- BFS: all one-away
 - then all two, then all three, ...

BFS becomes DFS

```
public Set<String> dfs(String start){
    Set<String> visited = new TreeSet<>();
    Queue<String> qu = new LinkedList<>();
    visited.add(start);
    qu.add(start);
```

```
while (qu.size() > 0) {
  String v = qu.remove();
  for(String adj : myGraph.getAdjacent(v)){
    if (! visited.contains(adj)) {
        visited.add(adj);
        qu.add(adj);
return visited;
```

DFS arrives

```
public Set<String> dfs(String start) {
    Set<String> visited = new TreeSet<>();
    Stack<String> qu = new Stack<>();
    visited.add(start);
    qu.push(start);
    while (qu.size() > 0) {
      String v = qu.pop();
      for(String adj : myGraph.getAdjacent(v)) {
        if (! visited.contains(adj)) {
            visited.add(adj);
            qu.push(adj);
    }
```

return visited;

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DFS Example

- DFS at A: B, C, D
 - then F, E



• DFS: goes deep one at a time

Example: Word Ladder Problem

- Change a word into another word
 - Change one letter at a time
 - Change COLD to WARM
 COLD -> CORD -> WORD -> WORM -> WARM



Algorithms + Data Structures

- BFS + Graphs = Word Ladder or Bacon Number
 - Getting from "above" to "zeros" in 17 steps!
 - above abode anode anole anile anise arise prise prime prims prams prats peats heats heads herds heros zeros
 - These edge weights are 1, so BFS works

- We can find the shortest path efficiently
 - Dijkstra's algorithm used in Internet today
 - Heuristics augment, absolute shortest needed?

Shortest Path and Longest Path

- We use breadth first search to find shortest path
 - Same code we saw in word-ladder problem
 - White, While, Whale, Shale, ... House
 - Efficient and polynomial time: edge-weight == 1
 - Need Dijkstra for positive edge-weight, still good
- No efficient algorithm for longest path, it's hard
 - If one found, every hard problem becomes easy
 - Most computer scientists don't think we'll find one

Connected Components: APT

- <u>https://www2.cs.duke.edu/csed/newapt/internet.html</u>
- Given a graph, a set of connected vertices
 - Which are important aka articulation points
 - Remove one? disconnect graph
- In example, removing 2 means ...
 - Disconnect 3 from 0 and 1



Connected Components: APT

- https://www2.cs.duke.edu/csed/newapt/internet.html
- What is this problem asking you to do?
 - What router, if removed, disconnects others?
- This is a graph problem! Vertices and edges?
 - Parse input, build graph, traverse graph
- Adjacency List: Map<String, Set<String>>
 - map.get("2") -- set of connected vertices

Toward All Green

- What part of this haven't you seen?
 - How is DFS or BFS used? Modify based on ...

```
public int articulationPoints(String[] routers) {
6∘
           makeGraph(routers);
 7
           int total = 0;
8
9
           for(int k=0; k < routers.length; k++) {</pre>
               String vertex = ""+k;
10
               String start = "0";
11
               if (k == 0) start ="1";
12
               Set<String> set = reachFromSkip(start,vertex);
13
14
               if (set.size() != routers.length-1) {
15
                    total += 1;
               }
16
17
           3
           return total;
18
19
       }
```

What is reachFromSkip method?

- Use BFS or DFS as provided, but ...
 - Do not push or enqueue skippable vertex/node
 - Can we reach everything from start? good!
 - Start from "0" unless skipping "0", ...
- Must create graph from input

WOTO (4 minutes)

http://bit.ly/201spring20-0417-1



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Jon Kleinberg

- Developed HITS, same time-frame as PageRank
- Professor at Cornell University
- MacArthur Genius award, Nevanlinna Prize, more
 - "It's much easier to make progress on a problem when you are enjoying what you are doing. In addition to finding work that is important, find work that has some personal interest for you....
 <u>ACM Infosys Interview</u>



Compsci 201 Graphs, APTS, and More Part 4 of 5





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Greedy Algorithms

- Which candles to burn?
 - The tallest ones: leads to more burning days
- Which votes to steal?
 - Opponent with the most: fewer "steals" to win
- Which weighted nodes to join in Huffman coding?
 - Smallest weights first: save bits, optimal!

A friend of a friend: APT

- <u>https://www2.cs.duke.edu/csed/newapt/friendscore.html</u>
- Model as a graph? Vertex: number, Edge? == 'Y'
 - 0: has one friend: 1
 - 1:0 and 2
 - 2: 1 and 3
 - 3: 2 and 4
 - 4:3
- So 2 has four two-friends
 - 1 has three two-friends
 - 0 has two two-friends





General Framework to Solve

- How to write twoFriends?
 - Make graph, find two-friends via ...
 - Find 1 friends? index of each 'Y'. Repeat

```
Map<Integer,Set<Integer>> myGraph;
 4
 5
       public int highestScore(String[] friends) {
 6∘
 7
           makeGraph(friends);
           int max = 0;
8
           for(int k=0; k < friends.length; k++) {</pre>
9
                Set<Integer> set = twoFriends(k);
10
11
               max = Math.max(set.size(), max);
           }
12
13
           return max;
       }
14
```

Set.addAll --- all my friends

- Model data using graph: parse via makeGraph
 4 Map<Integer,Set<Integer>> myGraph;
- My friends: myGraph.get(my_number))
 - Friend of a friend? for each of my friends ...

19	<pre>for(int friend: myGraph.get(my_index))</pre>	{
20	<pre>set.addAll(myGraph.get(friend));</pre>	
21	}	

Compsci 201 Graphs, APTS, and More Part 5 of 5





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Mathematics and Computer Science

- How do we solve differential equations?
 - It depends
- How do we estimate percolation threshold?
 - It depends
- How do we model cardiac behavior? ...
- Use simulation when no analytic solutions
 - Monte Carlo simulation for many problems
 - https://en.wikipedia.org/wiki/Monte_Carlo_method

Thinking about math+compsci <u>https://www2.cs.duke.edu/csed/newapt/bstcount.html</u>

- How many different binary search trees are there?
 - Size = 4, Size = 5 ... Size = N?
 - What about N = 6?

Ν	# trees
0	1
1	1
2	2
3	5
4	14
5	42

Combinatorics and Catalan

- Binary search trees with 6 nodes
 - Left subtree has: 0,1,2,3,4,5 nodes
 - What will right subtree have?
 - For each left, there is a right...
 - Count how many ways this happens

Ν	# trees
0	1
1	1
2	2
3	5
4	14
5	42

(1*42)+(1*14)+(2*5)+(5*2)+(14*1)+(42*1)= 132 Verify via: <u>https://en.wikipedia.org/wiki/Catalan_number</u>

Aside: From Catalan to Fibonacci

- Read about the Golden Ratio and Fibonacci #'s
 - 1,1,2,3,5,8,13,21, ... it's about rabbits?
 - Inevitable we discuss this, factorial, Bubble sort
- Do not do this at home, see classwork on Git <u>https://coursework.cs.duke.edu/201spring20/classcode/</u>

15⊜	<pre>public static long rfibo(int n) {</pre>
16	<pre>if (n <= 2) return 1;</pre>
17	<pre>return rfibo(n-1) + rfibo(n-2);</pre>
18	}

Exponential number of calls

- Since fib(8) calls fib(7) and fib(6)
 - And fib(6) calls ... which calls ... which ...
 - What is the recurrence? ~ T(n) = 2T(n-1) + O(1)
 - Solution to this is $O(2^n)$
- Actual fib isn't 2ⁿ, is exponential
 - Golden ratio: φ^{n} $\lim_{n \to \infty} \frac{F_{n+1}}{F_{n}} = \varphi$ $\underset{fib(1) \ fib(0)}{fib(1) \ fib(1) \ fib(0) \ fib(2) \ fib(1) \ fib(1) \ fib(0) \ fib(1) \ fib(1) \ fib(0) \ fib(1) \ fib(1) \ fib(0) \ fib(1) \ fib(1)$

fib(5)

Memoize aka Caching

- Caching in computer science is ... store to re-use
 - Similar to dynamic programming, but top-down
- If already seen? use that result, no recursion
 - Otherwise, recurse, store, return

```
static long [] memo = new long[5000];
20
       public static long rfib(int n) {
21⊖
           if (n <= 2) return 1;</pre>
22
23
           if (memo[n] != 0) {
24
                return memo[n];
25
            }
26
           memo[n] = rfib(n-1) + rfib(n-2);
27
           return memo[n];
       }
28
```

Look at this tree again

• Instead of doing this.....



Avoid repeated recursion ...

- Store calculated values in a map
 - Look up first, re-use what's already done
 - Use Map<Integer,Long> or long[]
 - An array is a map of index to value



All Green? Which one ...

- This solution will time out, too many helper calls
 - Use memoization to get all green
 - Add array or map, store, re-use

```
8∍
       public long helper(int n) {
 9
           if (n == 0 || n == 1) return 1;
           long total = 0;
10
           for(int leftCount = 0; leftCount < n; leftCount++) {</pre>
11
                total += helper(leftCount)*helper(n-leftCount-1);
12
           3
13
14
           return total;
       }
15
```

All Green? Do NOT turn this in

- Catalan via Wikipedia: this should NOT be used.
 - Notice **6564120420L**, long constant

```
public long howMany(int[] values) {
 30
           long[] catalan = {
 4
 5
               1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796,
 6
               58786, 208012, 742900, 2674440, 9694845, 35357670,
7
               129644790, 477638700, 1767263190, 6564120420L,
8
               24466267020L, 91482563640L, 343059613650L, 12899041473
           };
 9
           return catalan[values.length];
10
       }
11
```

WOTO

http://bit.ly/201spring20-0417-2

