Deterministic Finite Accepter (or Automata)

A DFA = \((Q, \Sigma, \delta, q_0, F)\)

where
- \(Q\) is finite set of states
- \(\Sigma\) is tape (input) alphabet
- \(q_0\) is initial state
- \(F \subseteq Q\) is set of final states.

\(\delta: Q \times \Sigma \rightarrow Q\)
Example: DFA that accepts even binary numbers.

Transition Diagram:

\[
M = (Q, \Sigma, \delta, q_0, F) = \\
(\{q_0, q_1\}, \{0, 1\}, \delta, q_0, \{q_1\})
\]

Tabular Format

<table>
<thead>
<tr>
<th></th>
<th>0</th>
<th>1</th>
</tr>
</thead>
<tbody>
<tr>
<td>q0</td>
<td>q1</td>
<td>q0</td>
</tr>
<tr>
<td>q1</td>
<td>q_1</td>
<td>q_0</td>
</tr>
</tbody>
</table>

Example of a move: \(\delta(q_0, 1) = q_0\)
Algorithm for DFA:

Start in start state with input on tape
q = current state
s = current symbol on tape
while (s ! = blank) do
    q = δ(q,s)
    s = next symbol to the right on tape
if q ∈ F then accept

Example of a trace: 11010
Pictorial Example of a trace for 100:

1) 1 0 0

2) 1 0 0

3) 1 0 0

4) 1 0 0
Definition:
\[ \delta^*(q, \lambda) = q \]
\[ \delta^*(q, wa) = \delta(\delta^*(q, w), a) \]

Definition The language accepted by a DFA \( M=(Q,\Sigma,\delta,q_0,F) \) is set of all strings on \( \Sigma \) accepted by \( M \). Formally,
\[ L(M) = \{ w \in \Sigma^* \mid \delta^*(q_0, w) \in F \} \]
Trap State

Example: \( L(M) = \{ b^n a \mid n > 0 \} \)

\( \Sigma = \{a, b\}\)
Example:

$L = \{ w \in \Sigma^* \mid w \text{ has an even number of a’s and an even number of b’s} \}$
Example: DFA that accepts even binary numbers that have an even number of 1’s.
Definition A language is regular iff there exists DFA $M$ s.t. $L=L(M)$. 
Chapter 2.2

Nondeterministic Finite Automata (or Accepter)

Definition

An NFA = (Q, Σ, δ, q₀, F)

where

Q is finite set of states
Σ is tape (input) alphabet
q₀ is initial state
F ⊆ Q is set of final states.

δ: Q × (Σ ∪ {λ}) → 2^Q
Example

Note: In this example $\delta(q_0, a) =$

$L =$
Example

$L = \{ (ab)^n \mid n > 0 \} \cup \{ a^n b \mid n > 0 \}$
Definition $q_j \in \delta^*(q_i, w)$ if and only if there is a walk from $q_i$ to $q_j$ labeled $w$.

Example From previous example:

$\delta^*(q_0, ab) =$

$\delta^*(q_0, aba) =$

Definition: For an NFA $M$,

$L(M) = \{ w \in \Sigma^* \mid \delta^*(q_0, w) \cap F \neq \emptyset \}$
2.3 NFA vs. DFA: Which is more powerful?

Example:
Theorem Given an NFA $M_N = (Q_N, \Sigma, \delta_N, q_0, F_N)$, then there exists a DFA $M_D = (Q_D, \Sigma, \delta_D, q_0, F_D)$ such that $L(M_N) = L(M_D)$.

Proof:

We need to define $M_D$ based on $M_N$.

$Q_D =$

$F_D =$

$\delta_D :$
Algorithm to construct $M_D$

1. start state is $\{q_0\} \cup \text{closure}(q_0)$

2. While can add an edge
   
   (a) Choose a state $A = \{q_i, q_j, \ldots q_k\}$ with missing edge for $a \in \Sigma$
   
   (b) Compute $B = \delta^*(q_i, a) \cup \delta^*(q_j, a) \cup \ldots \cup \delta^*(q_k, a)$
   
   (c) Add state $B$ if it doesn’t exist
   
   (d) add edge from $A$ to $B$ with label $a$

3. Identify final states

4. if $\lambda \in L(M_N)$ then make the start state final.
Example:
Properties and Proving - Problem 1

Consider the property
Replace_one_a_with_b or R1awb for short. If L is a regular, prove
R1awb(L) is regular.

The property R1awb applied to a language L replaces one a in each string with a b. If a string does not have an a, then the string is not in R1awb(L).

Example 1: Consider L={aaab, bbaa}
R1awb(L)=

Example 2: Consider \( \Sigma = \{a, b\} \), L = \{w \in \Sigma^* | w \) has an even number of a’s and an even number of b’s\}
R1awb(L)=

Proof:
Properties and Proving - Problem 2

Consider the property
Truncate_all_preceeding_b’s or
TruncPreb for short. If L is a regular, prove TruncPreb(L) is regular.

The property TruncPreb applied to a language L removes all preceeding b’s in each string. If a string does not have an preceeding b, then the string is the same in TruncPreb(L).

Example 1: Consider L = \{aaab, bbba\}

TruncPreb(L) =

Example 2: Consider L =
\{(bba)^n \mid n > 0\}

TruncPreb(L) =

Proof:
Minimizing Number of states in DFA

Why?

Algorithm

• Identify states that are indistinguishable
  These states form a new state

Definition Two states $p$ and $q$ are indistinguishable if for all $w \in \sum^*$

$$
\delta^*(q, w) \in F \Rightarrow \delta^*(p, w) \in F \\
\delta^*(p, w) \notin F \Rightarrow \delta^*(q, w) \notin F
$$

Definition Two states $p$ and $q$ are distinguishable if $\exists w \in \sum^*$ s.t.

$$
\delta^*(q, w) \in F \Rightarrow \delta^*(p, w) \notin F \text{ OR } \\
\delta^*(q, w) \notin F \Rightarrow \delta^*(p, w) \in F
$$
Example:
Example: