Section: LR Parsing

LR PARSING

LR(k) Parser

- bottom-up parser
- shift-reduce parser
- L means: reads input left to right
- R means: produces a rightmost derivation
- k - number of lookahead symbols

LR parsing process

- convert CFG to PDA
- Use the PDA and lookahead symbols
Convert CFG to PDA

The constructed NPDA:

• three states: s, q, f
  start in state s, assume z on stack

• all rewrite rules in state s, backwards
  rules pop rhs, then push lhs
  \((s,\text{lhs}) \in \delta(s,\lambda,\text{rhs})\)
  This is called a reduce operation.

• additional rules in s to recognize terminals
  For each \(x \in \Sigma, \ g \in \Gamma, \ (s,\text{xg}) \in \delta(s,\text{x},\text{g})\)
  This is called a shift operation.

• pop S from stack and move into state q

• pop z from stack, move into f, accept.
Example: Construct a PDA.

\[
S \rightarrow aSb \\
S \rightarrow b
\]
LR Parsing Actions

1. shift
   transfer the lookahead to the stack

2. reduce
   For $X \rightarrow w$, replace $w$ by $X$ on the stack

3. accept
   input string is in language

4. error
   input string is not in language

LR(1) Parse Table

- Columns:
  terminals, $\$ and variables

- Rows:
  state numbers: represent patterns in a derivation
LR(1) Parse Table Example

1) $S \rightarrow aSb$
2) $S \rightarrow b$

<table>
<thead>
<tr>
<th></th>
<th>a</th>
<th>b</th>
<th>$</th>
<th>S</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>s2</td>
<td>s3</td>
<td></td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td></td>
<td></td>
<td>acc</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>s2</td>
<td>s3</td>
<td></td>
<td>4</td>
</tr>
<tr>
<td>3</td>
<td></td>
<td>r2</td>
<td>r2</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td></td>
<td>s5</td>
<td></td>
<td></td>
</tr>
<tr>
<td>5</td>
<td></td>
<td>r1</td>
<td>r1</td>
<td></td>
</tr>
</tbody>
</table>

Definition of entries:

- $sN$ - shift terminal and move to state $N$
- $N$ - move to state $N$
- $rN$ - reduce by rule number $N$
- $acc$ - accept
- blank - error
state = 0
push(state)
read(symbol)
entry = T[state,symbol]
while entry.action ≠ accept do
    if entry.action == shift then
        push(symbol)
        state = entry.state
        push(state)
        read(symbol)
    else if entry.action == reduce then
        do 2*size_rhs times {pop()}
        state := top-of-stack()
        push(entry.rule.lhs)
        state = T[state,entry.rule.lhs]
        push(state)
    else if entry.action == blank then
        error
        entry = T[state, symbol]
end while
if symbol ≠ $ then error
Example:

Trace aabbb

5
b
3 4 4 5
b S S b
2 2 2 2 4 4
a a a a S S
2 2 2 2 2 2 2 2 1
a a a a a a a a a S
0 0 0 0 0 0 0 0 0 0
S: z z z z z z z z z z
L: a a b b b b b b b $ $
To construct the LR(1) parse table:

- Construct a dfa to model the top of the stack
- Using the dfa, construct an LR(1) parse table

To Construct the DFA

- Add $S' \to S$
- Place a marker “_” on the rhs $S' \to _S$
- Compute closure($S' \to _S$).

Def. of closure:

1. $\text{closure}(A \to v_{xy}) = \{A \to v_{xy}\}$ if $x$ is a terminal.

2. $\text{closure}(A \to v_{xy}) = \{A \to v_{xy}\}$
   $\cup (\text{closure}(x \to _w))$ for all $w$ if $x$ is a variable.
• The closure ($S' \rightarrow _S$) is state 0 and “unprocessed”.

• Repeat until all states have been processed
  – unproc = any unprocessed state
  – For each x that appears in $A \rightarrow ux_v$ do
    * Add a transition labeled “x” from state “unproc” to a new state with production $A \rightarrow ux_v$
    * The set of productions for the new state are: closure($A \rightarrow ux_v$)
    * If the new state is identical to another state, combine the states Otherwise, mark the new state as “unprocessed”

• Identify final states.
Example: Construct DFA

(0) $S' \rightarrow S$
(1) $S \rightarrow aSb$
(2) $S \rightarrow b$
Backtracking through the DFA

Consider aabbb

- Start in state 0.
- Shift “a” and move to state 2.
- Shift “a” and move to state 2.
- Shift “b” and move to state 3.
  Reduce by “S → b”
  Pop “b” and Backtrack to state 2.
  Shift “S” and move to state 4.
- Shift “b” and move to state 5.
  Reduce by “S → aSb”
  Pop “aSb” and Backtrack to state 2.
  Shift “S” and move to state 4.
- Shift “b” and move to state 5.
  Reduce by “S → aSb”
  Pop “aSb” and Backtrack to state 0.
Shift “S” and move to state 1.
• Accept. aabbb is in the language.
To construct LR(1) table from diagram:

1. If there is an arc from state1 to state2
   (a) arc labeled $x$ is terminal or $T[state1, x] = \text{sh state2}$
   (b) arc labeled $X$ is nonterminal $T[state1, X] = \text{state2}$

2. If state1 is a final state with $X \rightarrow w$
   For all $a$ in FOLLOW($X$), $T[state1, a] = \text{reduce by } X \rightarrow w$

3. If state1 is a final state with $S' \rightarrow \_S$
   $T[state1, \$] = \text{accept}$

4. All other entries are error
Example: LR(1) Parse Table

(0) S’ → S
(1) S → aSb
(2) S → b

Here is the LR(1) Parse Table with extra information about the stack contents of each state.

<table>
<thead>
<tr>
<th>Stack contents</th>
<th>State number</th>
<th>Terminals</th>
<th>Variables</th>
</tr>
</thead>
<tbody>
<tr>
<td>(empty)</td>
<td>0</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>1</td>
<td></td>
<td></td>
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<td></td>
<td>2</td>
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<tr>
<td></td>
<td>5</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Actions for entries in LR(1) Parse table \( T[\text{state}, \text{symbol}] \)

Let entry \( = T[\text{state}, \text{symbol}] \).

- If symbol is a terminal or $:
  - If entry is “shift state \(i\)” push lookahead and state \(i\) on the stack
  - If entry is “reduce by rule \(X \rightarrow w\)” pop \(w\) and \(k\) states (\(k\) is the size of \(w\)) from the stack.
  - If entry is “accept” Halt. The string is in the language.
  - If entry is “error” Halt. The string is not in the language.
• If symbol is nonterminal
  We have just reduced the rhs of a production $X \rightarrow w$ to a symbol. The entry is a state number, call it state$_i$. Push $T[\text{state}_i, X]$ on the stack.
Constructing Parse Tables for CFG’s with λ-rules

A → λ written as A → λ_

Example

S → ddX
X → aX
X → λ

Add a new start symbol and number the rules:

(0) S’ → S
(1) S → ddX
(2) X → aX
(3) X → λ

Construct the DFA:
Construct the LR(1) Parse Table

<table>
<thead>
<tr>
<th></th>
<th>a</th>
<th>d</th>
<th>$</th>
<th>S</th>
<th>X</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
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<td></td>
<td></td>
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<tr>
<td>1</td>
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<tr>
<td>6</td>
<td></td>
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<td></td>
</tr>
</tbody>
</table>
Possible Conflicts:

1. Shift/Reduce Conflict
   Example:
   
   \[
   A \rightarrow ab \\
   A \rightarrow abcd
   \]
   
   In the DFA:
   
   \[
   A \rightarrow ab_ \\
   A \rightarrow ab_ cd
   \]

2. Reduce/Reduce Conflict
   Example:
   
   \[
   A \rightarrow ab \\
   B \rightarrow ab
   \]
   
   In the DFA:
   
   \[
   A \rightarrow ab_ \\
   B \rightarrow ab_ \\
   \]

3. Shift/Shift Conflict