

CompSci 201, L5: Sets and Maps

Announcements, Coming up

- Today, Monday 1/30
 - Project 0: Person201 due
- This Wednesday, 2/1
 - APT2 due
- Next Monday, 2/6
 - Project 1: NBody due (future projects will be 2 week)

Person in CS: Shafi Goldwasser



- Born 1959 NYC, Israeli family
- Started grad school in CS at UC Berkeley without knowing what she wanted to study. PhD 1984.
- 2012 Turing award winner (and 2 Gödel prizes) along with Silvio Micali for theoretical computer science in the creation of zero-knowledge proofs in theoretical cryptography.
- Professor and Director of Simons Institute for Theory of Computing.

Wrapping up ArrayList: Analyzing Efficiency

Algorithmic tradeoffs depend on the implementation

Often, we are interested in how the **efficiency** of operations on data structures depends on **scale**. For an **ArrayList** with N values how efficient is...

- **get()**. Direct lookup in an Array. “Constant time” – does not depend on size of the list.
- **contains()**. Loops through Array calling `.equals()` at each element. Takes longer as list grows.
- **size()**. Returns value of an instance variable tracking size, does not depend on size of the list.
- **add()**. Depends.

How efficient is **ArrayList** add?

For an **ArrayList** with N values, 2 cases:

1. Space left – One Array assignment statement, *constant time*, does not depend on list size.
2. No space left – Copy entire list! Takes N array assignments!

How often are we in the second slow case? Depends on *how much we increase the Array size by in case 2.*

ArrayList Growth

Starting with a length 1 Array, if you add N elements one at a time and (when full) create a new Array that...

Is twice as large (geometric growth)

- Must copy at sizes:
 - 1, 2, 4, 8, 16, 32, ...
- Total values copied looks like:
 - $1+2+4+8+\dots+(N/4)+(N/2)$

Has 1 more position (arithmetic growth)

- Must copy at sizes:
 - 1, 2, 3, 4, ...
- Total values copied looks like:
 - $1+2+3+\dots+(N-2)+(N-1)$

Algebra to our rescue!

ArrayList Growth and Algebra

Geometric growth

$$1 + 2 + 4 + \cdots + (N/2)$$

$$= \sum_{i=0}^{\log_2 N - 1} 2^i$$

$$N - 1$$

Geometric series formula:

$$\sum_{i=0}^n r^i = \frac{1 - r^{n+1}}{1 - r}$$

Arithmetic growth

$$1 + 2 + 3 + \cdots + (N - 1)$$

$$= \sum_{i=1}^{N-1} i$$

$$= N(N - 1)/2$$

Arithmetic series formula:

$$\sum_{i=1}^n a_i = \left(\frac{n}{2}\right) (a_1 + a_n)$$

Math and Expectations in 201

- **Do not** expect you to formally derive closed form expressions / give proofs.
 - **Do** expect you to recognize:
 - $1 + 2 + 4 + \dots + N$ is **linear**, grows like $\approx N$.
 - $1 + 2 + 3 + \dots + N$ is **quadratic**, grows like $\approx N^2$.
 - Patterns like these show up again and again!
- Will make “like” more with asymptotic an

Will make “like” more formal
with asymptotic analysis

```
3      int n = 100;
4      int numIterations = 0;
5      for (int i=0; i<n; i++) {
6          for (int j=0; j<i; j++) {
7              numIterations += 1;
8          }
9      }
```

```
numIterations: 4950
n*(n-1)/2): 4950
```

Experiment to verify hypothesis

Live Coding



ArrayList add (to end) is (amortized) efficient

According to the Java 17 API documentation:

“The add operation runs in *amortized constant time*...” –
What does that mean?

- With geometric growth (e.g., double size of Array whenever out of space): Need a *linear* number of copies $\propto N$ copies to add N elements to ArrayList.
- The *average* number of copies per add is thus $\propto \frac{N}{N} = 1$, a *constant* that does not depend on N .

ArrayList add to the front is not efficient

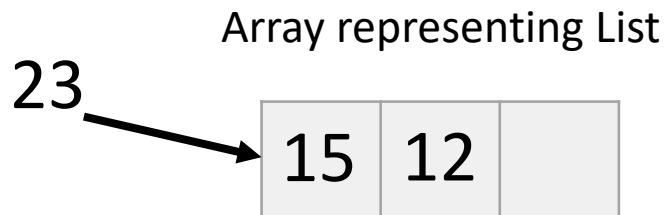
add

```
public void add(int index,  
                E element)
```

[Java 17 API documentation of add](#)

Inserts the specified element at the specified position in this list. Shifts the element currently at that position (if any) and any subsequent elements to the right (adds one to their indices).

Always requires shifting the entire Array, even if there is space available.



Sets

Set Review

```
public interface Set<E>  
    extends Collection<E>
```

A collection that contains no duplicate elements.

[Java API documentation](#)

- Stores UNIQUE elements
- Check if element in Set (using `.contains()`)
- Add element to set (using `.add()`)
 - Returns `false` if already there
- Remove element (with `.remove()`)
- Not guaranteed to store them in the order added

Set FAQs

```
[jshell> mySet  
mySet ==> [CS, 201]
```

1. How do I loop over a Set?

Enhanced for
loop

```
[jshell> for (String s : mySet) { System.out.println(s); }  
CS  
201
```

2. How do I convert between lists and sets?

```
[jshell> List<String> myList = new ArrayList<>();  
myList ==> []
```

```
[jshell> myList.addAll(mySet);  
$21 ==> true
```

addAll() method
convenient, same as looping
and adding one at a time

```
[jshell> myList  
myList ==> [CS, 201]
```

HashSet implementation of Set is very efficient

```
public class HashSet<E>  
    extends AbstractSet<E>  
    implements Set<E>, Cloneable, Serializable
```

Constant time = does not depend on the number of values stored in the Set.

This class implements the Set interface backed by a hash table (actually a HashMap instance). It makes no guarantees as to the iteration order of the set; in particular, it does not guarantee that the order will remain constant over time. This class permits the null element.

This class offers constant time performance for the basic operations (add, remove, contains and size), assuming the hash function disperses the elements properly among the buckets. Iterating over this set requires time proportional to the sum of the HashSet instance's size (the number of elements) plus the "capacity" of the backing HashMap instance (the number of buckets). Thus, it's very important not to set the initial capacity too high (or the load factor too low) if iteration performance is important.

[Java API documentation](#)

Under assumptions we will discuss next time

Count Unique Words?

```
public static int countWordsHashSet(String[] words) {  
    HashSet<String> mySet = new HashSet<>();  
    for (String w : words) {  
        mySet.add(w);  
    }  
    return mySet.size();  
}
```

For each word, constant time operation. "Linear complexity."

```
public static int countWordsArrayList(String[] words) {  
    ArrayList<String> myList = new ArrayList<>();  
    for (String w : words) {  
        if (!myList.contains(w)) {  
            myList.add(w);  
        }  
    }  
    return myList.size();  
}
```

For each word, must check all the words so far. "Quadratic complexity."

TreeSet stores sorted

Two important implementations of Set interface:

- HashSet – Very efficient add, contains
- TreeSet – Nearly as efficient, keeps values sorted.

```
5 String message = "computer science is so much fun";
6 char[] messageCharArray = message.toCharArray();
7 TreeSet<Character> uniqueChars = new TreeSet<>();
8 for (char c : messageCharArray) {
9     uniqueChars.add(c);
10 }
11 System.out.println(uniqueChars);
```

[, c, e, f, h, i, m, n, o, p, r, s, t, u]

Prints all unique
characters *in order*.

HashSet and TreeSet Implementations

```
public class HashSet<E>  
extends AbstractSet<E>  
implements Set<E>, Cloneable, Serializable
```

This class implements the `Set` interface, backed by a hash table (actually a `HashMap` instance). It makes no guarantees as to the iteration order of the set; in particular, it does not guarantee that the order will remain constant over time. This class permits the null element.

HashSet and HashMap both implemented with a hash table data structure, will discuss next time.

```
public class TreeSet<E>  
extends AbstractSet<E>  
implements NavigableSet<E>, Cloneable, Serializable
```

A `NavigableSet` implementation based on a `TreeMap`. The elements are ordered using their natural ordering, or by a `Comparator` provided at set creation time, depending on which constructor is used.

TreeSet and TreeMap both implemented using a special kind of *binary tree*, will discuss later in the course.

```
public class TreeMap<K,V>  
extends AbstractMap<K,V>  
implements NavigableMap<K,V>, Cloneable, Serializable
```

A Red-Black tree based `NavigableMap` implementation. The map

Maps

Map pairs keys with values

- Like an **address book**, lookup the value (address) of a key (person). Like a dictionary in Python.

Keys	Values
Bob	101 E. Main St.
Naomi	200 Broadway
Xi	121 Durham Ave.

- Map is an interface, must have methods like:
 - `put(k, v)`: Associate value `v` with key `k`
 - `get(k)`: Return the value associated with key `k`
 - `containsKey(k)`: Return true if key `k` is in the Map

Implementations: HashMap, TreeMap

```
1 import java.util.HashMap;  
2 import java.util.Map;  
3 import java.util.TreeMap;
```

Two major implementations:

- HashMap: Very efficient put, get, containsKey
- TreeMap: Nearly as efficient, keeps keys sorted

Map<KEY_TYPE, VALUE_TYPE>

Create a TreeMap to implement this Map

```
8 Map<String, String> addressBook = new TreeMap<>();  
9 addressBook.put("Bob", "101 E. Main St.");  
10 addressBook.put("Naomi", "200 Broadway");  
11 addressBook.put("Xi", "121 Durham Ave.");  
12 System.out.println(addressBook);
```

Sorted by keys due to
TreeMap

{Bob=101 E. Main St., Naomi=200 Broadway, Xi=121 Durham Ave.}

Check before you get

If you call `.get(key)` on a key not in the map, returns `null`, can cause program to crash.

```
6      Map<String, Integer> myMap = new HashMap<>();  
7      int val = myMap.get("hi");
```

Exception in thread "main" java.lang.NullPointerException: Cannot invoke "java.lang.Integer.intValue()" because the return value of "java.util.Map.get(Object)" is null

Instead, check first with `.containsKey()`.

```
6      Map<String, Integer> myMap = new HashMap<>();  
7      if (myMap.containsKey("hi")) {  
8          int val = myMap.get("hi");  
9      }
```

Adding “default” values

Often want a “default” value associated with new keys (examples: 0, empty list, etc.). Two options:

- `.putIfAbsent(key, val)`
- Check if does not contain key

```
6      Map<String, Integer> myMap = new HashMap<>();
7
8      myMap.putIfAbsent("hi", 0);
9
10     // Equivalent to line 8
11     if (!myMap.containsKey("hi")) {
12         myMap.put("hi", 0);
13     }
```

Updating maps

Single values

- `.get()` returns a *copy of the value*.
- Must use `.put()` again to update.

Collection values

- `.get()` returns *reference to collection*.
- Update the collection directly.

```
8   Map<String, Integer> myMap = new HashMap<>();
9   myMap.put("hi", 0);
10  int currentVal = myMap.get("hi");
11  myMap.put("hi", currentVal + 1);
```

```
14  Map<String, List<Integer>> otherMap = new HashMap<>();
15  otherMap.put("hi", new ArrayList<>());
16  otherMap.get("hi").add(0);
--
```

Counting with a Map

In this example we count how many of each character occur in message.

```
5 String message = "computer science is so much fun";
6 char[] messageCharArray = message.toCharArray();
7 TreeMap<Character, Integer> charCounts = new TreeMap<>();
8 for (char c : messageCharArray) {
9     if (!charCounts.containsKey(c)) {
10         charCounts.put(c, 1);
11     }
12     else {
13         int currentVal = charCounts.get(c);
14         charCounts.put(c, currentVal + 1);
15     }
16 }
17 System.out.println(charCounts);
```

Check if we have *not* seen c yet

Else get current value and increase

Comes in order because using TreeMap

{ =5, c=4, e=3, f=1, h=1, i=2, m=2, n=2, o=2, p=1, r=1, s=3, t=1, u=3}

Problem-Solving with Sets and Maps

Word Pattern Problem

Live Coding



<https://leetcode.com/problems/word-pattern/submissions/886368133/>