

L6: Hashing, HashMap, HashSet

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CompSci 201: Spring 2024
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Announcements, Coming up

- Today, Wednesday, 1/31
 - APT 2 due
- Monday, 2/5
 - Project 1: NBody due (future projects will be 2 week)
 - Project 2: Markov out (due in 2 weeks)
- Next Wednesday, 2/8
 - APT 3 due

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Finishing Maps API



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Map pairs keys with values

- Like an **address book**, lookup the value (address) of a key (person). Like a dictionary in Python.

| Keys | Values |
|---------|-----------------|
| Bob | 101 E. Main St. |
| Naomi | 200 Broadway |
| Stavros | 121 Durham Ave. |

- Map is an interface, must have methods like:
 - `put(k, v)`: Associate value `v` with key `k`
 - `get(k)`: Return the value associated with key `k`
 - `containsKey(k)`: Return true if key `k` is in the Map

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Implementations of Map

Two major implementations:

```
1 import java.util.HashMap;
2 import java.util.Map;
3 import java.util.TreeMap;
```

- HashMap**: Very efficient `put`, `get`, `containsKey`
- TreeMap**: Nearly as efficient, keeps **keys** sorted by their *"natural ordering"*

```
Map<KEY_TYPE, VALUE_TYPE>

8 Map<String, String> addressBook = new TreeMap<>();
9 addressBook.put("Bob", "101 E. Main St.");
10 addressBook.put("Naomi", "200 Broadway");
11 addressBook.put("Xi", "121 Durham Ave.");
12 System.out.println(addressBook);

Create a TreeMap to implement this Map

Sorted by keys due to TreeMap

{Bob=101 E. Main St., Naomi=200 Broadway, Xi=121 Durham Ave.}
```

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Check before you get

If you call `.get(key)` on a key not in the map, returns `null`, can cause program to crash.

```
6 Map<String, Integer> myMap = new HashMap<>();
7 int val = myMap.get("hi");

Exception in thread "main" java.lang.NullPointerException: Cannot invoke "java.lang.Integer.intValue()" because the return value of "java.util.Map.get(Object)" is null
```

Instead, check first with `.containsKey()`.

```
6 Map<String, Integer> myMap = new HashMap<>();
7 if (myMap.containsKey("hi")) {
8     int val = myMap.get("hi");
9 }
```

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Adding "default" values

Often want a "default" value associated with new keys (examples: 0, empty list, etc.). Two options:

- `.putIfAbsent(key, val)`
- Check if does not contain key before put

```

6      Map<String, Integer> myMap = new HashMap<>();
7
8      myMap.putIfAbsent("hi", 0);
9
10     // Equivalent to line 8
11     if (!myMap.containsKey("hi")) {
12         myMap.put("hi", 0);
13     }

```

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Updating maps

Immutable values:

- `.get()` returns a *copy* of the value.
- Must use `.put()` again to update.

Mutable values (e.g. collections)

- `.get()` returns *reference to collection*.
- Update the collection directly.

```

8      Map<String, Integer> myMap = new HashMap<>();
9      myMap.put("hi", 0);
10     int currentVal = myMap.get("hi");
11     myMap.put("hi", currentVal + 1);

14     Map<String, List<Integer>> otherMap = new HashMap<>();
15     otherMap.put("hi", new ArrayList<>());
16     otherMap.get("hi").add(0);

```

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Counting with a Map

In this example we count how many of each character occur in `message`.

```

5      String message = "computer science is so much fun";
6      char[] messageCharArray = message.toCharArray();
7      TreeMap<Character, Integer> charCounts = new TreeMap<>();
8      for (char c : messageCharArray) {
9          if (!charCounts.containsKey(c)) {
10             charCounts.put(c, 1);
11         }
12         else {
13             int currentVal = charCounts.get(c);
14             charCounts.put(c, currentVal + 1);
15         }
16     }
17     System.out.println(charCounts);

```

Check if we have not seen c yet

Else get current value and increase

Comes in order because using TreeMap

`{ =5, c=4, e=3, f=1, h=1, i=2, m=2, n=2, o=2, p=1, r=1, s=3, t=1, u=3}`

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HashSet/HashMap Implementation

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HashSet/Map efficiency

```
public class HashSet<E>
    extends AbstractSet<E>
    implements Set<E>, Cloneable, Serializable
```

This class implements the Set interface backed by a hash table (actually a HashMap instance). It makes no guarantees as to the iteration order of the set; in particular, it does not guarantee that the order will remain constant over time. This class permits the null element.

This class offers constant time performance for the basic operations (add, remove, contains and size), assuming the hash function disperses the elements properly among the buckets. Iterating over this set requires time proportional to the sum of the HashSet instance's size (the number of elements) plus the "capacity" of the backing HashMap instance (the number of buckets). Thus, it's very important not to set the initial capacity too high (for the load factor too low) if iteration performance is important.

[Java API documentation](#)

Constant time = does not depend on the number of values stored in the Set.

Under assumptions...

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Aside: Does constant time lookup (contains(), get(), etc.) matter?

- Social media: When you login, server needs to lookup to display the correct page for you.
 - Billions of accounts! Look it up in a List? NO! Constant time lookup with hashing.



- Routing/directions application: Need to lookup roads from a given intersection.
 - How many possible roads? Search through a list? NO! Constant time lookup with hashing.



- Could go on!

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Big questions about hashing

Last class: Usage of API HashSet/HashMap.

Today:

1. How does a hash table work to implement HashMap/HashSet?
2. Why do `.equals()` and `.hashCode()` matter?
3. Why are the `add()`, `contains()`, `put()`, `get()`, and `containsKey()`, etc., all constant time (and under what assumptions)?

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Hash Table Concept

- Implement HashMap with an Array also, storing `<key, value>` pairs
 - HashSet: A HashMap with only keys (no vals)
- Instead of always adding to next open spot (0, 1, 2, 3...)...
- **Big idea:** Calculate *hash* (an int) of key to determine where to store & lookup
 - Java OOP: Will use the `hashCode()` method of the key to get the hash
- Same hash to put and get, no looping over list

| | |
|---|-----------|
| 0 | |
| 1 | <"hi", 5> |
| 2 | |
| 3 | |
| 4 | <"ok", 3> |
| 5 | |
| 6 | |
| 7 | |

hash("ok") == 4

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HashMap methods at a high level

Always start by getting the **hash** =
`Math.abs(key.hashCode()) % list.size()`

Absolute value and % (remainder when dividing by) list size ensures valid index

- **put(key, value)**
 - Add (`<key, value>`) to list at index hash
 - If key already there, update value
- **get(key)**
 - Return value paired with key at index hash position of list
- **containsKey(key)**
 - Check if key exists at index hash position of list

| | |
|---|-----------|
| 0 | |
| 1 | <"hi", 5> |
| 2 | |
| 3 | |
| 4 | <"ok", 3> |
| 5 | |
| 6 | |
| 7 | |

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HashMap put/get example

- Suppose we have the <key, value> pair <"cs", 201>.

Hash:

```
[jshell] Math.abs("cs".hashCode()) % 8
$7 ==> 0
```

| | |
|---|-------------|
| 0 | <"cs", 201> |
| 1 | <"hi", 5> |
| 2 | |
| 3 | |
| 4 | <"ok", 3> |
| 5 | |
| 6 | |
| 7 | |

- put("cs", 201) in position 0
- get("cs") by looking up position 0, returns 201 returning the value

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Collisions

- Suppose now we want to put <"fain", 104>.

Hash:

```
[jshell] Math.abs("fain".hashCode()) % 8
$11 ==> 0
```

| | |
|---|-------------|
| 0 | <"cs", 201> |
| 1 | <"hi", 5> |
| 2 | |
| 3 | |
| 4 | <"ok", 3> |
| 5 | |
| 6 | |
| 7 | |

- put("fain", 104) in position 0
- But <"cs", 201> is already stored at position 0! Call this a **collision**.

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Dealing with collisions: concepts

- Think of the hash table as an Array of "buckets".
- Each bucket can store multiple <key, value> pairs.
- put(key, value)
 - Add to hash index bucket
 - Update value if key already in bucket
- get(key)
 - Loop over keys in hash index bucket
 - Return value of one that equals() key

| | |
|---|------------------------------|
| 0 | <"cs", 201> <"fain", 104> |
| 1 | <"hi", 5> |
| 2 | |
| 3 | |
| 4 | <"ok", 3> |
| 5 | |
| 6 | |
| 7 | |

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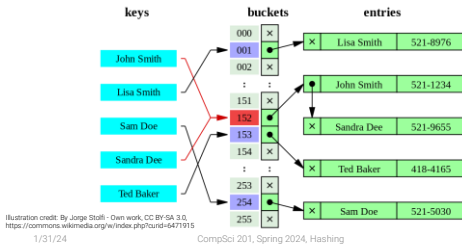
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Dealing with collisions: details

- Bucket is really another list
- Hash table is really an **array of lists** of <key, value> pairs.
- We call this technique for dealing with collisions **chaining**.



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2. HashSet and HashMap have constant time add, contains, put, get, and containsKey operations. That means that these methods... *

- ☐ Take the same amount of time to run
- ☐ Have the same number of operations
- ☒ Runtimes do not depend on number of elements of the Set/Map

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3. What is stored in each "bucket" in a HashMap? *

- ☐ A <key, value> pair
- ☐ A list of keys
- ☐ A list of values
- ☒ A list of <key, value> pairs

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4. Suppose we want to put <s, 1> into a HashMap where s.hashCode() = 12. If hash table has 4 buckets, in which bucket will we store <s, 1>? *

Select your answer **Answer: 0** ✓

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Where does equals() come in?

- If multiple <key, value> pairs in same bucket, need to know which to get() or update on a put() call.
- Always the pair where the key in the bucket equals() the key we put() or get().
- Need equals() to work correctly for the key type
 - String keys? Integer? Already implemented for you.
 - Storing objects of a class *you write*? Need to override and implement equals().

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What happens without equals()? Hashing cats

```

4 public class Cat {
5     String name;
6     int age;
7
8     @Override
9     public int hashCode() {
10         return 0;
11     }
12
13     Run | Debug
14     public static void main(String[] args) {
15         Set<Cat> myCats = new HashSet<>();
16         myCats.add(new Cat("kirk", 2));
17         myCats.add(new Cat("kirk", 2));
18         System.out.println(myCats.size());
19     }

```

Even though all cat objects have the same hashCode() of 0 and so go to the same bucket.

And these 2 Cat objects have the same values

Prints 2, cannot detect duplicates without equals()

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hashCode Correctness

- Need `hashCode()` to work correctly for the key type.
 - String keys? Already implemented for you.
 - Storing objects of classes *you write*? Need to override and implement `hashCode()`.
- What makes a `hashCode()` "correct" (not necessarily efficient)?
 - Any two objects that are `equals()` should have the same `hashCode()`.

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What happens without `hashCode()`? Hashing more cats

```

4 public class Cat {
5     String name;
6     int age;
7
8     @Override
9     public boolean equals(Object o) {
10         Cat other = (Cat) o;
11         if ((other.name.equals(this.name)) && (other.age == this.age)) {
12             return true;
13         }
14         return false;
15     }
16
17     Run | Debug
18     public static void main(String[] args) {
19         Set<Cat> myCats = new HashSet<>();
20         myCats.add(new Cat("kirk", 2));
21         myCats.add(new Cat("kirk", 2));
22         System.out.println(myCats.size());
23     }

```

Fixed `equals()` but removed `hashCode()`, using default

Still prints 2!

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Cat with `equals()` and `hashCode()`

```

4 public class Cat {
5     String name;
6     int age;
7
8     @Override
9     public boolean equals(Object o) {
10         Cat other = (Cat) o;
11         if ((other.name.equals(this.name)) && (other.age == this.age)) {
12             return true;
13         }
14         return false;
15     }
16
17     @Override
18     public int hashCode() {
19         return (name + Integer.toString(age)).hashCode();
20     }
21 }

```

`equals()` if have same name and age

Uses `String hashCode()` of name concat with age, if `equals()` will have same `hashCode()`

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Aside: toString()

Don't need for hashing, but `toString()` method allows "nice" printing.

```

4 public class Cat {
5     String name;
6     int age;
7
8     @Override
9     public String toString() {
10         return name;
11     }
12
13     Run | Debug
14     public static void main(String[] args) {
15         Set<Cat> myCats = new HashSet<>();
16         myCats.add(new Cat("kirk", 2));
17         System.out.println(myCats);
18     }

```

`toString()` method used for printing, including inside a Collection

Prints kirk instead of Cat@...

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What is the String hashCode()?

```

42 private int getBucket(String s) {
43     int val = Math.abs(s.hashCode()) % myTable.size();
44     return val;
45 }

```

Remember how `hashCode()` is used to get the bucket index.

```

hashCode [jshell] > "hello".hashCode();
public int hashCode() $4 ==> 99162322
Returns a hash code for this String. The hash code for a String object is (jshell) > "hello".hashCode();
s[0]*31^(n-1) + s[1]*31^(n-2) + ... + s[n-1] $5 ==> 99162323
using int arithmetic, where s[i] is the ith character of the string, n is the length of the string, and ^ indicates exponentiation. (The hash value of an empty string is 0.) (jshell) > "what".hashCode();
Overrides: hashCode in class Object $6 ==> 3648196
Returns: a hash code value for this object.
Java API String documentation
Interprets each character as an int, does arithmetic.

```

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Revisiting Hashing Efficiency

- Runtime of `get()`, `put()`, and `containsKey()`
= Time to get the hash

+ Time to search "bucket", calling `.equals()` on everything in the bucket

Constant, does not depend on number of pairs in Map

Depends on number of pairs per bucket

⇒ HashMaps are faster with more buckets

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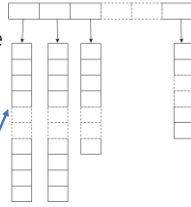
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"correct" but inefficient hashCode()

Correctness requirement: Any
.equals() keys should have the
same hashCode().

```
28  @Override
29  public int hashCode() {
30      return 0;
31  }
```



Still satisfies, but not good...

Stores everything in the first bucket!
No more efficient than a list!

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Correct and efficient hashCode()

From the [Java 17 API documentation](#):

- **Correctness:** "If two objects are equal...hashCode...must produce the same integer result."
- **Efficiency:** "...producing distinct integer results for unequal objects may improve the performance of hash tables."

- [String hashCode\(\)](#)
satisfies both

```
[jshell] > "hello".hashCode();
$4 ==> 99162322
[jshell] > "hellp".hashCode();
$5 ==> 99162323
[jshell] > "what".hashCode();
$6 ==> 3648196
```

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Cat hashCode() revisited

```
4  public class Cat {
5      String name;
6      int age;
7
8      @Override
9      public boolean equals(Object o) {
10         Cat other = (Cat) o;
11         if ((other.name == this.name) && (other.age == this.age)) {
12             return true;
13         }
14         return false;
15     }
16
17     @Override
18     public int hashCode() {
19         return (name + Integer.toString(age)).hashCode();
20     }
21 }
```

equals() if have same
name and age

If equals() will have
same hashCode()

If unequal? Unlikely (but
possible!) to have the
same hashCode()

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Simple uniform hashing assumption (SUHA)

- Suppose we hash N pairs to M buckets.
- **Simple uniform hashing assumption:**
Any element (i.e., key for HashMap, value for HashSet) is **equally likely** to hash into any bucket, **independently** of where any other element hashes to. [CLRS]
- Probability any two unequal elements hash into the same bucket: $1/M$
 - Spread of pairs to buckets **looks random** (but is not).
 - Ways to design such hash functions, not today
 - We make this assumption to analyze efficiency in theory, can verify runtime performance in practice

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Implications of SUHA

- Expected number of pairs per bucket under SUHA? N/M [N pairs, M buckets].
- Stronger statements are true: Very high probability that any bucket has approximately N/M pairs.
- Runtime implication?
 - Time to get the hash
 - Time to search over the hash index "bucket"
 - Calling `.equals()` on everything in the bucket

Constant, does not depend on N or M .Expect $\sim N/M$ pairs to search

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