

# L8: Asymptotic (Big-O) Analysis

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CompSci 201: Spring 2024  
5/8/24

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## Logistics, Coming up

- Today, 2/7
  - APT 3 due
- Next Monday 2/12
  - Midterm Exam 1
- Next Wednesday 2/14
  - APT 4 due

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## Person in CS: Alan Turing

- 1912-1954 (died at 41)
- English, PhD at Princeton in 1938
- Mathematician, cryptographer, pioneering thinker in AI
  - "Father of modern computer science"
  - Turing machine – helped formalize what is computable
  - Cryptography work in WW2
- Prosecuted in 1952 for homosexuality
  - Given choice of chemical "treatment" or prison, took former
  - Died 2 years later of cyanide poisoning, circumstances debated



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How many **total characters must be copied** by the code **on lines 8 and 9**? Remember that Strings are immutable in Java. \*

```
7 String s = "hi";
8 s += "hey";
9 s += s;
```

- ☐ 5
- ☐ 7
- ☐ 9
- ☐ 10
- ☒ 15
- ☐ 30

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Suppose method A has linear complexity and takes 10 ms to run on an input of size N. About what would you expect the runtime to be for an input of size 2\*N? \*

- ☐ 10 ms
- ☒ 20 ms
- ☐ 40 ms
- ☐ 100 ms

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Suppose method B has quadratic complexity and takes 10 ms to run on an input of size N. About what would you expect the runtime to be for an input of size 2\*N? \*

- ☐ 10 ms
- ☐ 20 ms
- ☒ 40 ms
- ☐ 100 ms

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Here is another String concatenation method. Suppose the input string  $s$  has a small number of characters, say 3. As a function of the parameter  $reps$ , how would you characterize the runtime complexity of the method? Hint: As a function of  $reps$ , how many total characters will be copied across all iterations of the loop? \*

```

7   public static String concatAloT(int reps, Str
8   for (int i=0; i<reps; i++) {
9       s += s;
10  }
11  return s;
12  }
```

- ☐ Constant  
☐ Linear  
☐ Quadratic  
☒ Exponential

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## Asymptotic Analysis and Big O Notation

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## Runtime and memory

- Two most fundamental resources on a computer:
  - Processor cycles: Number of operations per second machine can perform
    - (2 GHz = 2 billion operations per second).
  - Memory: space for storing variables, data, etc.
    - (esp. working memory, a.k.a. cache and RAM)
- We will mostly focus on runtime complexity
  - Often comes at expense of memory, e.g., HashMap
- Start by reasoning about empirical runtimes, but...

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## Constant Time Examples

- Index into an array (ar[0] or ar[201])
- Arithmetic (+, -, \*, /, %, etc.)
- Primitive comparison <, ==, etc.
- Accessing an object attribute (e.g. .length)
- ArrayList .get(), .size(), .add() [to end, amortized]
- Non-constant time usually has a loop or method call, may depend on implementation of data structures at hand

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## Big-O (limit definition)

- Given  $N$  (for example, the size of the input)
- Function  $T(N)$  (for example, the number of **constant time** operations in the code)

**Definition (big O notation).**

$T(N)$  is  $O(g(N))$  if  $\lim_{N \rightarrow \infty} \frac{T(N)}{g(N)} \leq c$  for some constant  $c$  that does not depend on  $N$ .

In other words:  $T(N)$  is  $O(g(N))$  if it is at most a constant factor times slower than  $g(N)$  for large input  $N$ .

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## Two general rules

1. Can drop constants
  - $2N+3 \rightarrow O(N)$
  - $0.001N + 1,000,000 \rightarrow O(N)$
2. Can drop lower order terms
  - $2N^2+3N \rightarrow O(N^2)$
  - $N+\log(N) \rightarrow O(N)$
  - $2^N + N^2 \rightarrow O(2^N)$

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Hierarchy of some common complexity class

Big O	Name	Example
$O(2^N)$	Exponential	Calculate all subsets of a set
$O(N^3)$	Cubic	Multiply NxN matrices
$O(N^2)$	Quadratic	Loop over all <i>pairs</i> from N things
$O(N \log(N))$	Nearly-linear	Sorting algorithms
$O(N)$	Linear	Loop over N things
$O(\log(N))$	Logarithmic	Binary search a sorted list
$O(1)$	Constant	Addition, array access, etc.

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Some common complexity classes and their growth

N	$O(\log(N))$	$O(N)$	$O(N^2)$	$O(N^3)$	$O(2^N)$	If you double N...
1	1	1	1	1	2	• $O(\log(N))$ adds ~1
2	2	2	4	8	4	• $O(N)$ roughly doubles
4	3	4	16	64	16	• $O(N^2)$ roughly quadruples
8	4	8	64	512	256	• $O(N^3)$ roughly multiplies by 8
16	5	16	256	4k	65k	• $O(2^N)$ squares each time
32	6	32	1k	32k	4.2E+9	
64	7	64	4k	262k	1.8E+19	

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Relation to Empirical Timing and Lower Order Terms

N	$N^2 + 19N + 200$	factor increase	
10	490	NA	Looks linear?
20	980	2.00	
40	2560	2.61	Asymptotic analysis describe behavior <i>in the limit as N becomes large</i> , lower order terms may dominate at small input sizes.
80	8120	3.17	
160	28840	3.55	
320	108680	3.77	
640	421960	3.88	
1280	1662920	3.94	
2560	6602440	3.97	
5120	26311880	3.99	
10240	105052360	3.99	
20480	419819720	4.00	

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$n^2 + n\log(n) + (\log(n))^2$  is... \*

- ☒  $O(n^2)$
- ☐  $O(n\log(n))$
- ☐  $O((\log(n))^2)$

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$n^2 + 2^n$  is... \*

- ☐  $O(n^2)$
- ☒  $O(2^n)$

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$\log(n^2)$  is... \*  
 $= 2\log(n) = O(\log(n))$

- ☐  $O(n^2)$
- ☐  $O(n)$
- ☐  $O((\log(n))^2)$
- ☒  $O(\log(n))$

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Suppose you time an algorithm for different values of N and get the results shown in the table. What is the best characterization of the asymptotic runtime complexity observed in the data? \*

N	Time (s)
100	0.03
200	0.08
400	0.24
800	0.80
1600	2.87
3200	10.85
6400	42.18
12800	166.28

- ☐  $O(N^3)$
- ☒  $O(N^2)$
- ☐  $O(N)$
- ☐  $O(\log(N))$
- ☐  $O(1)$

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Suppose you time an algorithm for different values of N and M and get the results shown in the table. What is the best characterization of the asymptotic runtime complexity observed in the data? \*

N	M	Time (s)
100	100	0.81
100	200	1.21
100	400	2.01
200	100	1.11
200	200	1.51
200	400	2.31
400	100	1.71
400	200	2.11
400	400	2.91

- ☐  $O(N)$
- ☐  $O(M)$
- ☐  $O(NM)$
- ☒  $O(N+M)$

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### Big-Oh for Runtime: Algorithms & Code

- What is the runtime complexity of `stuff(n)`?
- How many times does the loop iterate?
  - In terms of `n`, the parameter
- Loop body is  $O(1)$ ?
  - Constant time
  - Independent of `n`
  - Add `n` same as add 1

```
7- public int stuff(int n) {
8   int sum = 0;
9   for(int k=0; k < n; k += 1) {
10    sum += n;
11  }
12  return sum;
13 }
```

Linear,  $O(n)$

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### General strategy for determining Big-O runtime complexity

Most general: Determine  $T(N)$ , the number of **constant time operations** as a function of the **size of the input,  $N$** . Then simplify using Big-O.

Practically, covers common cases:

1. For each line of code, label:
  - a) Complexity of that line, and
  - b) Number of times the line is executed
2. Add up over all lines, multiplying the two labels

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### Nested loop example

What about the big-O runtime complexity of this code as a function of  $n$ ?

6	public int nested(int n) {	Line	Complexity	Iterations
7	int result = 0;	7	$O(1)$	1
8	for (int i=0; i<n; i++) {	8	$O(1)$	$n$
9	for (int j=0; j<i; j++) {	9	$O(1)$	?
10	result += 1;	10	$O(1)$	?
11	}			
12	}	13	$O(1)$	1
13	return result;			

How many times does line 10 execute?

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### Nested loop example

How many times does line 10 execute?

6	public int nested(int n) {	when i is	Line 10 executes this many times
7	int result = 0;		
8	for (int i=0; i<n; i++) {	1	1
9	for (int j=0; j<i; j++) {	2	2
10	result += 1;	...	...
11	}	$n-2$	$n-2$
12	}	$n-1$	$n-1$
13	return result;		

In total?  $1 + 2 + \dots + (n - 2) + (n - 1) \approx \frac{n^2}{2}$  is  $O(n^2)$  iterations

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### Nested loop example

Putting it together:

		Line	Complexity	Iterations
6	public int nested(int n) {	7	O(1)	1
7	int result = 0;	8	O(1)	n
8	for (int i=0; i<n; i++) {	9	O(1)	O(n <sup>2</sup> )
9	for (int j=0; j<i; j++) {	10	O(1)	O(n <sup>2</sup> )
10	result += 1;	13	O(1)	1
11	}			
12	}			
13	return result;			

Total runtime complexity: (1) + (n) + (n<sup>2</sup>) + (n<sup>2</sup>) + (1)  
is O(n<sup>2</sup>)

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### Not all nested loops are quadratic

What about the big-O runtime complexity of this code as a function of n?

		Line	Complexity	Iterations
16	public int nested2(int n) {	17	O(1)	1
17	int result = 0;	18	O(1)	n
18	for (int i=0; i<n; i++) {	19	O(1)	100n
19	for (int j=0; j<100; j++)	20	O(1)	100n
20	result += 1;	23	O(1)	1
21	}			
22	}			
23	return result;			
24	}			

Total runtime complexity: (1) + (n) + (200n) + (1)  
is O(n)

Reminder: 200n is 200 times slower than n, but their runtimes both scale linearly

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### Not all loops are nested

What about the big-O runtime complexity of this code as a function of n?

		Line	Complexity	Iterations
28	public int parallel(int n) {	29	O(1)	1
29	int result = 0;	30	O(1)	n
30	for (int i=0; i<n; i++) {	31	O(1)	n
31	result += 1;	33	O(1)	n
32	}	34	O(1)	n
33	for (int i=0; i<n; i++) {	36	O(1)	1
34	result += 1;			
35	}			
36	return result;			
37	}			

Total runtime complexity: (1) + (4n) + (1)  
is O(n)

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## Not all loops increment by 1

Big-O Runtime complexity of calc(N) is...

- How many times does the loop iterate?
  - Concrete to abstract: calc(16), calc(32), ...
- Inside loop?  $O(1)$  operations

```

143 public int calc(int n) {
144     int sum = 0;
145     for(int k=1; k < n; k *= 2) {
146         sum += k;
147     }
148     return sum;
149 }

```

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## Generalizing: Concrete to Abstract

N	# loop iterations
1	0
2	1
4	k=1,2...2 iters
8	k=1,2,4...3 iters

N	# loop iterations
16	5 iterations
32	k=1,2,4,8,16..5 iters
33	6 iterations
63	6 iterations

```

143 public int calc(int n) {
144     int sum = 0;
145     for(int k=1; k < n; k *= 2) {
146         sum += k;
147     }
148     return sum;
149 }

```

$O(\log(N))$

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## Accounting for iteration and non-constant time operations

What about the big-O runtime complexity of this code as a function of  $n = \text{words.size}()$ ?

```

40 public ArrayList<String> uniqueWords (ArrayList<String> words) {
41     ArrayList<String> unique = new ArrayList<>();
42     for (String w : words) {
43         if (!unique.contains(w)) {
44             unique.add(w);
45         }
46     }
47     return words;
48 }

```

Annotations:

- Loop n times...
- $O(1)$ , initialized empty List
- Amortized  $O(1)$ , add to end
- $O(n)$  in worst case
- Still  $O(1)$ , does not copy list to return

Total: Make  $n$  calls to  $O(n)$  contains:  $O(n^2)$

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## Exponential time algorithm?

Problem from previous WOTO: What is the runtime complexity of `concatA lot` as a function of `reps`?

```

12 public static String concatA lot(int reps, String s) {
13     for (int i=0; i<reps; i++) {
14         s += s;
15     }
16     return s;
17 }

```

Loop reps times

Same as:  
`s = s + s;`

Runtime of line 14 is  $O(s.length())$ . And this doubles every iteration through the loop.

Examine how the length of `s` grows by iterations.

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## Exponential time algorithm?

```

12 public static String concatA lot(int reps, String s) {
13     for (int i=0; i<reps; i++) {
14         s += s;
15     }
16     return s;
17 }

```

Loop reps times

Same as:  
`s = s + s;`

Examine how the length of `s` grows by iterations.

Iteration	<code>s.length()</code>	$O(1)$ operations (char copies)
0 (input <code>s</code> )	1 (suppose)	2
1	2	4
2	4	8
3	8	16
...	...	...
<code>reps-1</code>	$2^{\text{reps}-1}$	$(2)(2^{\text{reps}-1}) = 2^{\text{reps}}$

So runtime has to be at least  $2^{\text{reps}}$ , exponential complexity!

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2. What is the big  $O$  runtime complexity of the `keepHalving` method as a function of the parameter `n`? \*

```

102 public int keepHalving(int n) {
103     int numIterations = 0;
104     while (n > 1) {
105         n = n / 2;
106         numIterations++;
107     }
108     return numIterations;
109 }

```

- ☐  $O(1)$
- ☒  $O(\log(n))$
- ☐  $O(n)$
- ☐  $O(n^2)$
- ☐  $O(n^3)$
- ☐  $O(2^n)$

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What is the big O runtime complexity of the moreLooping method as a function of the parameter n? \*

```
86 public int moreLooping(int n) {
87     int result = 0;
88     for (int i=n-1; i<n; i++) {
89         for (int k=0; k<10; k++) {
90             result += 1;
91         }
92     }
93     return result;
94 }
```

- ☒ O(1)
- ☐ O(log(n))
- ☐ O(n)
- ☐ O(n^2)
- ☐ O(n^3)
- ☐ O(2^n)

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What is the big O runtime complexity of the reverse method as a function of n where n is the size of the List parameter input? add(0, s) adds s to the front of the list. \*

```
94 public static List<String> reverse(List<String> input) {
95     ArrayList<String> result = new ArrayList<>();
96     for (String s : input) {
97         result.add(0, s);
98     }
99     return result;
100 }
```

- ☐ O(1)
- ☐ O(log(n))
- ☐ O(n)
- ☒ O(n^2)
- ☐ O(n^3)
- ☐ O(2^n)

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